Local Naming and Scope

These slides borrow heavily from Ben Wood's Fall '15 slides, some of which are in turn based on Dan Grossman's material from the University of Washington.

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Motivation for local bindings

We want local bindings = a way to name things locally in functions and other expressions.

Why?

- For style and convenience
- Avoiding duplicate computations
- A big but natural idea: nested function bindings
- Improving algorithmic efficiency (*not* "just a little faster")

`let` expressions

2 questions:

a new keyword!

- Syntax: ***(let {[id1 e1] ... [idn en]} e_body)***
  - Each ***xi*** is any *variable*, and ***e_body*** and each ***ei*** are any *expressions*
- Evaluation:
  - Evaluate each ***ei*** to ***vi*** in the current dynamic environment.
  - Evaluate ***e_body[v1,…vn/id1,…,idn]***in the current dynamic environment.

  Result of whole **`let`** expression is result of evaluating ***e_body***.

Example

```
> (let {[a (+ 1 2)] [b (* 3 4)]} (list a b))
'(3 12)
```

Pretty printed form

```
> (let {[a (+ 1 2)]
        [b (* 3 4)]}
    (list a b))
'(3 12)
```

Parens vs. Braces vs. Brackets

As matched pairs, they are interchangeable.
Differences can be used to enhance readability.

```
> (let {[a (+ 1 2)] [b (* 3 4)]} (list a b))
'(3 12)

> (let ((a (+ 1 2)) (b (* 3 4))) (list a b))
'(3 12)

> (let [[a (+ 1 2)] [b (* 3 4)]] (list a b))
'(3 12)

> (let [{a (+ 1 2)} (b (* 3 4))] (list a b))
'(3 12)
```


`let` is an expression

A let-expression is ***just an expression***, so we can use it ***anywhere*** an expression can go.

Silly example:

```
(+ (let {[x 1]} x)
   (let {[y 2]
         [z 4]}
     (- z y)))
```


`let` is just syntactic sugar!

```
(let {[id1 e1] … [idn en]} e_body)
```

desugars to

```
((lambda (id1 … idn) e_body) e1 … en)
```

Example:

```
(let {[a (+ 1 2)] [b (* 3 4)]} (list a b))
```

desugars to

```
((lambda (a b) (list a b)) (+ 1 2) (* 3 4))
```


Scope and Lexical Contours

scope = area of program where declared name can be used.

Show scope in Racket via ***lexical contours*** in ***scope diagrams***.

```
(define add-n  (λ ( x ) (+  n  x )) )
(define add-2n (λ ( y ) (add-n (add-n  y ))))
(define n 17)
(define f (λ ( z )
            (let {[ c (add-2n z ) ]
                  [ d (- z 3)     ]}
              (+ z (* c d )))  )   )
```

Declarations vs. References

A **declaration** introduces an identifier (variable) into a scope.

A **reference** is a use of an identifier (variable) within a scope.

We can box declarations, circle references, and draw a line from each reference to its declaration. Dr. Racket does this for us (except it puts ovals around both declarations and references).

An identifier (variable) reference is **unbound** if there is no declaration to which it refers.

Scope and Define Sugar

```
(define (add-n  x ) (+  n  x ) )
(define (add-2n y ) (add-n (add-n  y )) )
(define n 17)
(define (f z )
          (let {[ c (add-2n z ) ]
                [ d (- z 3)     ]}
            (+ z (* c d )))     )     )
```


Shadowing

An inner declaration of a name ***shadows*** uses of outer declarations of the same name.

```
(let {[x 2]}
  (- (let {[x (* x x)]}
       (+ x 3))   Can't refer to
                  outer x here.
     x ))
```


Alpha-renaming

Can consistently rename identifiers as long as it doesn't change the connections between uses and declarations.

```
(define (f w z)
  (* w
     (let {[c (add-2n z)]
           [d (- z 3)]}
       (+ z (* c d))))))
```

OK

```
(define (f c d)
  (* c
     (let {[b (add-2n d)]
           [c (- d 3)]}
       (+ d (* b c))))))
```

Not OK

```
(define (f x y)
  (* x
     (let {[x (add-2n y)]
           [y (- d y)]}
       (+ y (* x y))))))
```

Scope, Free Variables, and Higher-order Functions

In a lexical contour, an identifier is a ***free variable*** if it is not defined by a declaration within that contour.

Scope diagrams are especially helpful for understanding the meaning of free variables in higher order functions.

```
(define (make-sub n )
  (λ ( x ) (- x n)) )

(define (map-scale factor ns )
  (map (λ ( num ) (* factor num )) ns) )
```


Your Turn: Compare the Following

```
(let {[a 3] [b 12]}
  (list a b
        (let {[a (- b a)]
              [b (* a a)]}
          (list a b))))
```

```
(let {[a 3] [b 12]}
  (list a b
        (let {[a (- b a)]}
          (let {[b (* a a)]}
            (list a b)))))
```


New sugar: `let*`

`(let* {} e_body)` desugars to `e_body`

`(let* {[id1 e1] …} e_body)`

desugars to `(let {[id1 e1]}`
`(let* {…} e_body))`

Example:

```
(let {[a 3] [b 12]}
  (list a b
        (let* {[a (- b a)]
               [b (* a a)]}
          (list a b)))))
```


`and` and `or` sugar

`(and)` desugars to **`#t`**
`(and e1)` desugars to **`e1`**
`(and e1 …)` desugars to **`(if e1 (and …) #f)`**

`(or)` desugars to **`#f`**
`(or e1)` desugars to **`e1`**
`(or e1 …)` desugars to

```
(let ((id1 e1))
  (if e1 e1 (or …)))
```

where **`id1`** must be **fresh** – i.e., not used elsewhere in the program.

- Why is let needed in `or` desugaring but not `and`?
- Why must `id1` be fresh?

Avoid repeated recursion

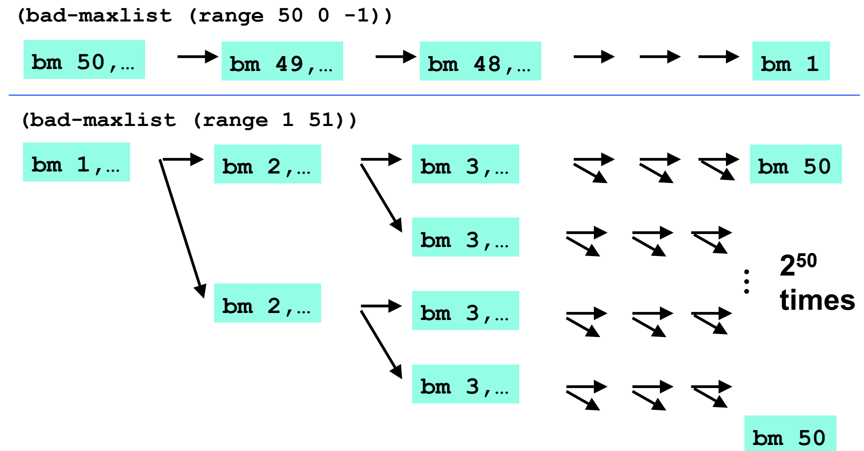
Consider this code and the recursive calls it makes

- Don't worry about calls to **`first`**, **`rest`**, and **`null?`** because they do a small constant amount of work

```
(define (bad-maxlist xs)
   (if (null? xs)
       -inf.0
       (if (> (first xs) (bad-maxlist (rest xs)))
           (first xs)
           (bad-maxlist (rest xs)))))
```


Fast vs. unusable

```
(if (> (first xs)
       (bad-maxlist (rest xs)))
    (first xs)
    (bad-maxlist (rest xs)))
```



Some calculations

Suppose one **`bad-maxlist`** call's **`if`** logic and calls to **`null?`**, **`first?`**, **`rest`** take 10^{-7} seconds total

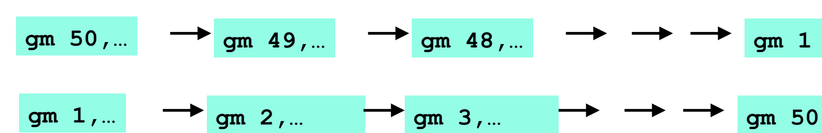
- Then (**`bad-maxlist (list 50 49 … 1))`** takes 50×10^{-7} sec
- And (**`bad-maxlist (list 1 2 … 50))`** takes $(1 + 2 + 2^2 + 2^3 + \ldots + 2^{49}) \times 10^{-7}$ $= (2^{49} - 1) \times 10^{-7} = 1.12 \times 10^8$ sec
  - over 3.5 years
  - **`(bad-maxlist (list 1 2 … 55))`** takes over 1 century
  - Buying a faster computer won't help much ☺

The key is not to do repeated work that might do repeated work that might do…

- Saving recursive results in local bindings is essential…

Efficient maxlist

```
(define (good-maxlist xs)
  (if (null? xs)
      -inf.0
      (let {[rest-max (good-maxlist (rest xs))]}
        (if (> (first xs) rest-max)
            (first xs)
            rest-max))))
```



Transforming good-maxlist

```
(define (good-maxlist xs)
  (if (null? xs)
      -inf.0
      (let {[rest-max (good-maxlist (rest xs))]}
        (if (> (first xs) rest-max)
            (first xs)
            rest-max))))
```

```
(define (good-maxlist xs)
  (if (null? xs)
      -inf.0
      ((λ (fst rest-max) ; name fst too!
         (if (> fst rest-max) fst rest-max))
       (first xs)
       (good-maxlist (rest xs)))))
```

```
(define (good-maxlist xs)
  (if (null? xs)
      -inf.0
      (max (first xs) (good-maxlist (rest xs)))))
```

```
(define (max a b)
  (if (> a b) a b))
```


Local function bindings with `let`

- Silly example:

```
(define (quad x)
    (let ([square (lambda (x) (* x x))])
        (square (square x))))
```

- Private helper functions bound locally = good style.
- But can't use let for local recursion. Why not?

```
(define (up-to-broken x)
  (let {[between (lambda (from to)
                   (if (> from to)
                       null
                       (cons from
                             (between (+ from 1) to))))]}
    (between 1 x)))
```


`letrec` to the rescue!

```
(define (up-to x)
  (letrec {[between (lambda (from to)
                      (if (> from to)
                          null
                          (cons from
                                (between (+ from 1) to))))]}
    (between 1 x)))
```

In ***(let {[id1 e1] ... [idn en]} e_body)***,
e1 … en are in the scope of id1 … idn.

Better

```
(define (up-to-better x)
  (letrec {[up-to-x (lambda (from)
                      (if (> from x)
                          null
                          (cons from
                                (up-to-x (+ from 1)))))]}
    (up-to-x 1)))
```

- Functions can use bindings in the environment where they are defined:
  - Bindings from "outer" environments
    - Such as parameters to the outer function
  - Earlier bindings in the let-expression
- Unnecessary parameters are usually bad style
  - Like **to** in previous example

Mutual Recursion with letrec

```
(define (test-even-odd num)
  (letrec {[even? (λ (x)
                    (if (= x 0)
                        #t
                        (not (odd? (- x 1)))))]
           [odd? (λ (y)
                    (if (= y 0)
                        #f
                        (not (even? (- y 1)))))]}
    (list (even? num) (odd? num))))

> (test-even-odd 17)
'(#t #f)
```


Local definitions are sugar for `letrec`

```
(define (up-to-alt2 x)
  (define (up-to-x from)
    (if (> from x)
        null
        (cons from
              (up-to-x (+ from 1)))))
  (up-to-x 1))

(define (test-even-odd-alt num)
  (define (even? x)
    (if (= x 0) #t (not (odd? (- x 1)))))
  (define (odd? y)
    (if (= y 0) #f (not (even? (- y 1)))))
  (list (even? num) (odd? num)))
```


Nested functions: style

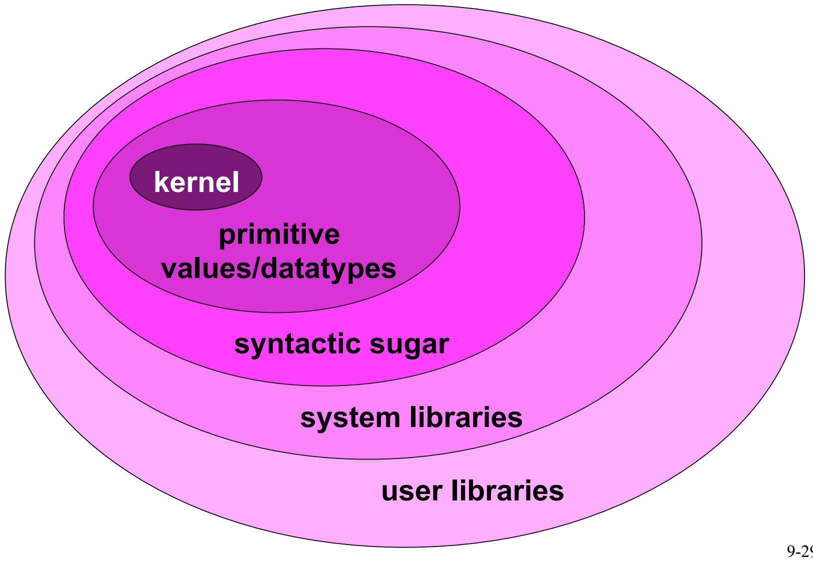
- Good style to define helper functions inside the functions they help if they are:
  - Unlikely to be useful elsewhere
  - Likely to be misused if available elsewhere
  - Likely to be changed or removed later
- A fundamental trade-off in code design: reusing code saves effort and avoids bugs, but makes the reused code harder to change later

Local Scope in other languages

What support is there for local scope in Python?
JavaScript?
Java?

You will explore this in PS5!

Pragmatics: Programming Language Layers



Where We Stand

Kernel	Sugar	Built-in library functions	User-defined library functions