



# CS 240 Stage 2!

## Hardware-Software Interface

Memory addressing, C language, pointers

Assertions, debugging

Machine code, assembly language, program translation

Control flow

Procedures, stacks

Data layout, security, linking and loading



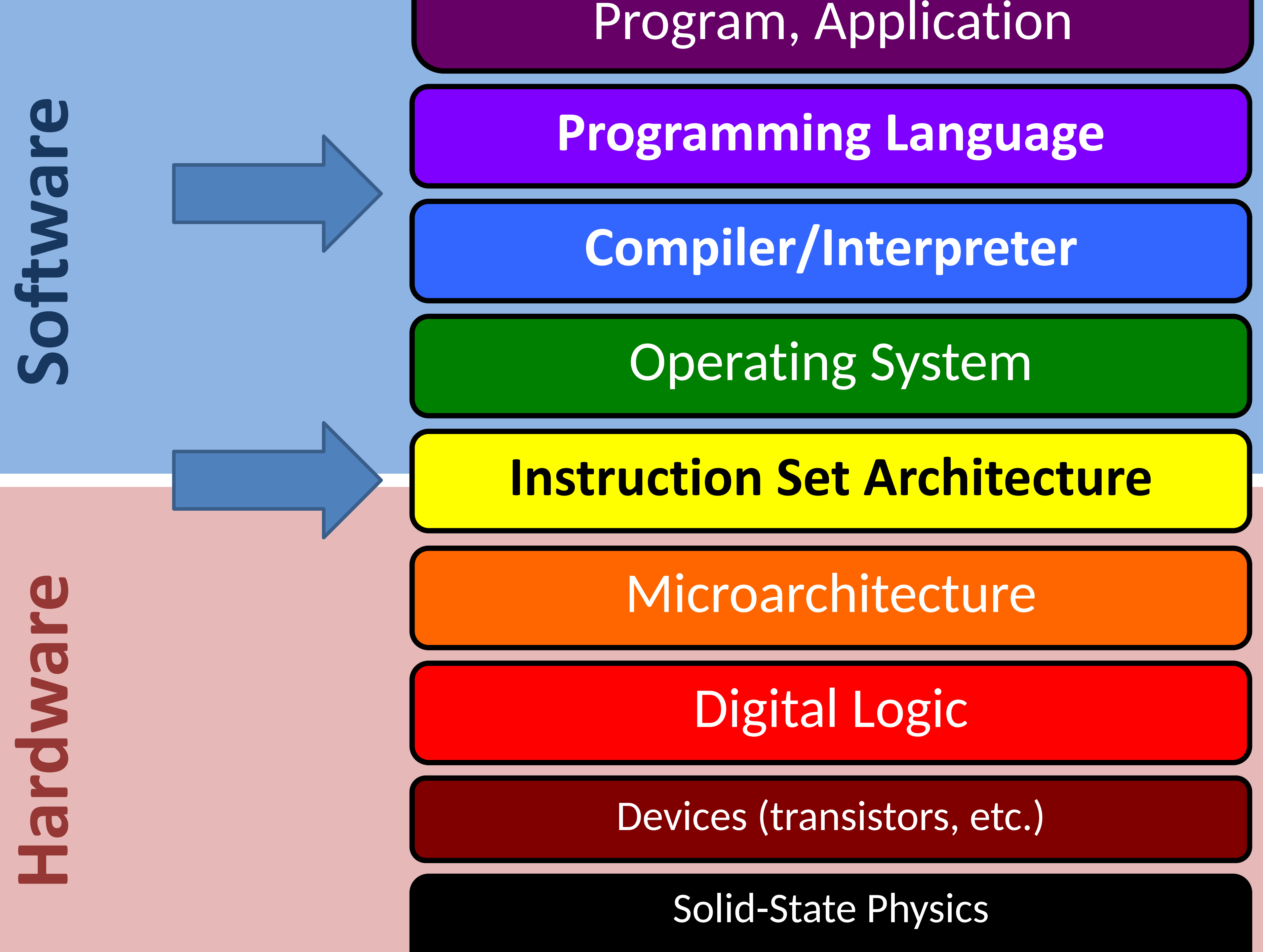
**CS 240**

Foundations of Computer Systems

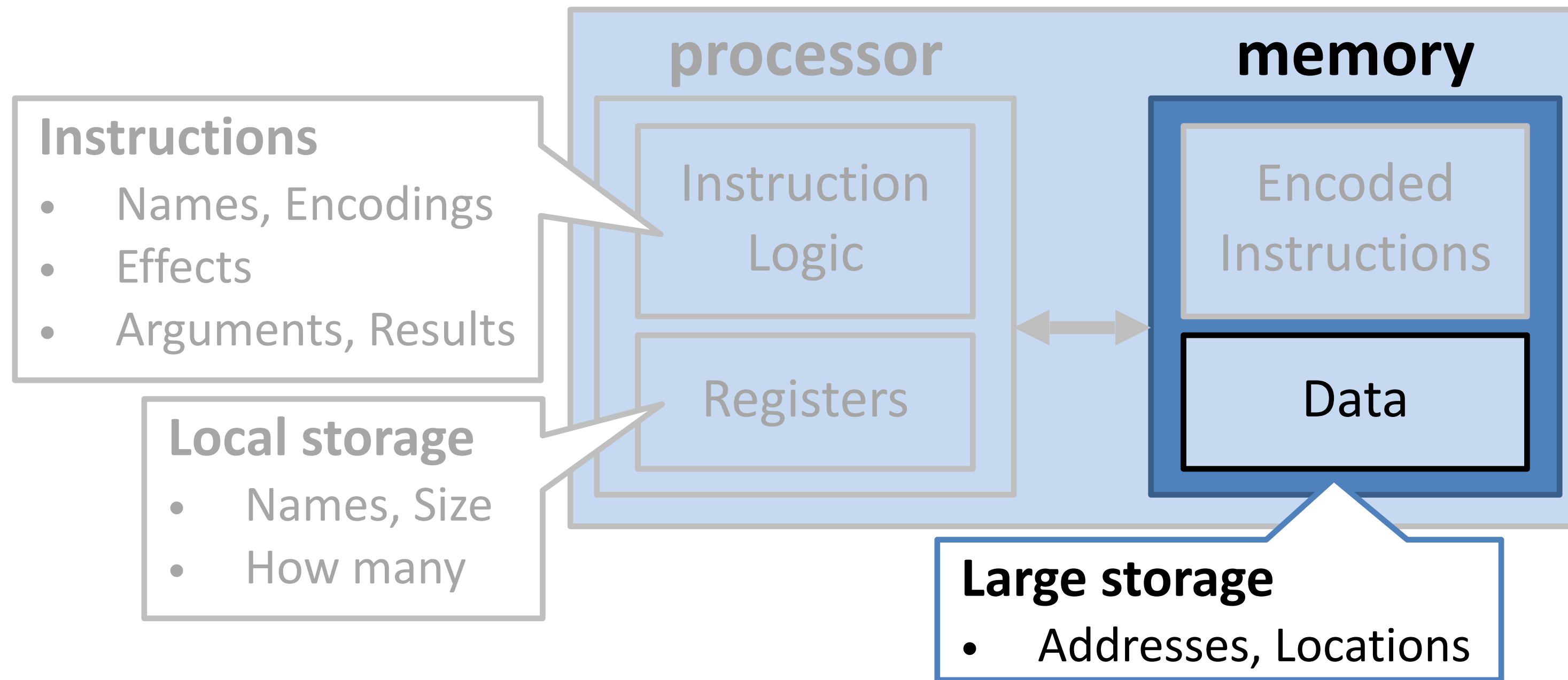


# Programming with Memory

the memory model  
pointers and arrays in C



# Instruction Set Architecture (HW/SW **Interface**)

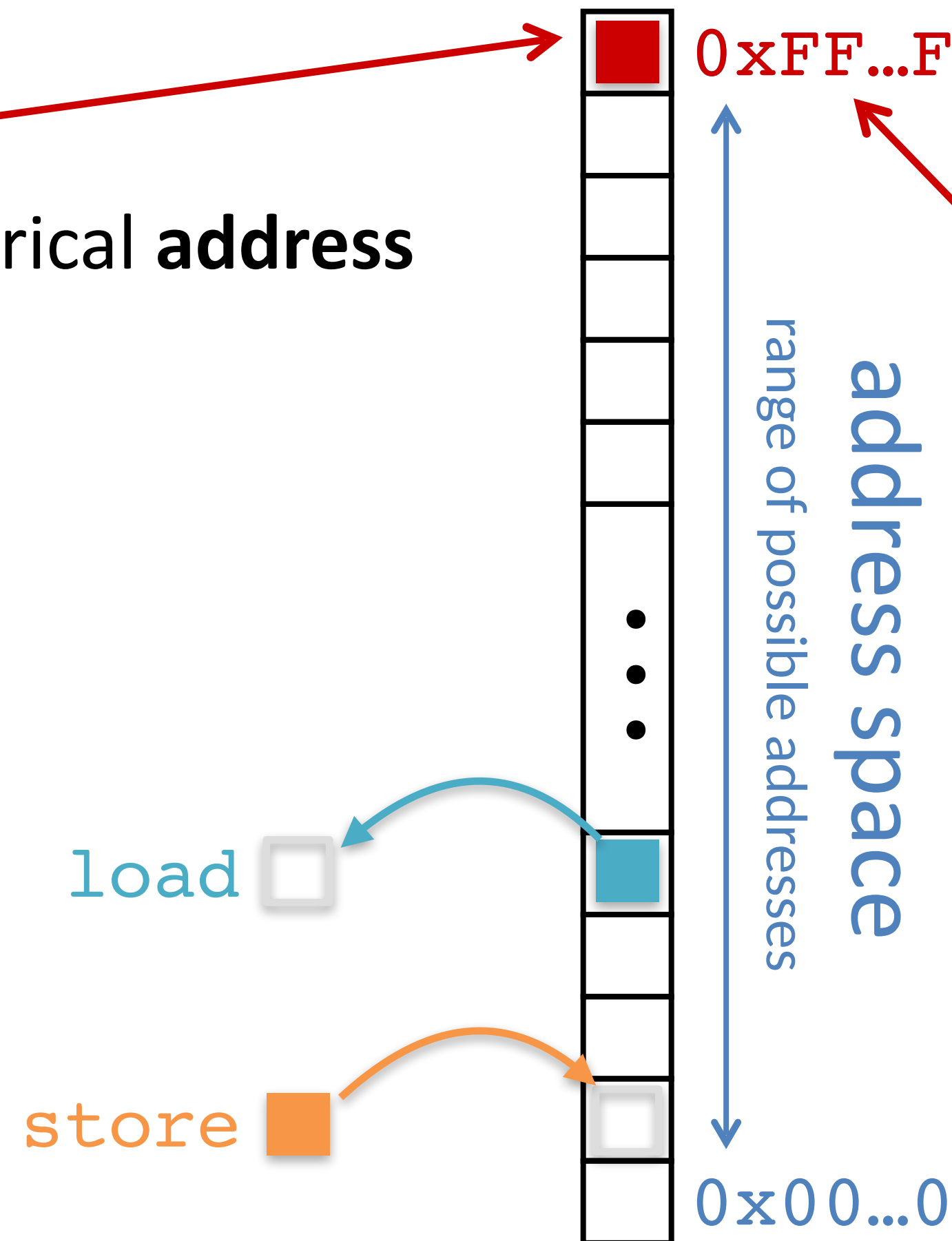


**Computer**

# Byte-addressable memory = mutable byte array

**Location / cell = element**

- Identified by unique numerical **address**
- Holds **one byte (8 bits)**



**Address = index**

- Unsigned number
- Represented by one word
- Computable and storable as a value

**Operations:**

- **Load:** read contents at given address
- **Store:** write contents at given address

# Multi-byte values in memory

Store across contiguous byte locations.

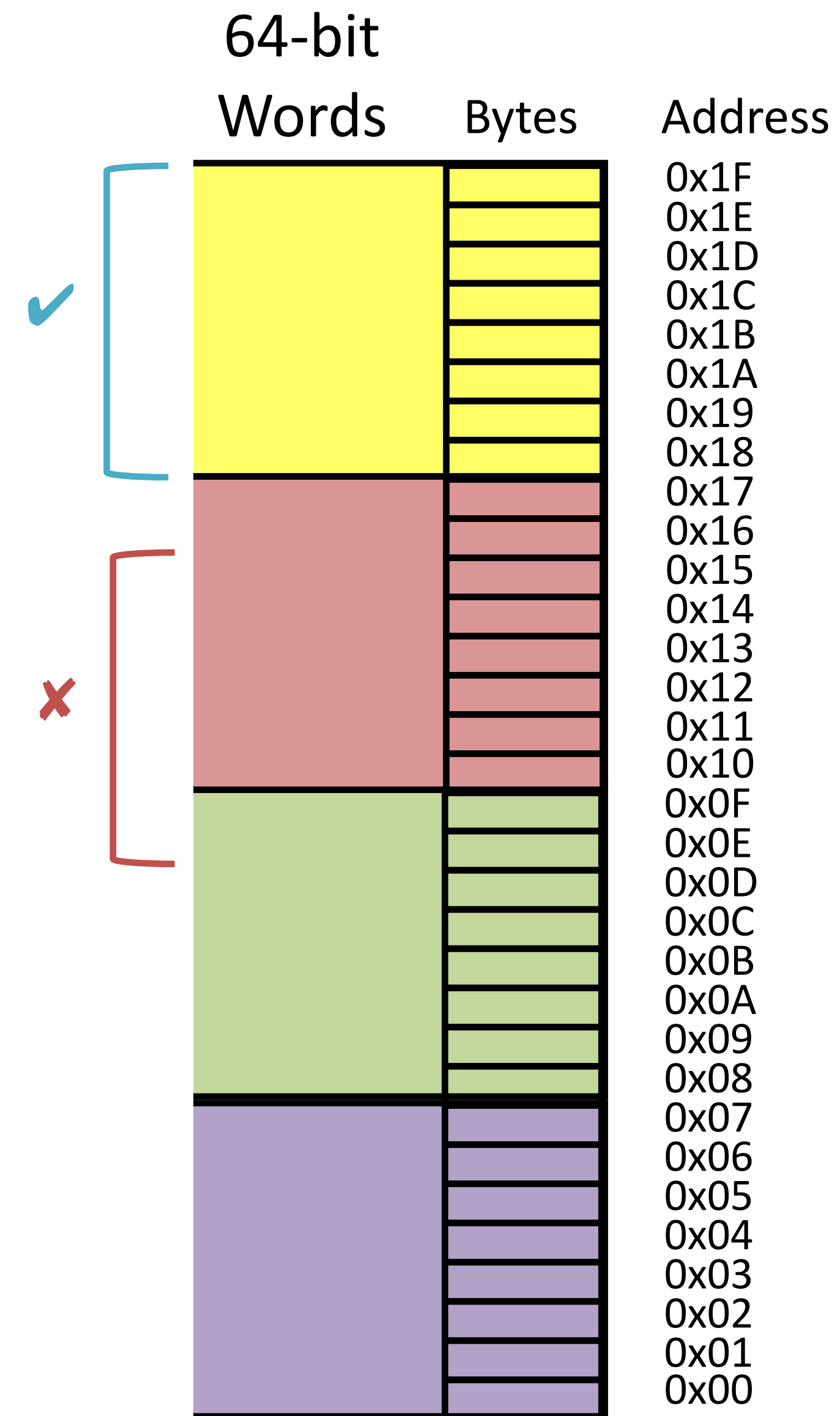
Example: 8 byte (64 bit) values

## Alignment

Multi-byte values start at addresses that are multiples of their size

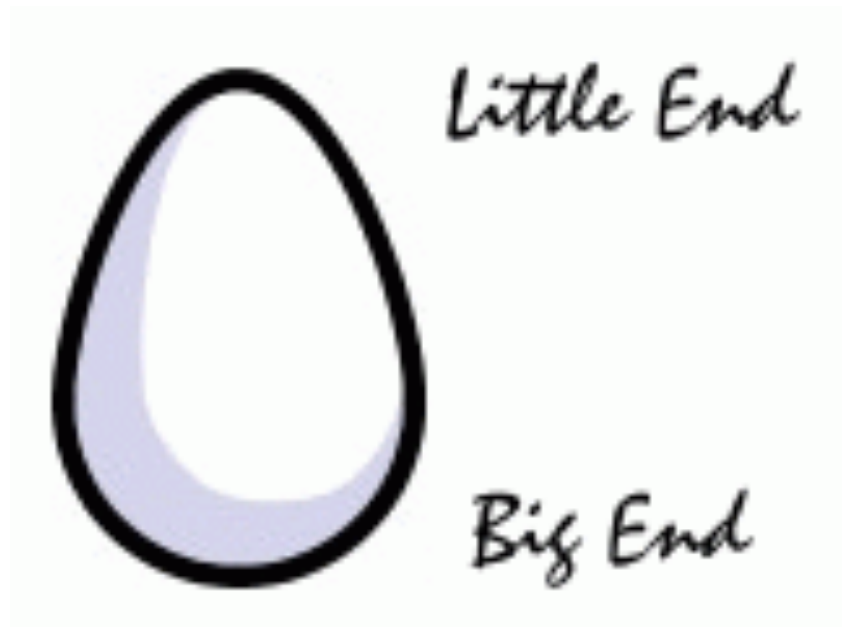
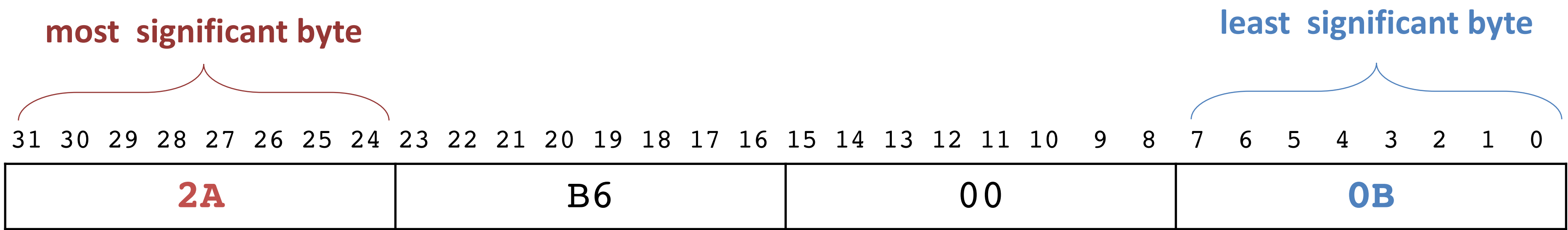
Bit order within byte always same.

Recall: byte ordering within larger value?



# Endianness: details

In what order are the individual bytes of a multi-byte value stored in memory?



Address	Contents
03	2A
02	B6
01	00
00	0B

**Little Endian:** least significant byte first

- low order byte at low address
- high order byte at high address
- used by **x86**, ... and **CS240**!

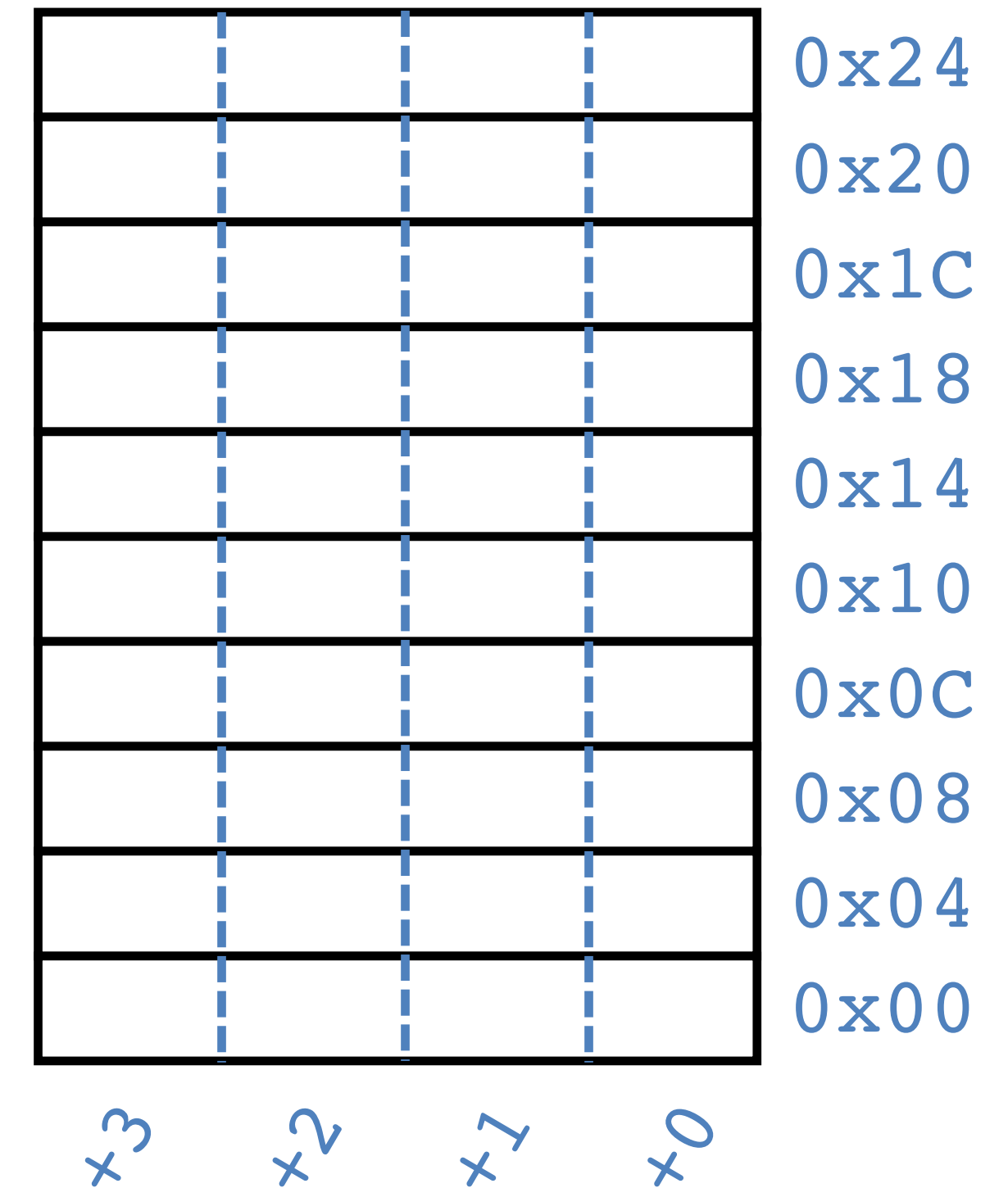
Address	Contents
03	0B
02	00
01	B6
00	2A

**Big Endian:** most significant byte first

- high order byte at low address
- low order byte at high address
- used by networks, SPARC, ...

# Data, addresses, and pointers

For these slides, we'll draw the bytes in this reverse order so that multi-byte values can be read directly



memory drawn as 32-bit values,  
little endian order



# Data, addresses, and pointers

**address** = index of a location in memory

**pointer** = a reference to a location in memory,  
represented as an address stored as data

Let's store the number 240 at address **0x20**.

$$240_{10} = F0_{16} = 0x00\ 00\ 00\ F0$$

At address **0x08** we store a pointer to the contents at address **0x20**.

At address **0x00**, we store a pointer to a pointer.

The number 12 is stored at address **0x10**.

Is it a pointer?

How do we know if values are pointers or not?

How do we manage use of memory?

				0x24
00	00	00	F0	0x20
				0x1C
				0x18
				0x14
00	00	00	0C	0x10
				0x0C
00	00	00	20	0x08
				0x04
00	00	00	08	0x00
3	2	1	0	

memory drawn as 32-bit values,  
little endian order

# C: Variables are locations

The compiler creates a map from variable name → location.

Declarations do not initialize!

```
int x; // x @ 0x20
```

```
int y; // y @ 0x0C
```

```
x = 0; // store 0 @ 0x20
```

```
// store 0x3CD02700 @ 0x0C
```

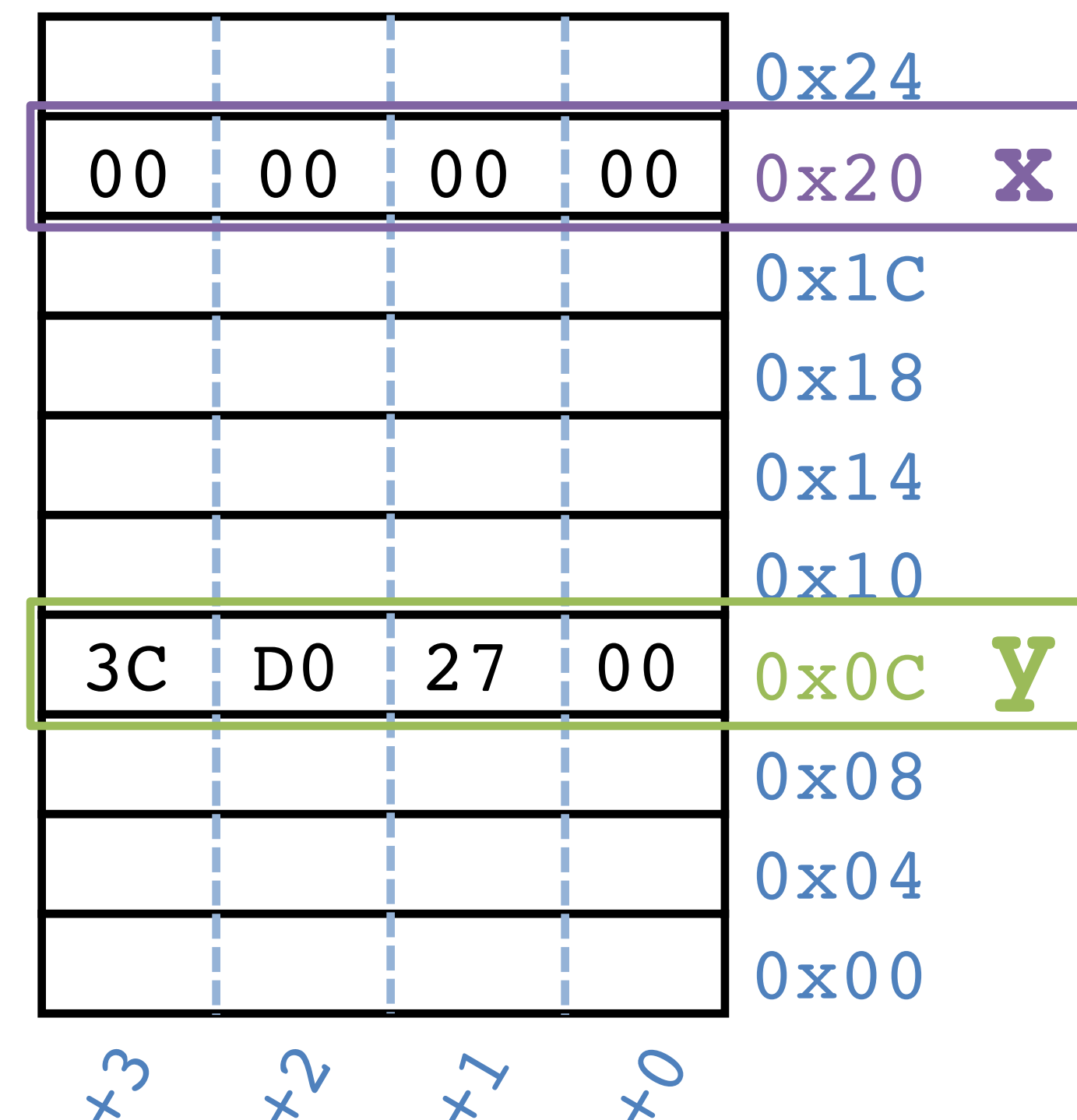
```
y = 0x3CD02700;
```

```
// 1. load the contents @ 0x0C
```

```
// 2. add 3
```

```
// 3. store sum @ 0x20
```

```
x = y + 3;
```



# C: Variables are locations

The compiler creates a map from variable name → location.

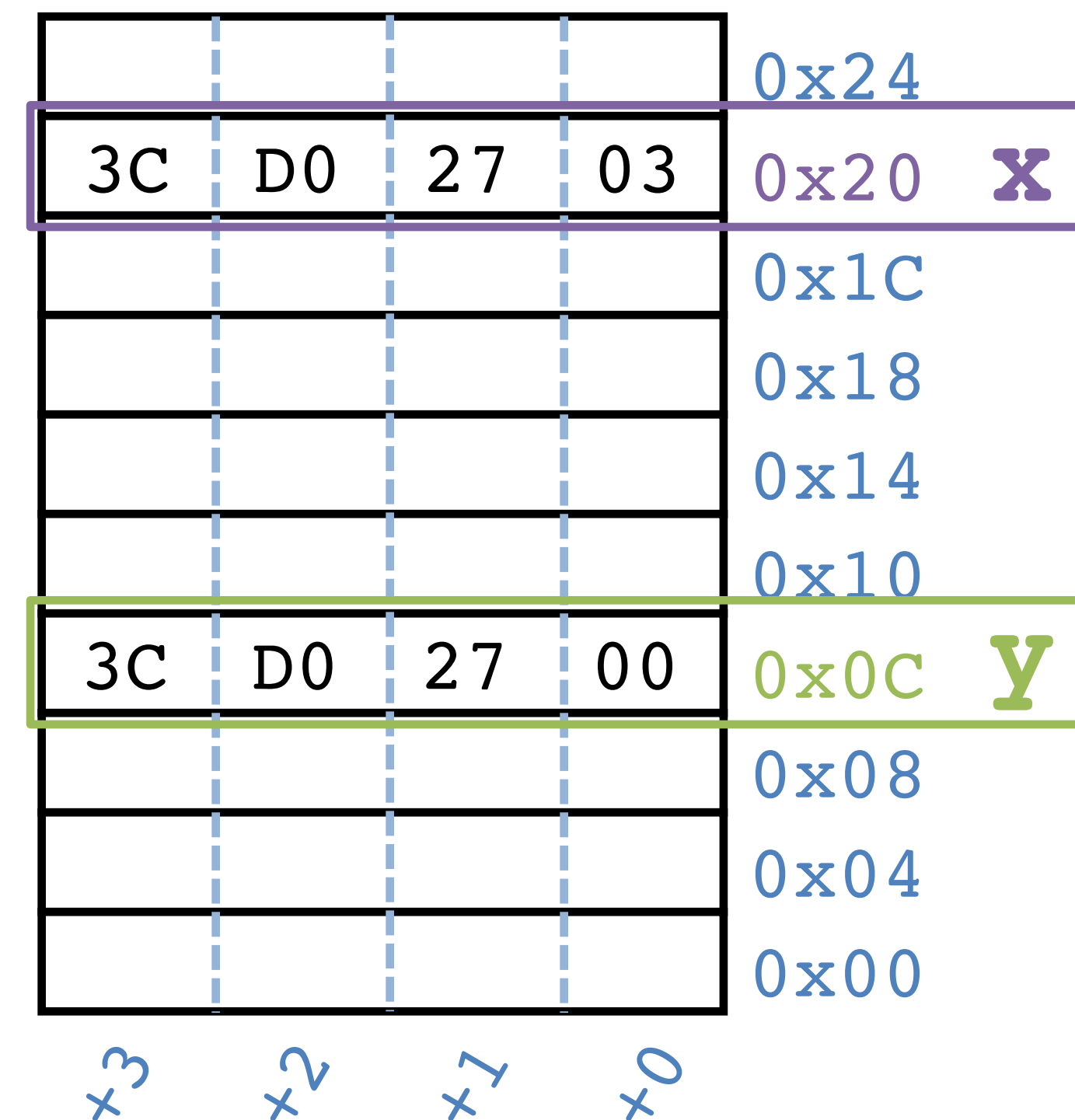
Declarations do not initialize!

```
int x; // x @ 0x20
int y; // y @ 0x0C

x = 0; // store 0 @ 0x20

// store 0x3CD02700 @ 0x0C
y = 0x3CD02700;

// 1. load the contents @ 0x0C
// 2. add 3
// 3. store sum @ 0x20
x = y + 3;
```



# C: Pointer operations and types

*address* = index of a location in memory

*pointer* = a *reference* to a location in memory, an address stored as data

## Expressions using addresses and pointers:

&\_\_\_\_ **address of** the memory location representing \_\_\_\_  
a.k.a. "reference to \_\_\_\_"

\*\_\_\_\_ **contents at** the memory address given by \_\_\_\_  
a.k.a. "dereference \_\_\_\_"

## Pointer types:

\_\_\_\_\* address of a memory location holding a \_\_\_\_  
a.k.a. "a reference to a \_\_\_\_"

# C: Types determine sizes

## Sizes of data types (in bytes)

Java Data Type	C Data Type
boolean	<i>bool</i>
byte	char
char	
short	short int
int	int
float	float
	long int
double	double
long	long long
	long double
(reference)	(pointer) *

Used by CS Linux, most modern machines

32-bit word	64-bit word
1	1
1	1
2	2
2	2
4	4
4	4
4	8
8	8
8	8
8	16
4	8

address size = word size

# C: Pointer example

*& = address of*  
*\* = contents at*

```
int* p;
```

Declare a variable, p

that will hold the address of a memory location holding an int

```
int x = 5;
```

```
int y = 2;
```

Declare two variables, x and y, that hold ints, and store 5 and 2 in them, respectively.

```
p = &x;
```

Take the address of the memory

representing x

... and store it in the memory location representing p.  
Now, "*p points to x.*"

Add 1 to

the contents of memory at the address

```
y = 1 + *p;
```

given by the contents of the  
memory location representing p

... and store it in the memory location representing y.

# C: Pointer example

& = address of  
\* = contents at

location

C assignment:

*Left-hand-side = right-hand-side;*

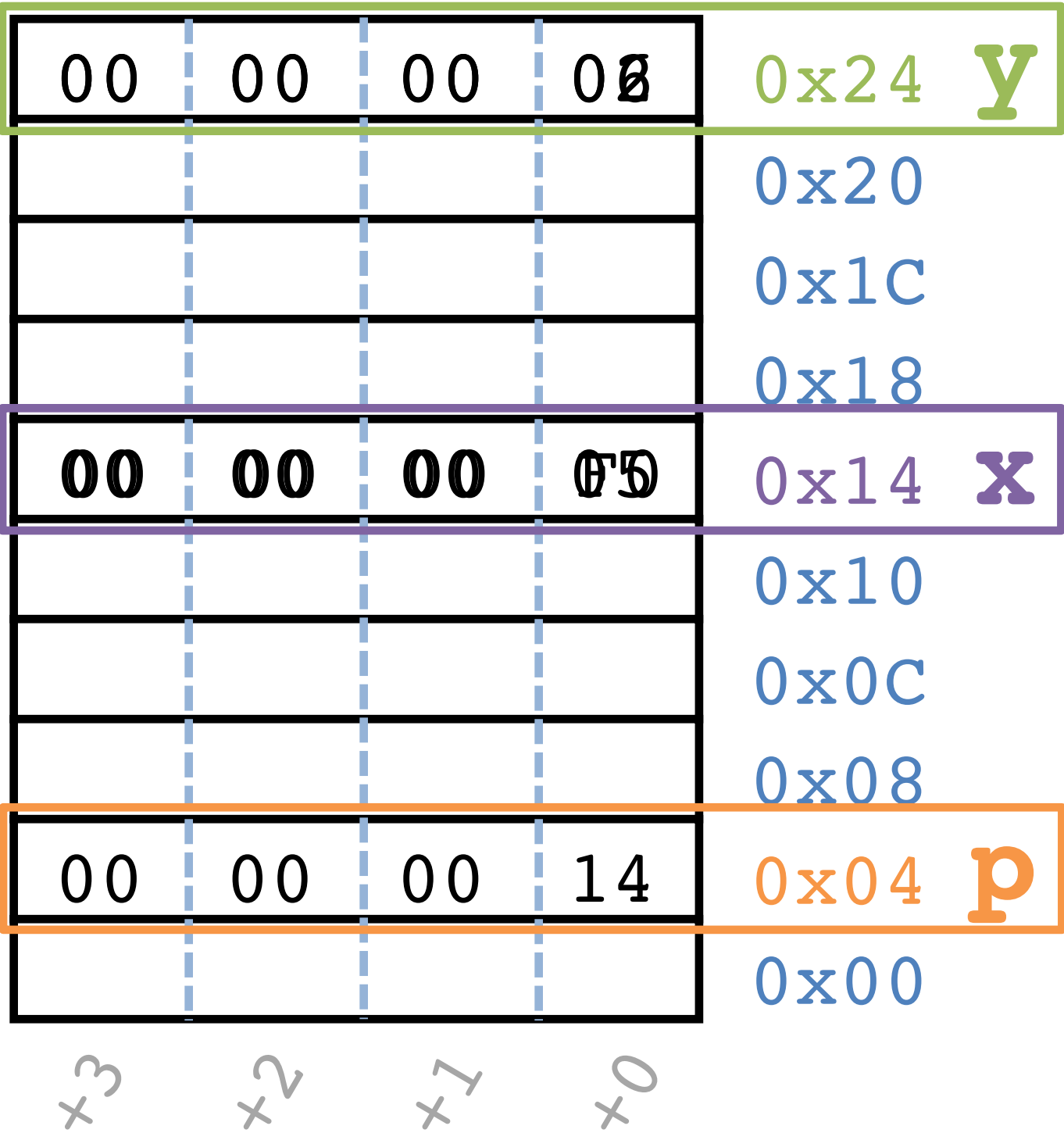
value

What is the type of \*p?  
What is the type of &x?  
What is \*(&y) ?

```
int* p; // p @ 0x04
int x = 5; // x @ 0x14, store 5 @ 0x14
int y = 2; // y @ 0x24, store 2 @ 0x24
p = &x; // store 0x14 @ 0x04

// 1. load the contents @ 0x04 (=0x14)
// 2. load the contents @ 0x14 (=0x5)
// 3. add 1
// 4. store sum as contents @ 0x24
y = 1 + *p;

// 1. load the contents @ 0x04 (=0x14)
// 2. store 0xF0 as contents @ 0x14
*p = 240;
```





# C: Pointer type syntax

Spaces between base type, `*`, and variable name mostly do not matter.

The following are **equivalent**:

```
int* ptr;
```

I see: "The variable **ptr** holds an **address of an int** in memory."

```
int * ptr;
```

```
int *ptr;
```

more common C style

Looks like: "Dereferencing the variable **ptr** will yield an **int**."

Or "The **memory location** where the variable **ptr** points holds an **int**."

Caveat: do not declare multiple variables unless using the last form.

```
int* a, b; means int *a, b; means int* a; int b;
```



# C: Arrays

Declaration:

```
int a[6];
```

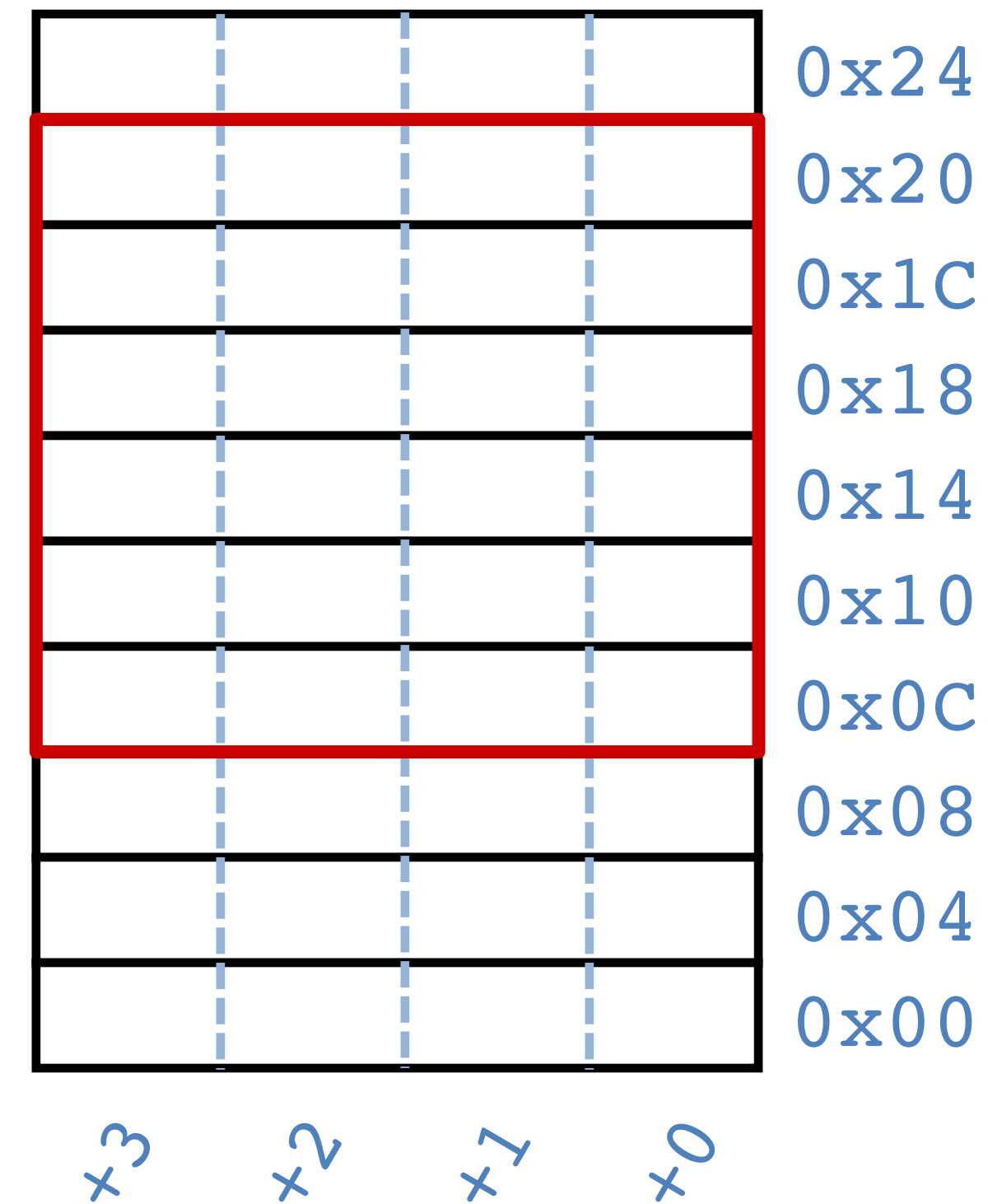
element type

name

number of  
elements

Arrays are adjacent memory locations storing the same type of data.

**a** is a name for the array's base address, can be used as an *immutable* pointer.



# C: Arrays

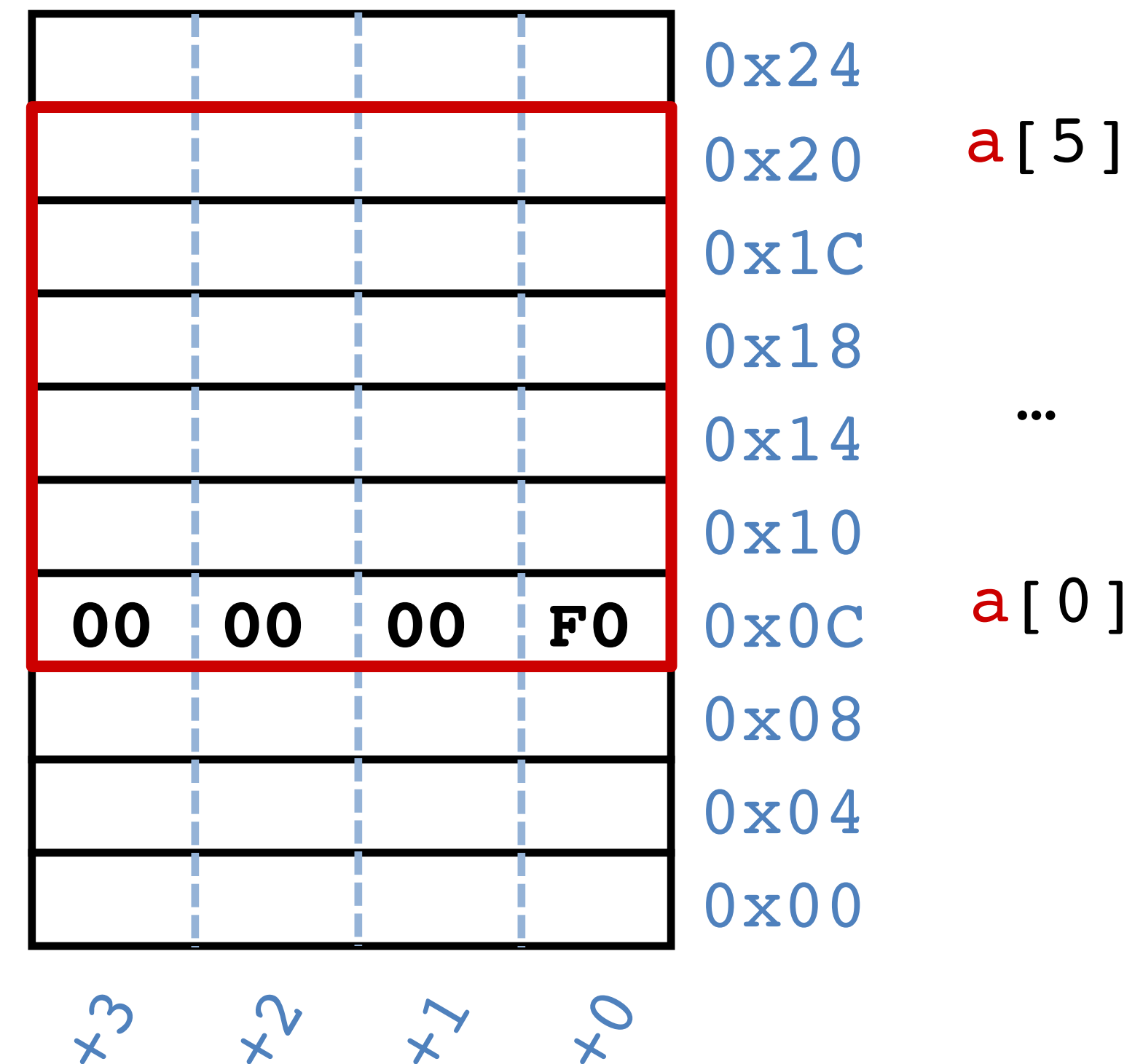
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Indexing: `a[0] = 0xf0;`

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Address of `a[i]` is base address `a` plus `i` times element size in bytes.



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```
int a[6];
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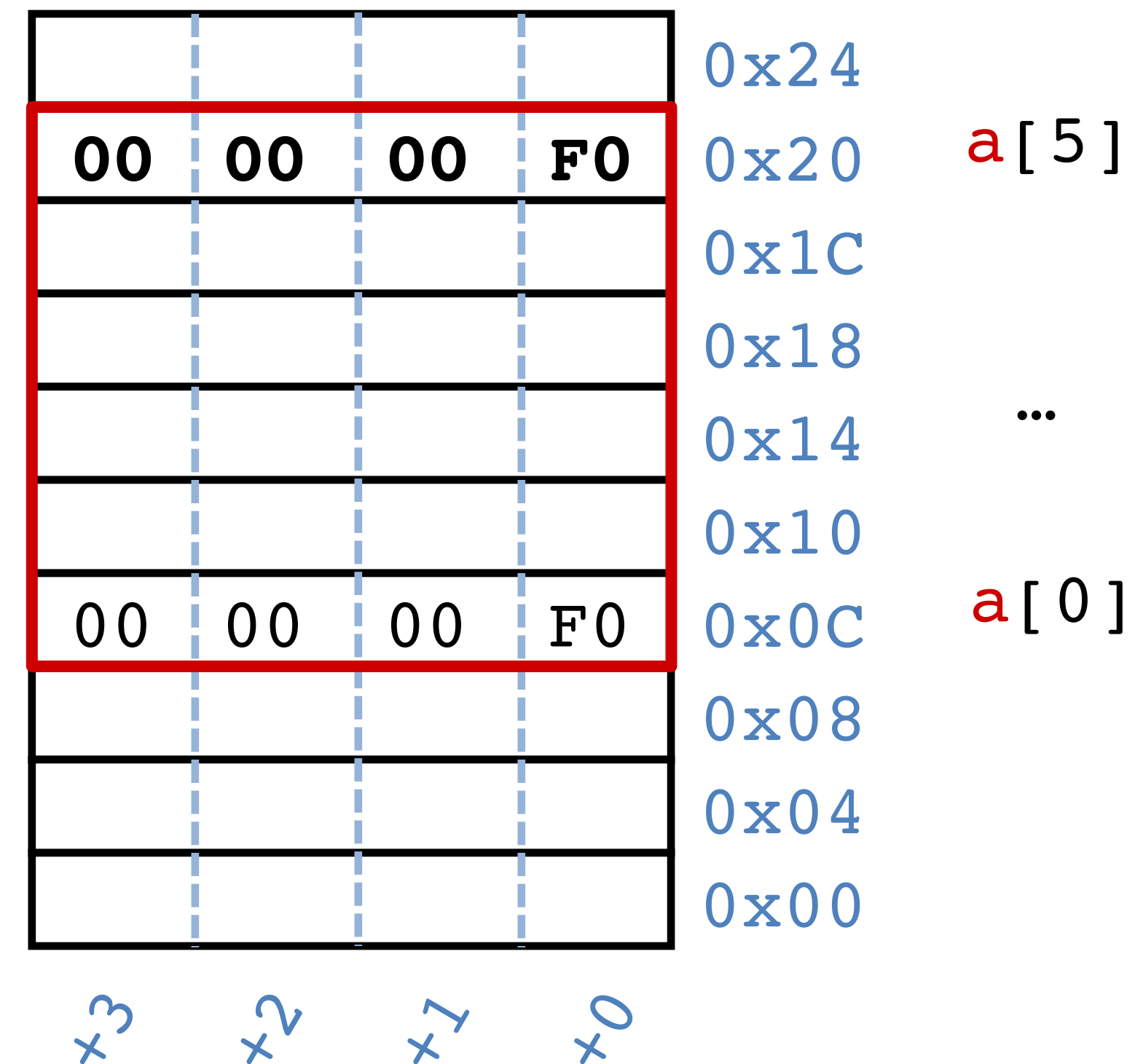
Indexing:

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a[0] = 0xf0;  
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Indexing:

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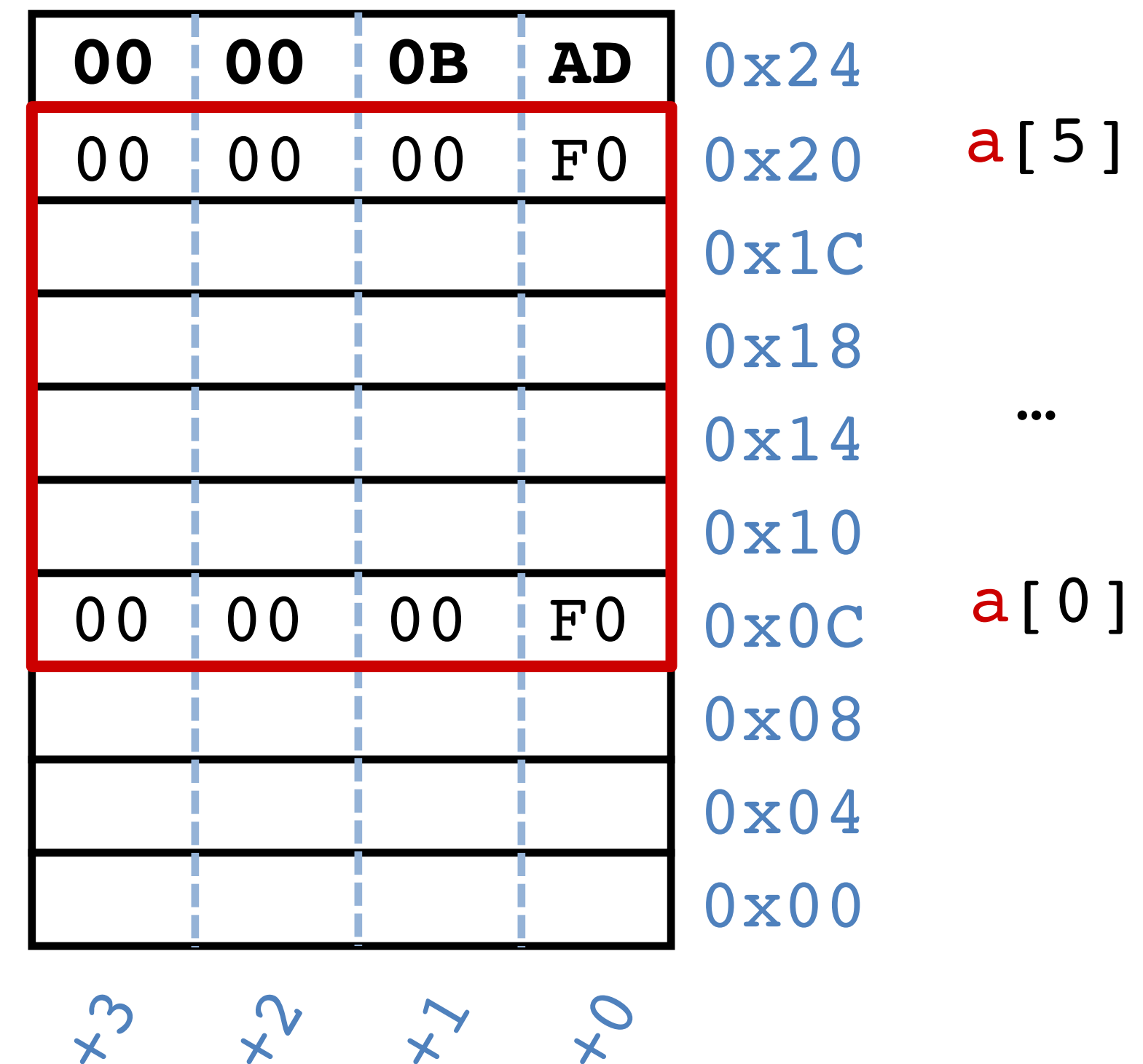
No bounds  
check:

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a[6] = 0xBAD;
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a[6] = 0xBAD;  
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00	00	0B	AD	0x24	
00	00	00	F0	0x20	<b>a[5]</b>
				0x1C	
				0x18	
				0x14	...
				0x10	
00	00	00	F0	0x0C	<b>a[0]</b>
00	00	0B	AD	0x08	
				0x04	
				0x00	
x3	x2	x1	x0		

# C: Arrays

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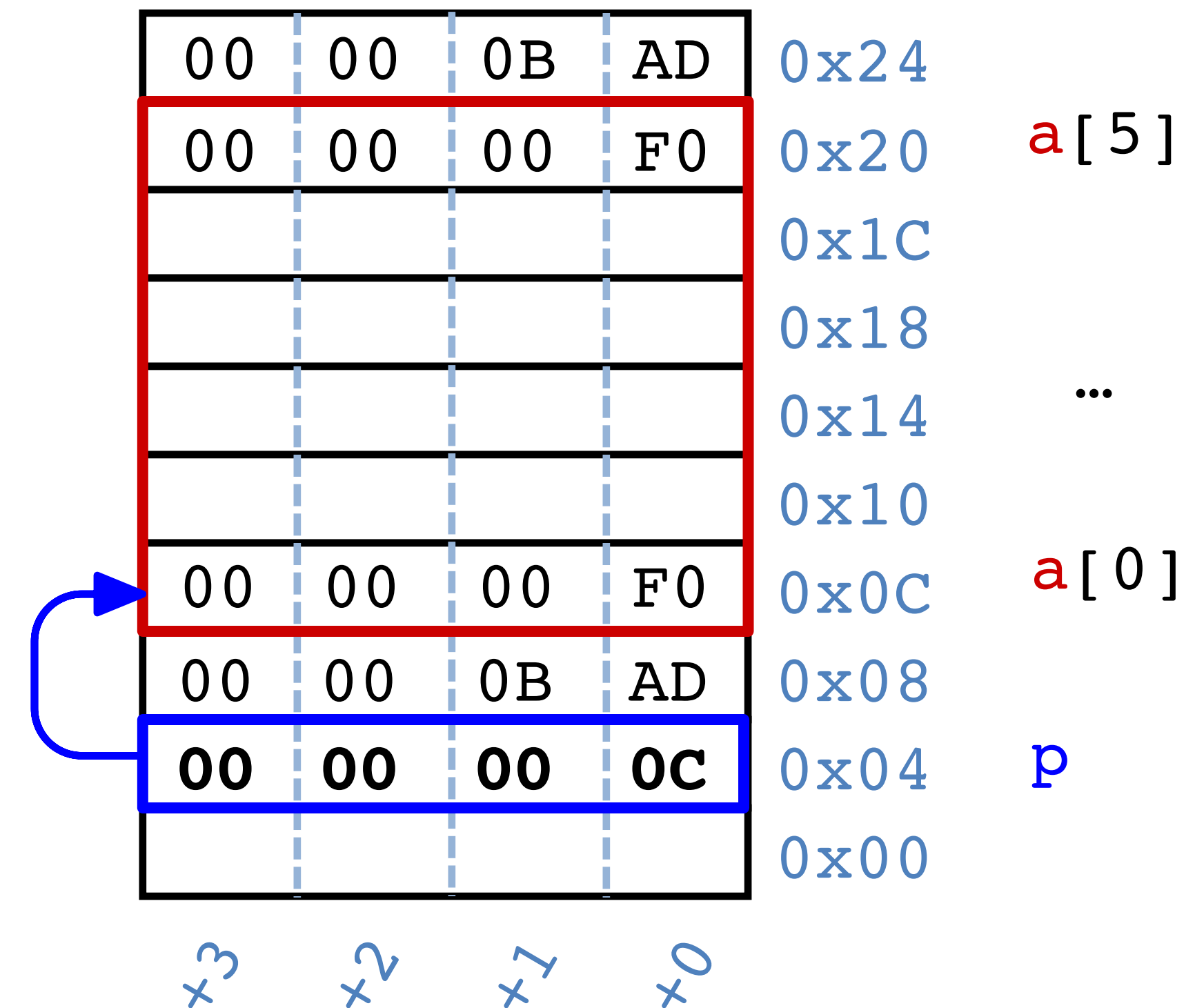
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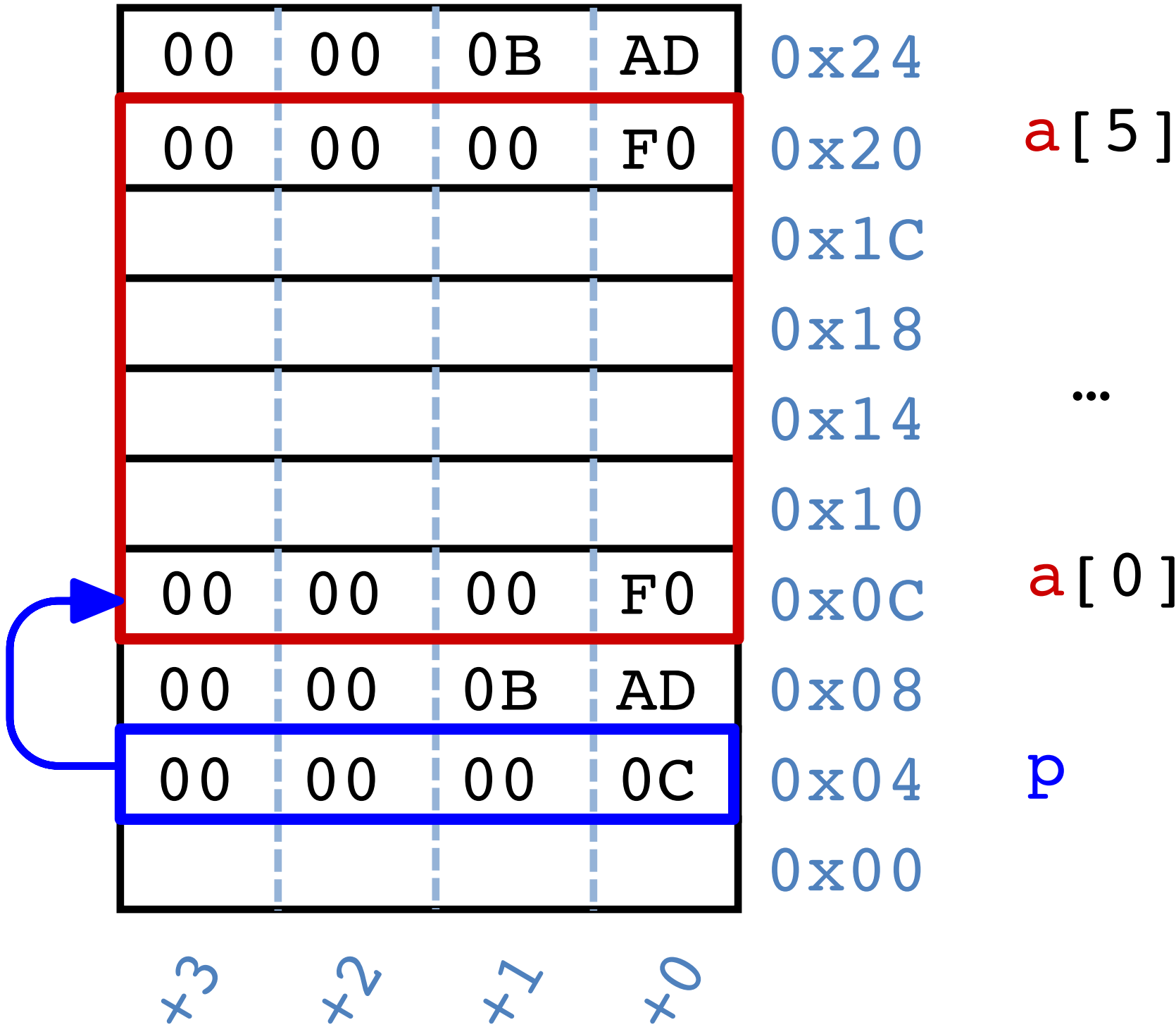
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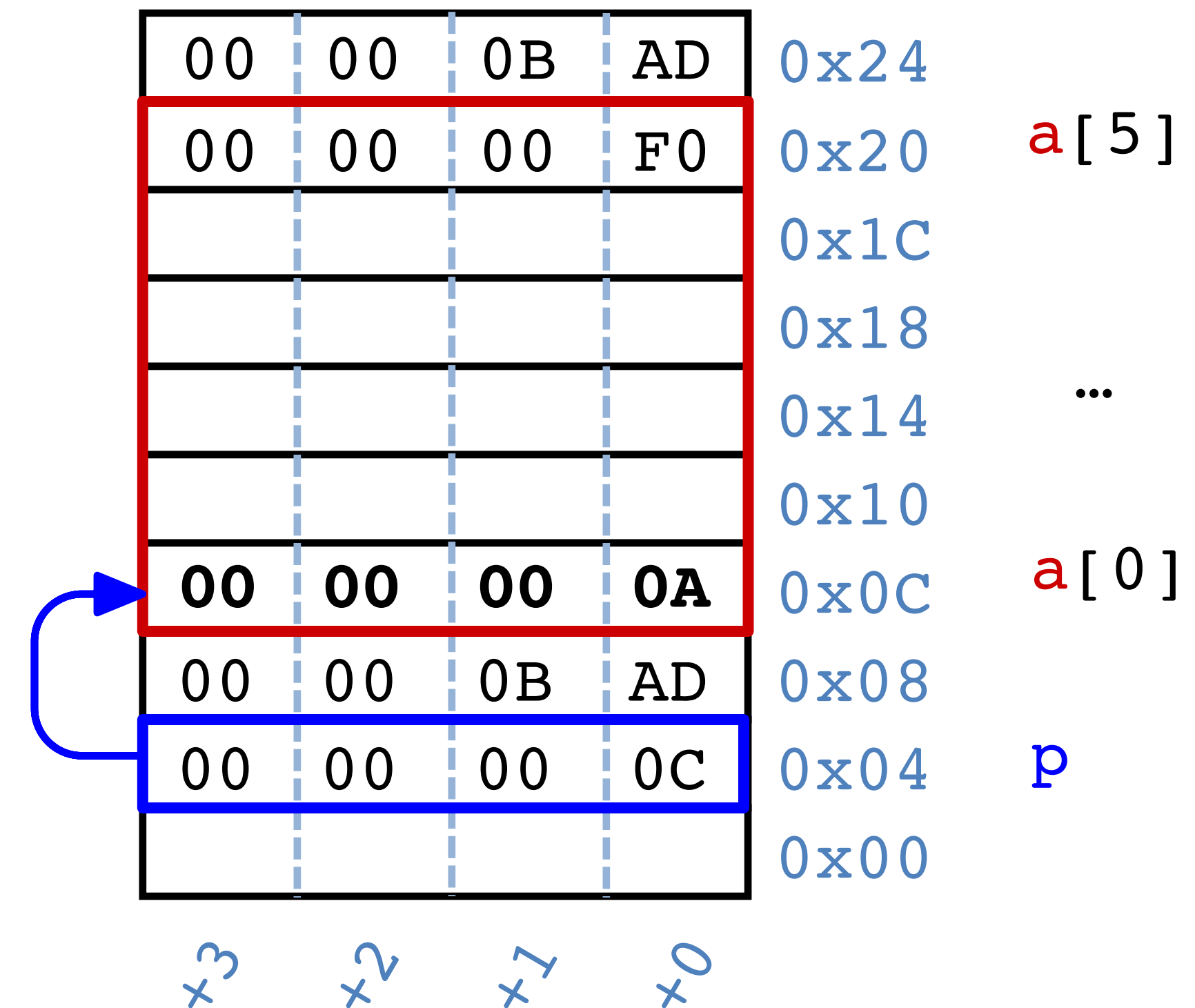
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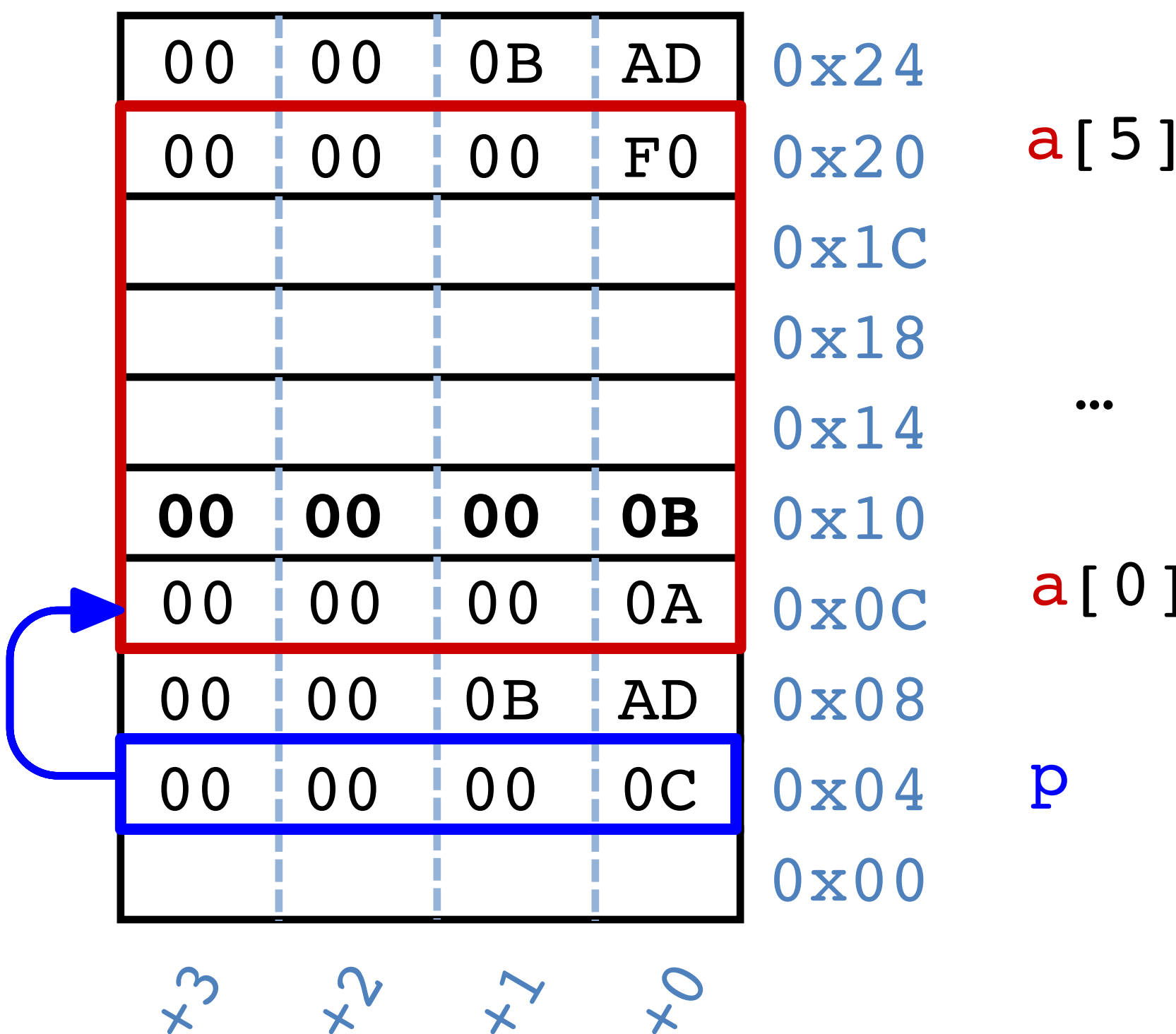
equivalent  $\begin{cases} p[1] = 0xB; \\ *(p + 1) = 0xB; \end{cases}$

*array indexing = address arithmetic*  
Both are scaled by the size of the type.

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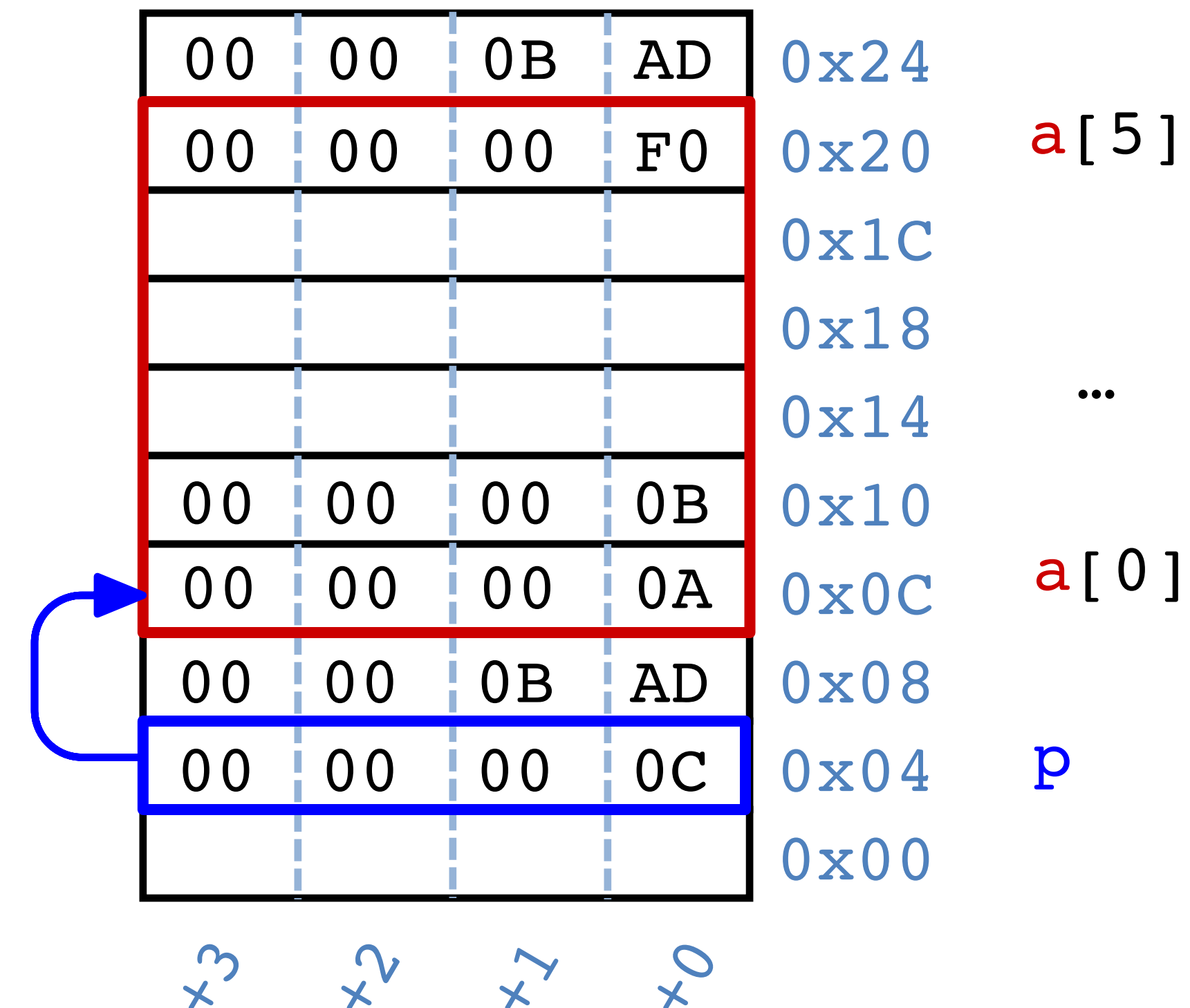
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No bounds check:

`a[6] = 0xBAD;`  
`a[-1] = 0xBAD;`

Pointers:

`int* p;`

equivalent

`p = a;`  
`p = &a[0];`  
`*p = 0xA;`

equivalent

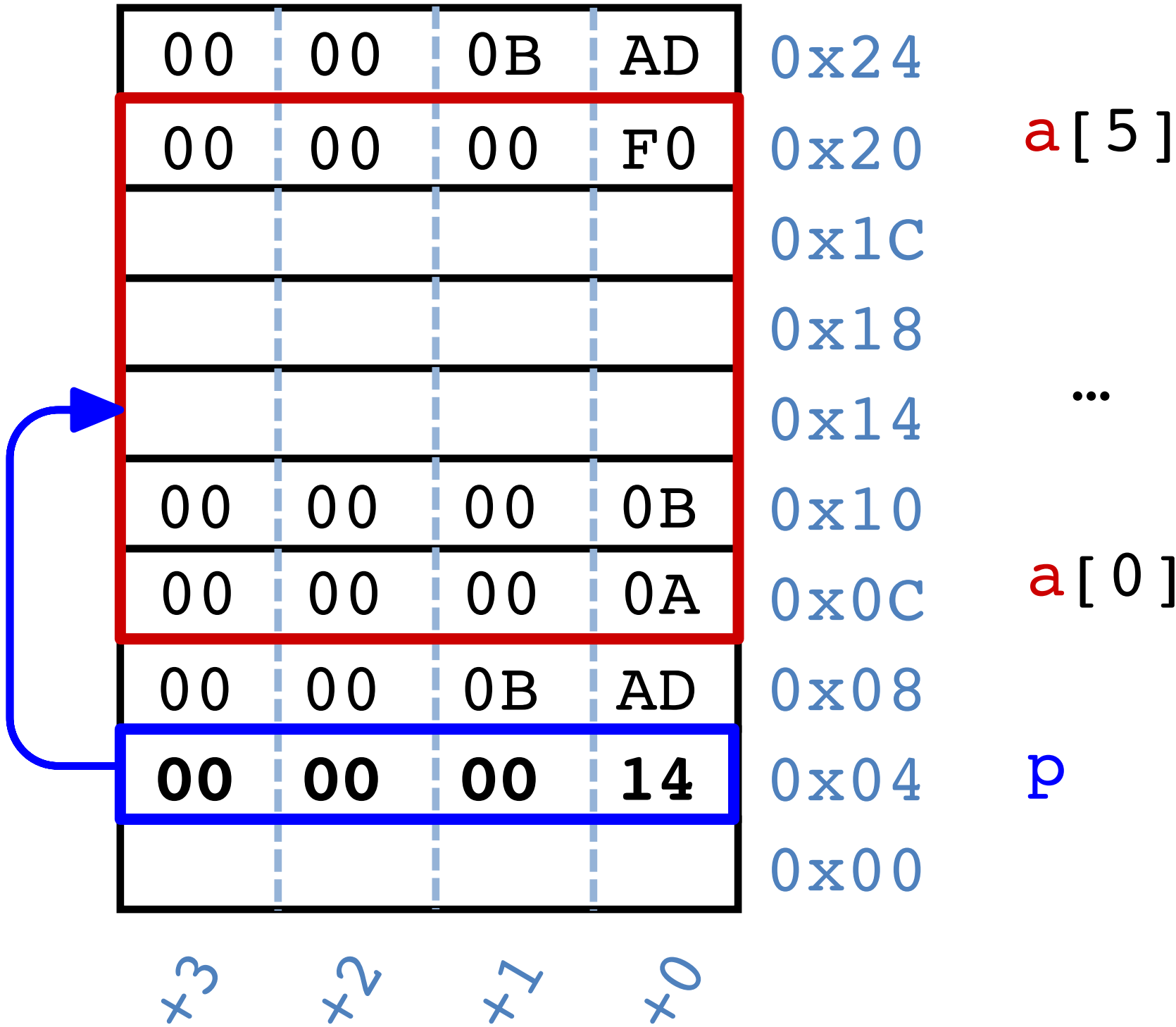
`p[1] = 0xB;`  
`*(p + 1) = 0xB;`  
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array indexing = address arithmetic  
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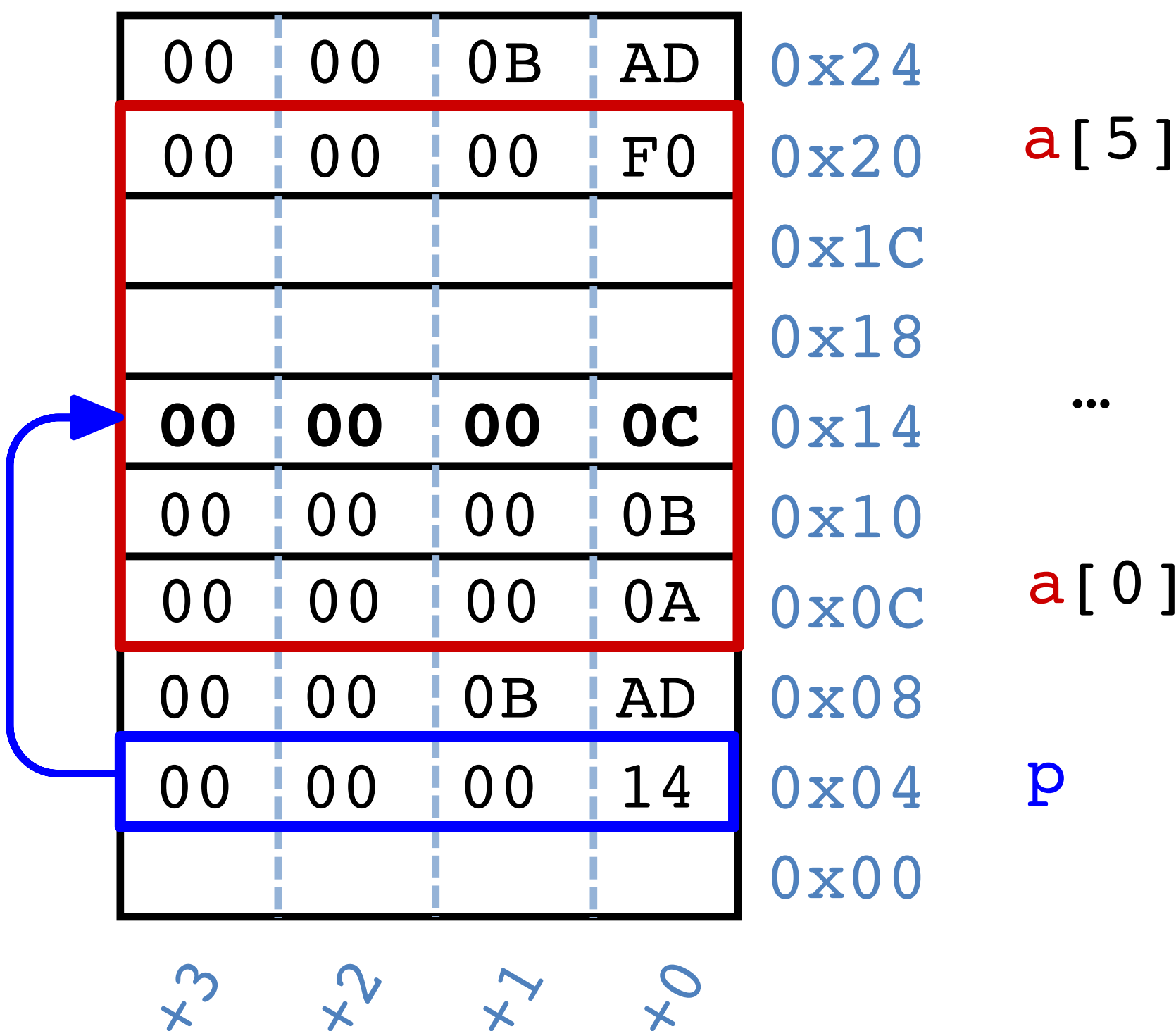
*array indexing = address arithmetic*  
Both are scaled by the size of the type.

`*p = a[1] + 1;`

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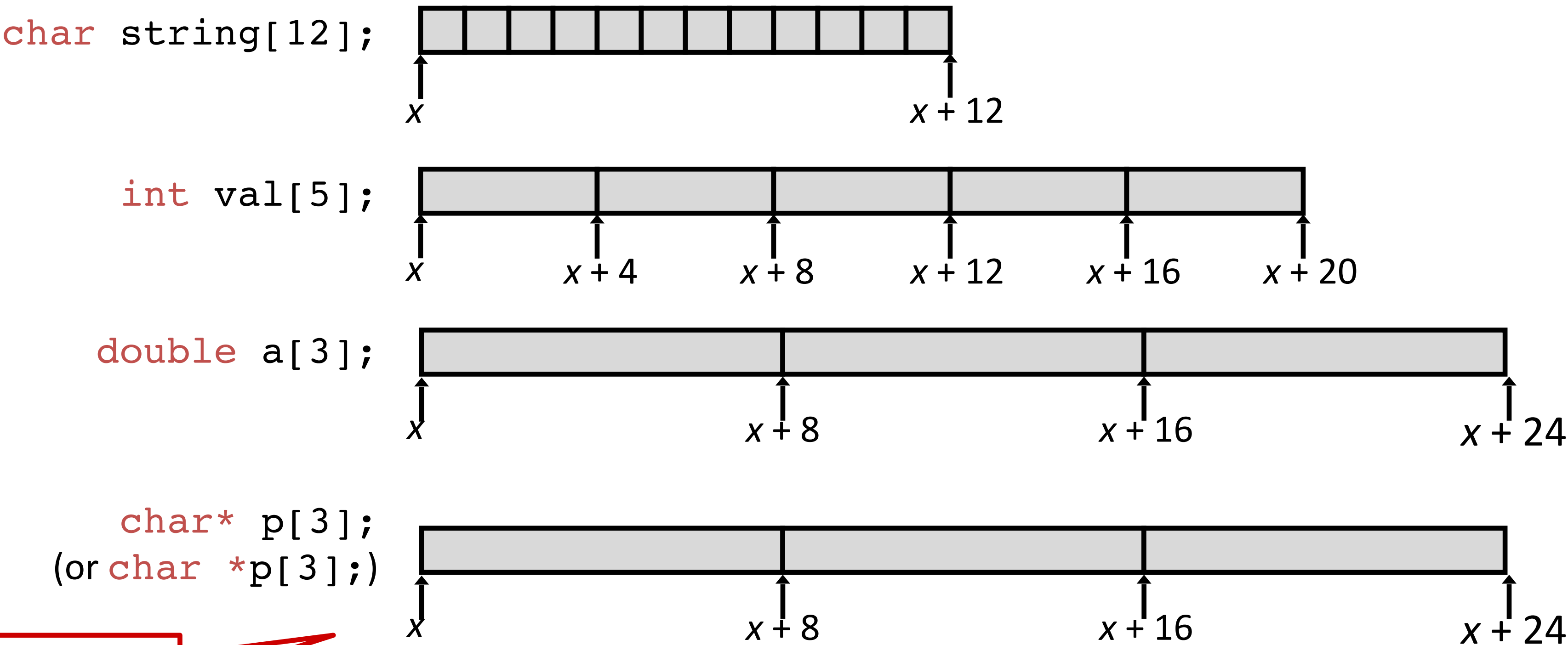


# C: Array allocation

## Basic Principle

$T$   $A[N];$   
Array of length  $N$  with elements of type  $T$  and name  $A$   
Contiguous block of  $N * \text{sizeof}(T)$  bytes of memory

Use *sizeof* to determine proper size in C.



x86-64

size depends on the machine word size

# C: Array access

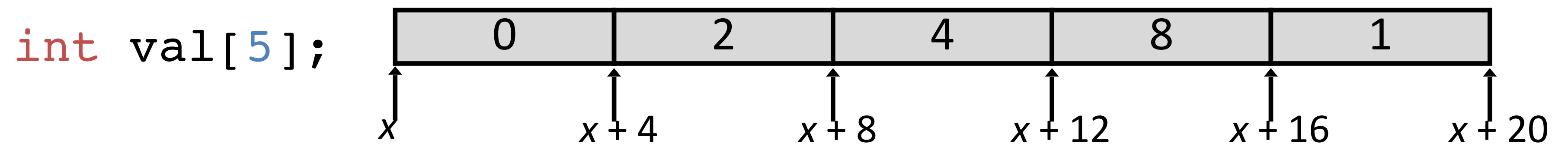
ex

## Basic Principle

$T$   $A[N]$ ;

Array of length  $N$  with elements of type  $T$  and name  $A$

Identifier  $A$  has type  $T^*$



Expression	Type	Value
<code>val[4]</code>	$\text{int}$	1
<code>val</code>	$\text{int}^*$	
<code>val+1</code>	$\text{int}^*$	
<code>&amp;val[2]</code>	$\text{int}^*$	
<code>val[5]</code>	$\text{int}$	
<code>*(val+1)</code>	$\text{int}$	
<code>val + i</code>	$\text{int}^*$	



# Representing strings

A C-style string is represented by an array of bytes (`char`).

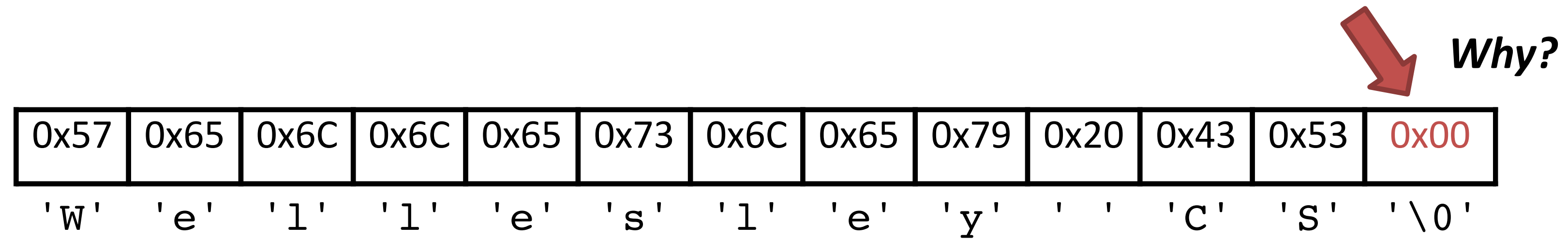
- Elements are one-byte **ASCII codes** for each character.
- ASCII = American Standard Code for Information Interchange

32	space	48	0	64	@	80	P	96	`	112	p
33	!	49	1	65	A	81	Q	97	a	113	q
34	"	50	2	66	B	82	R	98	b	114	r
35	#	51	3	67	C	83	S	99	c	115	s
36	\$	52	4	68	D	84	T	100	d	116	t
37	%	53	5	69	E	85	U	101	e	117	u
38	&	54	6	70	F	86	V	102	f	118	v
39	'	55	7	71	G	87	W	103	g	119	w
40	(	56	8	72	H	88	X	104	h	120	x
41	)	57	9	73	I	89	Y	105	i	121	y
42	*	58	:	74	J	90	Z	106	j	122	z
43	+	59	;	75	K	91	[	107	k	123	{
44	,	60	<	76	L	92	\	108	l	124	
45	-	61	=	77	M	93	]	109	m	125	}
46	.	62	>	78	N	94	^	110	n	126	~
47	/	63	?	79	O	95	_	111	o	127	del

# C: Null-terminated strings

ex

C strings: arrays of ASCII characters ending with *null character*.



Does Endianness matter for strings?

```
int string_length(char str[]) {
```

```
}
```



# C: *\** and *[]*

ex

C programmers often use *\** where you might expect *[]*:

*e.g.*, *char\**:

- pointer to a *char*
- pointer to the first *char* in a string of unknown length

```
int strcmp(char* a, char* b);
```

# C: 0 vs. ' \0 ' vs. NULL

## 0

Name: zero  
Type: `int`  
Size: 4 bytes  
Value: `0x00000000`  
Usage: The integer zero.

## ' \0 '

Name: null character  
Type: `char`  
Size: 1 byte  
Value: `0x00`  
Usage: Terminator for C strings.

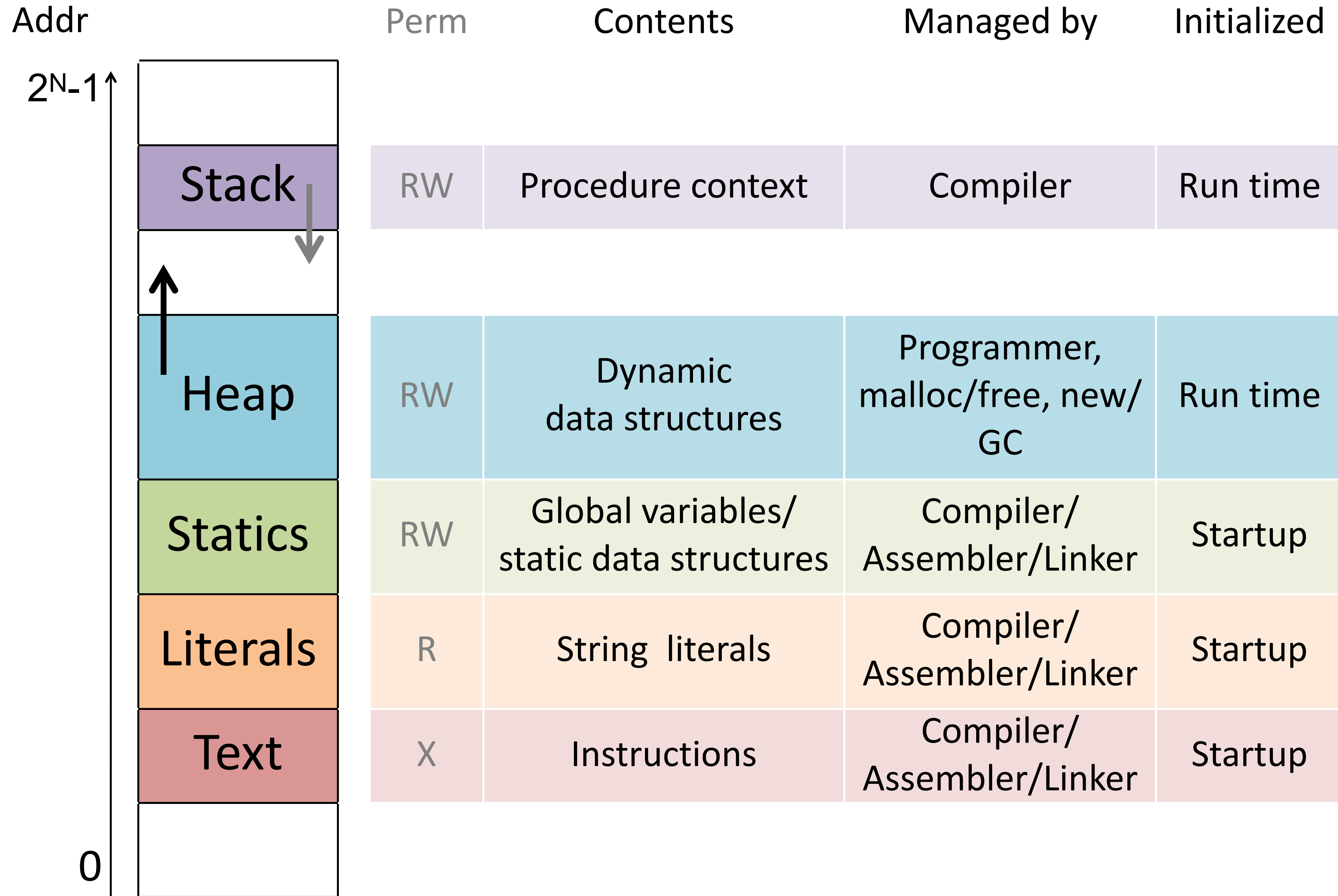
## NULL

Name: null pointer / null reference / null address  
Type: `void*`  
Size: 1 word (= 8 bytes on a 64-bit architecture)  
Value: `0x0000000000000000`  
Usage: The absence of a pointer where one is expected.  
Address 0 is inaccessible, so `*NULL` is invalid; it crashes.

*Is it important/necessary to encode the null character or the null pointer as 0x0?*

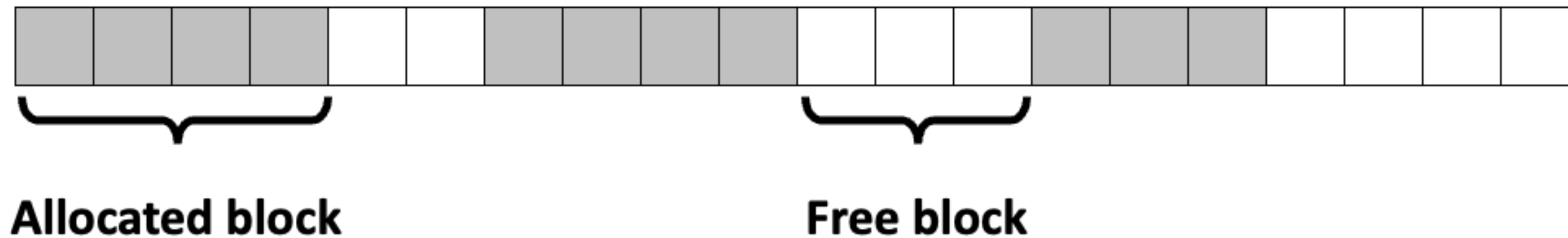
*What happens if a programmer mixes up these "zeroey" values?*

# Memory address-space layout



# C: Dynamic memory allocation in the heap

**Heap:**



**Managed by memory allocator:**

pointer to newly allocated block  
of at least that size

number of contiguous bytes required

```
void* malloc(size_t size);
```

```
void free(void* ptr);
```

pointer to allocated block to free

# C: standard memory allocator

```
#include <stdlib.h>          // include C standard library
```

```
void* malloc(size_t size)
```

Allocates a memory block of at least `size` bytes and returns its address.

If memory error (e.g., allocator has no space left), returns `NULL`.

## Rules:

- Check for error result.

- Cast result to relevant pointer type.

- Use `sizeof(...)` to determine size.

```
void free(void* ptr)
```

Deallocates the block referenced by `ptr`, making its space available for new allocations.

`ptr` **must** be a `malloc` result that has not yet been freed.

## Rules:

- `ptr` must be a `malloc` result that has not yet been freed.

- Do not use `*ptr` after freeing.

# C: Dynamic array allocation

```
#define ZIP_LENGTH 5
int* zip = (int*)malloc(sizeof(int)*ZIP_LENGTH);
if (zip == NULL) {    // if error occurred
    perror("malloc"); // print error message
    exit(0);          // end the program
}

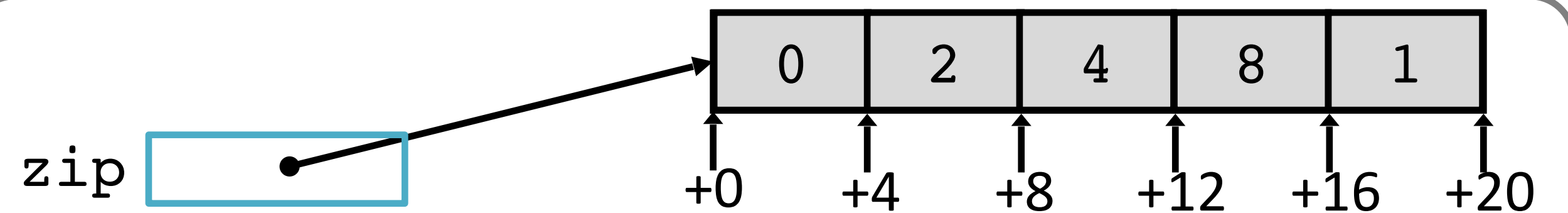
zip[0] = 0;
zip[1] = 2;
zip[2] = 4;
zip[3] = 8;
zip[4] = 1;

printf("zip is");
for (int i = 0; i < ZIP_LENGTH; i++) {
    printf(" %d", zip[i]);
}
printf("\n");

free(zip);
```

zip 0x7fedd2400dc0 0x7fff58bdd938

1	0x7fedd2400dd0
8	0x7fedd2400dcc
4	0x7fedd2400dc8
2	0x7fedd2400dc4
0	0x7fedd2400dc0



# C: Array of pointers to arrays of `ints`

```
int** zips = (int**)malloc(sizeof(int*) * 3);
```

```
zips[0] = (int*)malloc(sizeof(int)*5);
```

```
int* zip0 = zips[0];
```

```
zip0[0] = 0;
```

```
zips[0][1] = 2;
```

```
zips[0][2] = 4;
```

```
zips[0][3] = 8;
```

```
zips[0][4] = 1;
```

```
zips[1] = (int*)malloc(sizeof(int)*5);
```

```
zips[1][0] = 2;
```

```
zips[1][1] = 1;
```

```
zips[1][2] = 0;
```

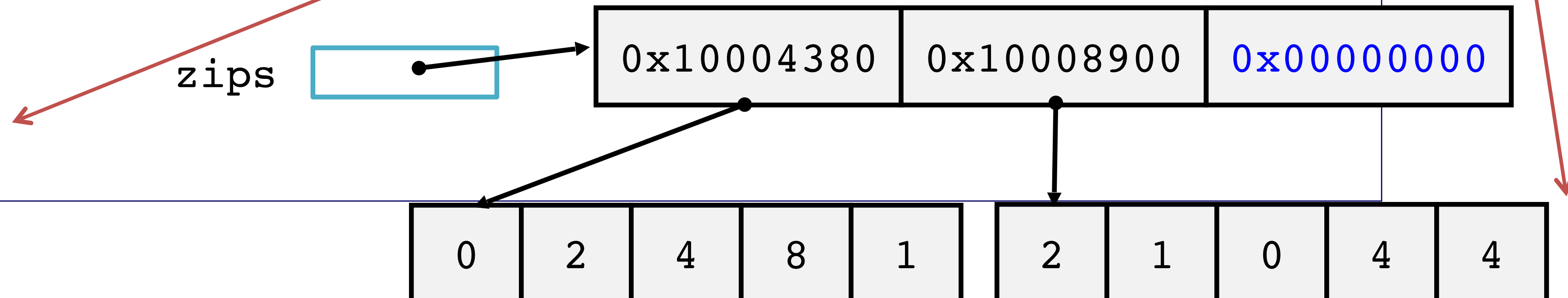
```
zips[1][3] = 4;
```

```
zips[1][4] = 4;
```

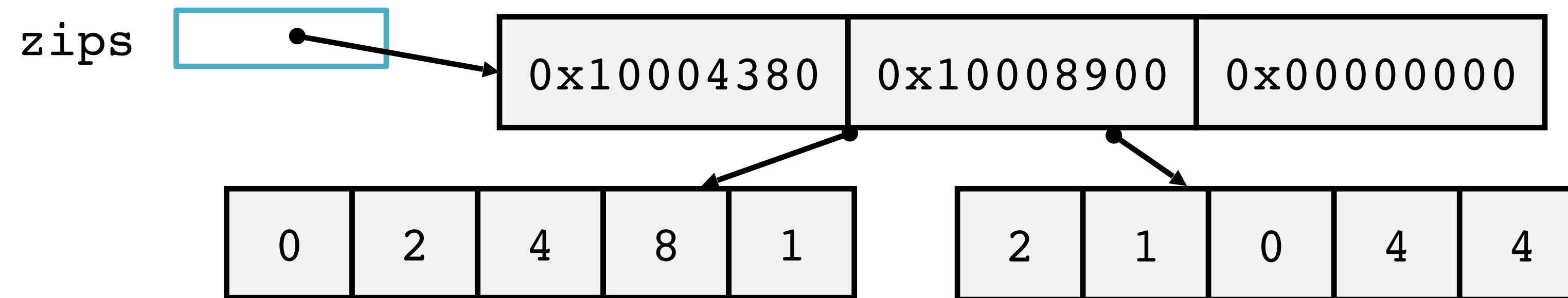
```
zips[2] = NULL;
```

Why  
no NULL?

Why terminate  
with NULL?



# Zip code



```
// return a count of all zips that end with digit endNum  
int zipCount(int* zips[], int endNum) {
```

```
}
```



# scanf reads formatted input

```
int val;
```

Declared, but not initialized.  
Holds anything.

```
...
```

```
scanf( "%d", &val );
```

Read one **int**  
in decimal<sub>10</sub> format  
from input.

Store it in memory  
at this address.

Store in memory at the address  
given by the **address of val**:  
store input @ 0x7F...F38.

val

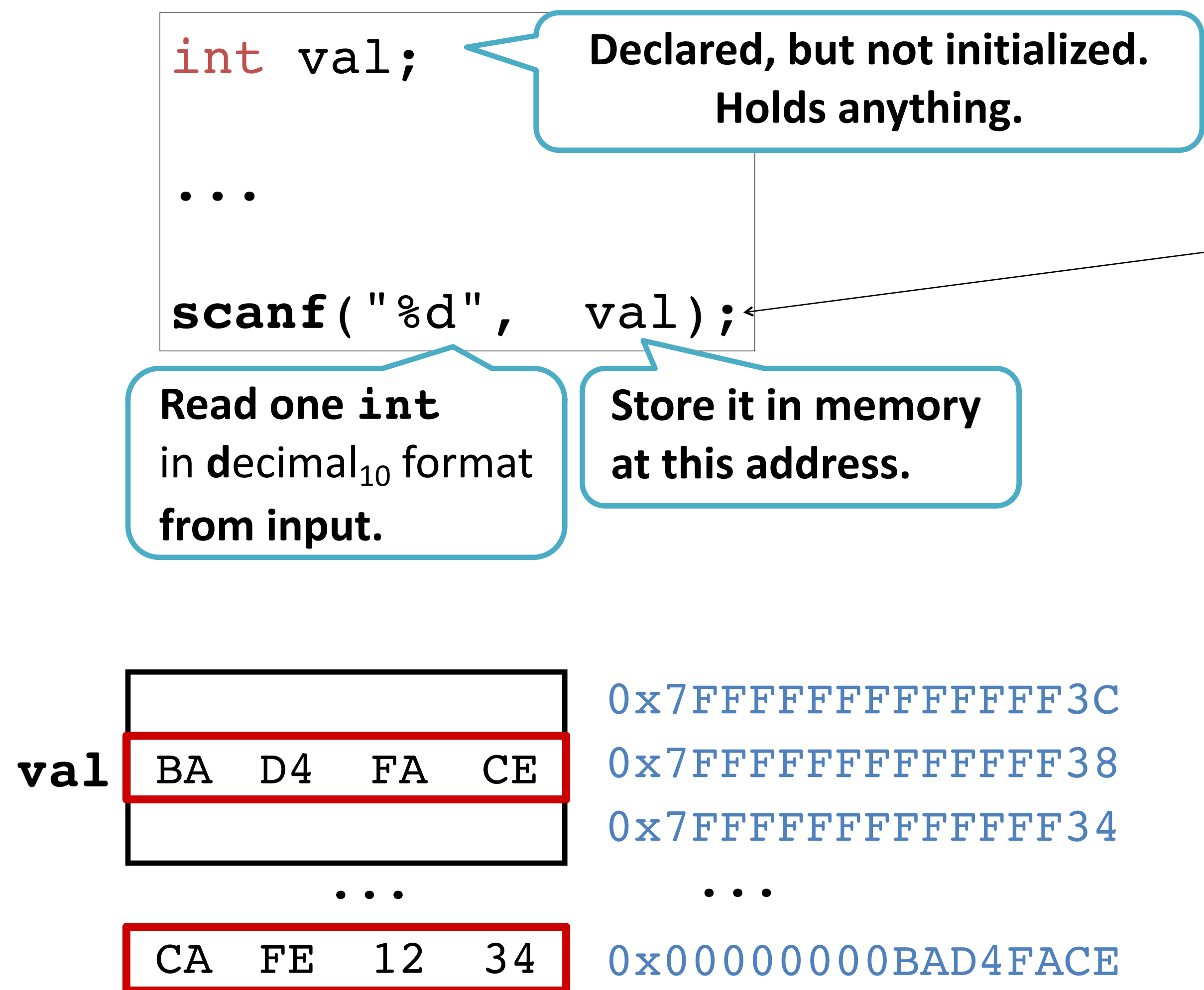
CA	FE	12	34
----	----	----	----

0x7FFFFFFFFFFFFFFF3C

0x7FFFFFFFFFFFFFFF38

0x7FFFFFFFFFFFFFFF34

# C: Classic bug using `scanf`



Store in memory at the address given by the **contents of `val`** (implicitly cast as a pointer): store input @ 0xBAD4FACE.

**Best case:** 🙄 crash immediately with segmentation fault/bus error.

**Bad case:** 🤡 silently corrupt data stored @ 0xBAD4FACE, fail to store input in `val`, and keep going.

**Worst case:** 💻🔥💣🚀 program does literally anything.

# C: Memory error messages

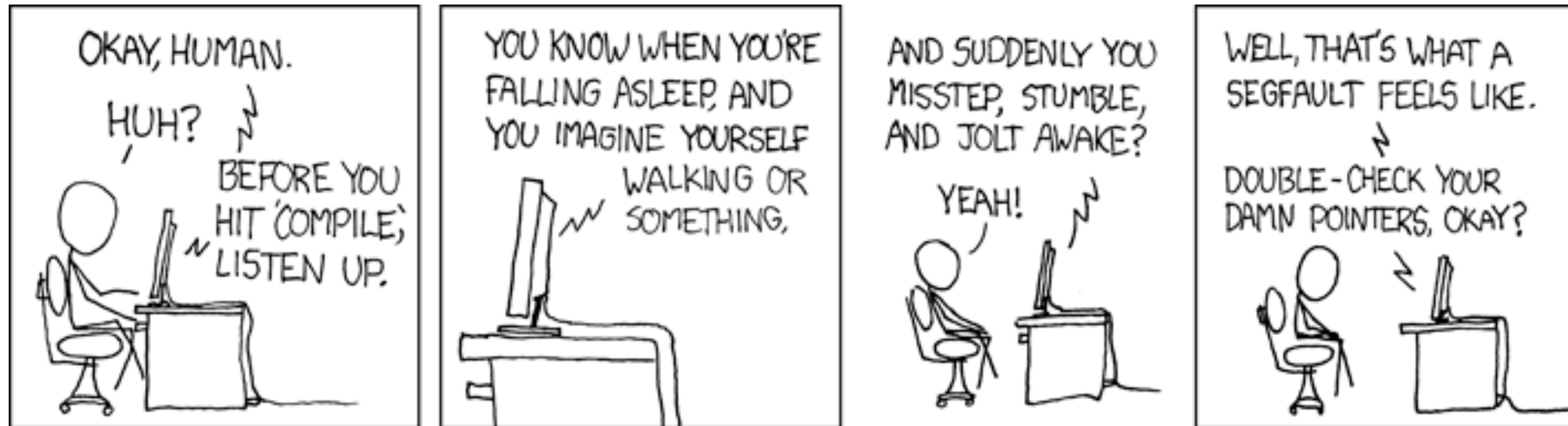
11: **segmentation fault** ("segfault", SIGSEGV)

accessing address outside legal area of memory

10: **bus error** (SIGBUS)

accessing misaligned or other problematic address

More to come on debugging!



<http://xkcd.com/371/>

# C: Why?

## Why learn C?

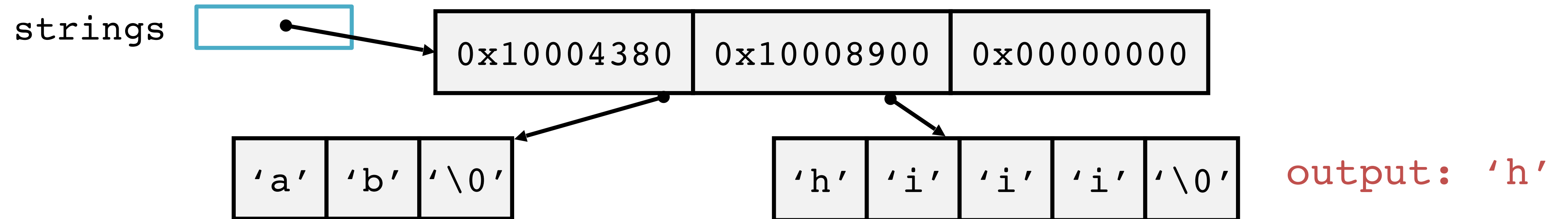
- Think like actual computer (abstraction close to machine level) without dealing with machine code.
- Understand just how much Your Favorite Language provides.
- Understand just how much Your Favorite Language might cost.
- Classic.
- Still (more) widely used (than it should be).
- Pitfalls still fuel devastating reliability and security failures today.

## Why not use C?

- Probably not the right language for your next personal project.
- It "gets out of the programmer's way" ... even when the programmer is unwittingly running toward a cliff.
- Advances in programming language design since the 70's have produced languages that fix C's problems while keeping strengths.

# Group example: longest string starts with

ex



```
// Return the starting character of the longest string in the
// null-terminated strings array.
// You can use: int strlen(char *str)
char longest_string_starts_with(char ** strings) {
```