WELLESLEY

Practice problems

For Exam 2: ISA

Struct practice problem (similar to CSAPP 3.45)

ex

```
struct s {
  char *a;
  short b;
  int *c;
  char d;
  int e;
  char f;
};
```

1. Draw a picture of how this struct is laid out in memory, labeling the byte offset of each field (starting with `a` at offset `+0`);

2. Modify your picture to show how much space a single element of this struct would take if used as an element of an array (e.g., the total size).

Recall: a short is 2 bytes in C

3. Rearrange the fields of the struct to minimize wasted space. Draw the new offsets and the total size.

Struct practice problem (similar to CSAPP 3.45)

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struct s {
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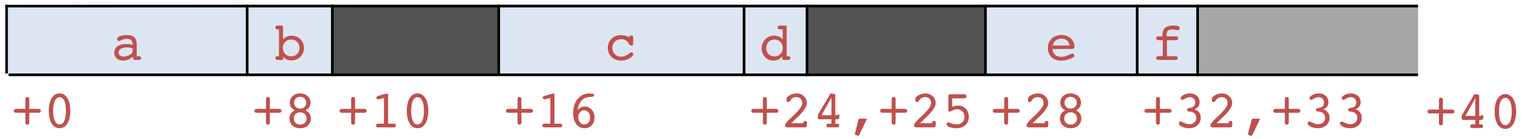

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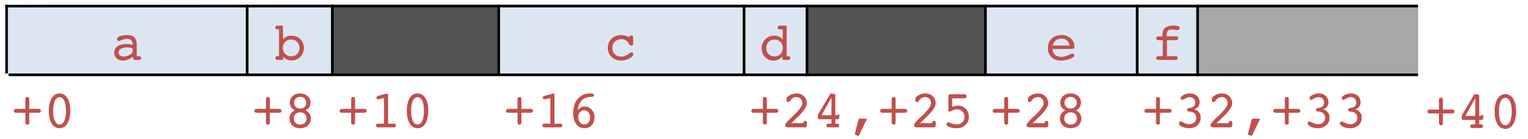

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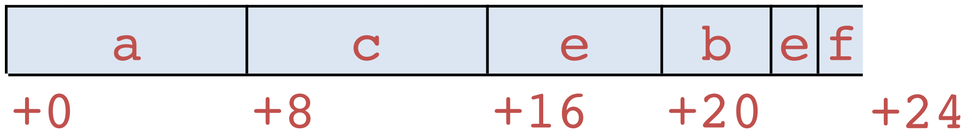

2-D array practice problem

ex

```
long a[2][3];
```

1. Draw a picture of how this array is laid out in memory, labeling the indices and byte offset of each element (starting with `a[0][0]` at offset `+0`);

Recall: `index = C*r + c`

scale by element size

```
long get_elem_1_2(long a[2][3]){
  return a[1][2];
}
```

2. Write x86 assembly code to implement this function.

2-D array practice problem: solution

ex

```
long a[2][3];
```

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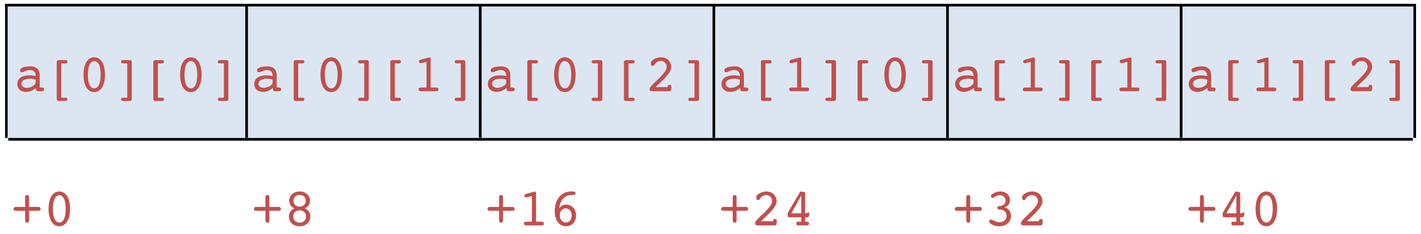


```
long get_elem_1_2(long a[2][3]){
  return a[1][2];
}
```

2. Write x86 assembly code to implement this function.

Since we know the size, we can calculate
C*r+c = 3*1+2 = 5, 5*sizeof(long) = 5*8 = 40

```
movq 40(%rdi),%rax
retq
```

x86 struct/LinkedList practice problem

ex

```
nodeFunc2:
        pushq   %rbp
        pushq   %rbx
        subq    $8, %rsp
        movl    %esi, %ebx
        movslq  %esi, %rax
        testq   %rdi, %rdi
        je      .L1
        movq    %rdi, %rbp
        movl    8(%rdi), %esi
        cmpl    %esi, %ebx
        jb      .L5
.L3:    movq    0(%rbp), %rdi
        movl    %ebx, %esi
        call    nodeFunc2
.L1:    addq    $8, %rsp
        popq    %rbx
        popq    %rbp
        ret
.L5:    movl    %esi, %ebx
        jmp     .L3
```

```
typedef struct Node {
  struct Node* next;
  unsigned int value;
} Node;

long nodeFunc2(Node* node, unsigned int x) {
    // ???
}

long nodeFunc1(Node* node) {
    nodeFunc2(node, 0);
}
```

Consider the above function that calculates something useful about a linked list of unsigned integers using a helper function.

1. Identify which pieces of x86 refer to `next` and `value`.
2. Identify the base case of the recursive function `nodeFunc2`. What is returned in this case?
3. Identify the recursive case of `nodeFunc2`. What is the argument passed to the recursive call?
4. What is `nodeFunc1` calculating with helper `nodeFunc2`?

x86 struct/LinkedList practice problem

ex

```
nodeFunc2:
        pushq   %rbp
        pushq   %rbx
        subq    $8, %rsp
        movl    %esi, %ebx
        movslq  %esi, %rax
        testq   %rdi, %rdi
        je      .L1
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        movl    8(%rdi), %esi
        cmpl    %esi, %ebx
        jb      .L5
.L3:    movq    0(%rbp), %rdi
        movl    %ebx, %esi
        call    nodeFunc2
.L1:    addq    $8, %rsp
        popq    %rbx
        popq    %rbp
        ret
.L5:    movl    %esi, %ebx
        jmp     .L3
```

At call, `%rsp` must be a multiple of 16

node = `%rdi`

base case

recursive case

```
typedef struct Node {
  struct Node* next;
  unsigned int value;
} Node;
long nodeFunc2(Node* node, unsigned int max) {
  if (node == 0) {
    return max;
  }
  if (node->value > max) {
    max = node->value;
  }
  nodeFunc2(node->next, max);
}
long nodeFunc1(Node* node) {
    nodeFunc2(node, 0);
}
```

`8(%rdi)` accesses `node->value`, `(%rdi)` accesses `node->next`,
if `(node->value > x)`, jump to `.L5`, sets `%ebx` to `node->value`
`%ebx` calculates the max of `node->value` and `x`

in the base case, returns second arg, `x` (the maximum value found so far)

`nodeFunc1` uses its helper to **find the maximum value** within a linked list.

x86 recursive procedure practice problem

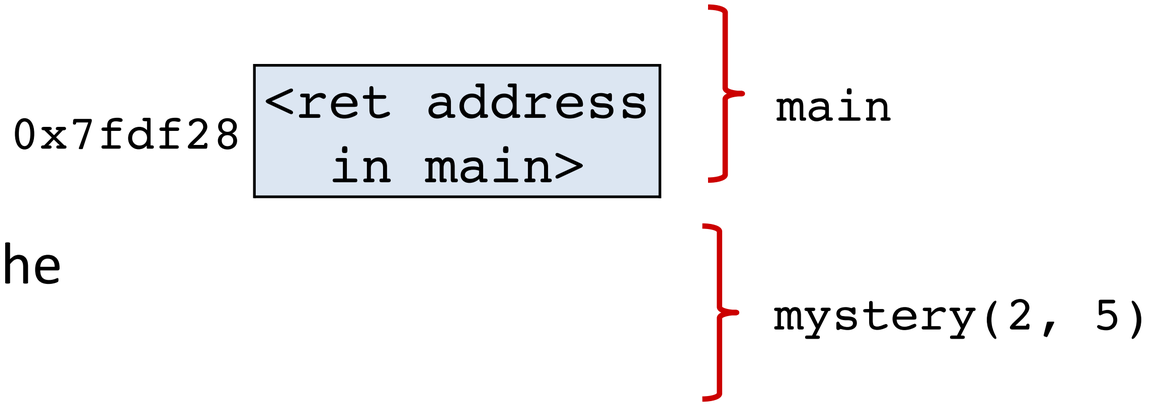
ex

```
        mystery:
401106 mov    $0x0,%eax
40110b test   %edi,%edi
40110d jne    401110 <mystery+0xa>
40110f ret
401110 push   %rbx
401111 mov    %esi,%ebx
401113 sub    $0x1,%edi
401116 call   401106 <mystery>
40111b movslq %ebx,%rsi
40111e add    %rsi,%rax
401121 pop    %rbx
401122 ret
```

1. What registers is being saved to the stack? Why?

2. What instruction address gets saved to the stack? Why?

3. What is this function computing?



4. Fill in the top of this stack after the function returns to main for `mystery(2, 5)`.

What is each value returned, in order?

x86 recursive procedure practice problem

ex

```
        mystery:
401106  mov     $0x0,%eax
40110b  test    %edi,%edi
40110d  jne     401110 <mystery+0xa>
40110f  ret
401110  push    %rbx
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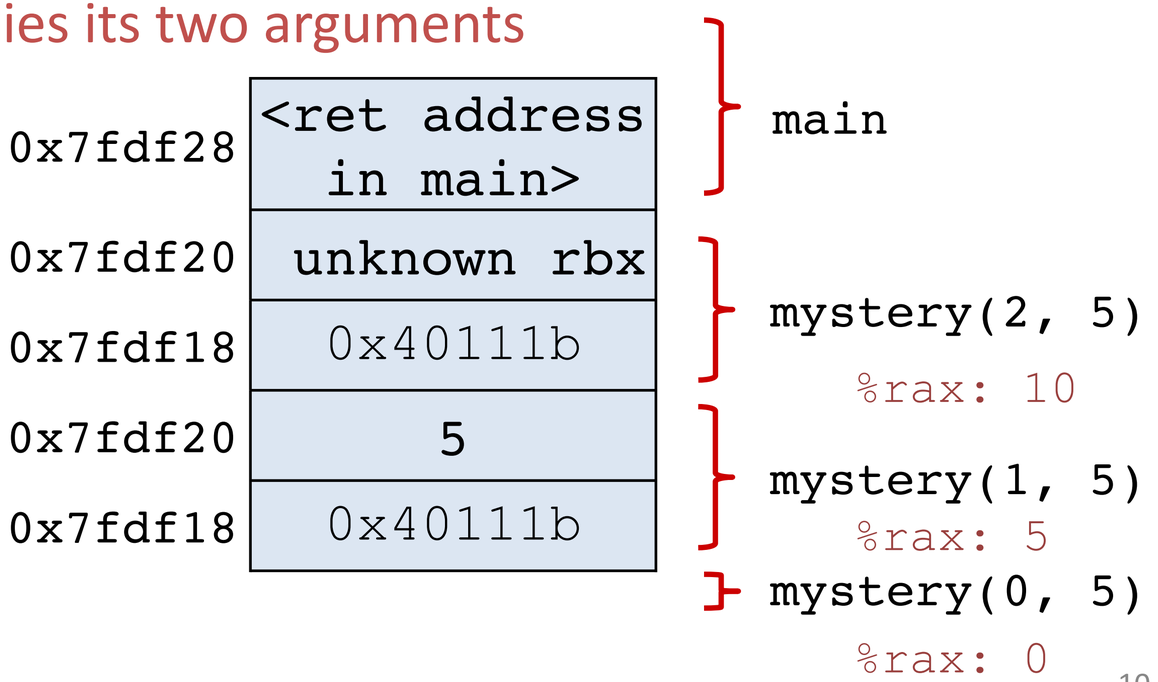
401116  call    401106 <mystery>
40111b  movslq  %ebx,%rsi
40111e  add     %rsi,%rax
401121  pop     %rbx
401122  ret
```

1. What registers is being saved to the stack? Why?
   %rbx, so that it is not overwritten in the recursive call
2. What instruction address gets saved to the stack? Why?
   0x40111b, return address after recursive call
3. What is this function computing?
   Multiplies its two arguments

```
int mult(int x, int y) {
    if (x == 0)  return 0;
    return y + mult(x - 1, y);
}
```

%rax: 0



x86 short answer practice problems

ex

1. Which x86 instructions implicitly change the stack pointer? How do they change it?

2. What are some things defined by the *word size*? What is the word size we have been using for x86 in class?

3. Describe the general idea of a buffer overflow exploit in C code compiled to x86.

4. Describe how a child process's memory is related to the memory of the parent process.

x86 short answer practice problems

ex

1. Which x86 instructions implicitly change the stack pointer? How do they change it?

pushq	popq	call	ret
`%rsp -= 8`	`%rsp += 8`	`%rsp -= 8`	`%rsp += 8`

2. What are some things defined by the *word size*? What is the word size we have been using for x86 in class?

Register size, address size, pointer size

NOT instruction size (variable-width instruction size)

3. Describe the general idea of a buffer overflow exploit in C code compiled to x86.

Buffer overflow occurs when code lacks bounds checking in writing untrusted input to a destination region of memory that is too small. Buffer overflow attacks can overwrite the return addresses on the stack to point to further exploit code.

4. Describe how a child process's memory is related to the memory of the parent process.

The child process starts with a copy of the state of the parent's memory. It is a private copy: the child and the parent do not share memory once the child is created.

x86 arithmetic practice problem

ex

```
long funmath0(long x, long y) {
    return x + 4*y + 21;
}
```

```
long funmath1(long x, long y) {
    return 2*x + 4*y + 21;
}
```

```
long funmath2(long x, long y) {
    return 6*x + 5*y + 21;
}
```

Implement the above functions in x86 *without* `addq` or `mulq`.

You can use `leaq` and any other x86 instruction.

Recall: addressing modes can only multiply by 1, 2, 4, or 8.

x86 arithmetic practice problem

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    return x + 4*y + 21;
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}
```

```
long funmath2(long x, long y) {
    return 6*x + 5*y + 21;
}
```

3 possible answers:

```
funmath0:
  leaq    21(%rdi,%rsi,4), %rax
  ret

funmath1:
  leaq    (%rdi,%rsi,2), %rax
  leaq    21(%rax,%rax), %rax
  ret

funmath2:
  leaq    (%rdi,%rdi,2), %rdx
  leaq    (%rsi,%rsi,4), %rax
  leaq    21(%rax,%rdx,2), %rax
  ret
```

Implement the above functions in x86 *without* `addq` or `mulq`.

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