

PROBLEM SET 2
Due Friday, February 25, 2000

Notes:

This is the final version of Problem Set 2. The due date has been extended from Friday, February 18 to Friday February 25. The problem set is now worth 200 points rather than the usual 100 points.

This problem set contains 5 mandatory problems and 3 (optional) extra credit problems

- Problem 1 is the same as posted previously.
- Problem 2 is the same as posted previously except:
 - (1) the description of bits (part **g**) notes that `(bits 0)` is a special case;
 - (2) the previous definition of `fast-expt` had a bug: `(* power ...)` should have been `(* base ...)`. This has been corrected;
 - (3) there are now example applications of `fast-expt` in part **h**;
 - (4) there are **five** new parts (**i**, **j**, **k**, **l**, and **m**);
 - (5) The previous definition of `fold1` has been modified to be consistent with the definition presented in class. In particular, `(op init (car lst))` in the original PS2 has now been replaced by `(op (car lst) init)`;
 - (6) the higher-order list operations `map2` and `map-append` have been removed from Appendix A (if you need them, they are easy to define in terms of the other operators).
- Problems 3 through 5 are new.
- There are 2 extra credit problems worth 60 points of extra credit.

Reading: First-Class Functions handout; *SICP* 1.3, 2.2.3 – 2.2.4.

Problem 1 [20]: Using the Substitution Model to Reason About Higher-Order Functions

Consider the following definitions:

```
(define apply-to-5 (lambda (f) (f 5)))  
  
(define create-subtractor (lambda (n) (lambda (x) (- x n))))
```

Use the substitution model to show the evaluation of the following expressions:

- a. `(apply-to-5 (create-subtractor 2))`
- b. `(apply-to-5 create-subtractor)`
- c. `((apply-to-5 create-subtractor) 2)`
- d. `(create-subtractor apply-to-5)`

Problem 2 [100]: Aggregate Data Paradigm

Implement the following functions in terms of the higher-order list operations in Appendix A. (These can be found on-line in the `cs251` download folder in `ps2/higher-order-list-ops.scm`) You should **not** use recursion in any of your definitions, though you may want to define some auxiliary functions. You will recognize some of these functions from PS1.

a [5] `(append lst1 lst2)`

Return a list containing all the elements of `lst1` followed by the elements of `lst2`.

```
> (append '(1 2 3) '(4 5 6))
(1 2 3 4 5 6)

>(append '((a b) (c d)) '(e (f g) h))
((a b) (c d) e (f g) h)
```

b [5] `(reverse lst)`

Return a list containing the elements of `lst` in reverse order.

```
> (reverse '(a b c d))
(d c b a)

> (reverse '((a b) (c d)))
((c d) (a b))

> (reverse '())
()
```

c [5] `(unzip lst)`

Assume that `lst` is a list of length `len` whose *i*th element is a list of the form `(ai bi)`. Return a list of the form `(lst1 lst2)` where `lst1` and `lst2` are length `len` lists whose *i*th elements are `ai` and `bi`, respectively.

```
> (unzip '((1 a) (2 b) (3 c)))
((1 2 3) (a b c))

> (unzip '((1 a)))
((1) (a))

> (unzip '())
(() ())
```

d [6] (sum-multiples-of-3-or-5 *m n*)

Assume *m* and *n* are integers. Returns the sum of all integers from *m* up to *n* (inclusive) that are multiples of 3 and/or 5.

```
> (sum-multiples-of-3-or-5 0 10)
33 ; 3 + 5 + 6 + 9 + 10

> (sum-multiples-of-3-or-5 -9 12)
22

> (sum-multiples-of-3-or-5 18 18)
18

> (sum-multiples-of-3-or-5 10 0)
0 ; The range "10 up to 0" is empty.
```

e [7] (all-contain-multiple? *n intss*)

Assume that *n* is an integer and *intss* is a list of lists of integers. Returns #t if each list of integers in *intss* contains at least one integer that is a multiple of *n*; returns #f if some list of integers in *intss* does not contain a multiple of *n*. (Note that some Scheme interpreters use the empty list () to stand for #f.)

```
> (all-contain-multiple? 5 '((17 10 12) (25) (3 7 5)))
#t

> (all-contain-multiple? 3 '((17 10 12) (25) (3 7 5)))
#f

> (all-contain-multiple? 3 '())
#t
```

f [8] (cartesian-product *lst1 lst2*)

Returns a list of all duples (*a b*) where *a* ranges over the elements of *lst1* and *b* ranges over the elements of *lst2*. The duples should be sorted first by the *a* entry (relative to the order in *lst1*) and then by the *b* entry (relative to the order in *lst2*).

```
> (cartesian-product '(1 2) '(a b c))
((1 a) (1 b) (1 c) (2 a) (2 b) (2 c))

> (cartesian-product '(2 1) '(c a b))
((2 c) (2 a) (2 b) (1 c) (1 a) (1 b))

> (cartesian-product '(c a b) '(2 1))
((c 2) (c 1) (a 2) (a 1) (b 2) (b 1))

> (cartesian-product '(1) '(a))
((1 a))

> (cartesian-product '() '(a b c))
()
```

g [10] (bits *int*)

Assume *int* is a non-negative integer. Returns a list of bits (i.e, binary digits -- 0s and 1s) in the binary representation of *int*, where the bits are ordered from most significant digit to least significant.

```
> (bits 0)
(0)

> (bits 1)
(1)

> (bits 2)
(1 0)

> (bits 3)
(1 1)

> (bits 10)
(1 0 1 0)

> (bits 20)
(1 0 1 0 0)

> (bits 26)
(1 1 0 1 0)

> (bits 42)
(1 0 1 0 1 0)

> (bits 52)
(1 1 0 1 0 0)
```

Hint: Consider the sequence of numbers obtained by successive integer division by 2 (using Scheme's `quotient` function) until reaching a number less than 1, as shown in the following examples. Do you see a relationship to `bits`?

```
0: () ; Need a special case for 0.
1: (1)
2: (2 1)
3: (3 1)
10: (10 5 2 1)
20: (20 10 5 2 1)
26: (26 13 6 3 1)
42: (42 21 10 5 2 1)
52: (52 26 13 6 3 1)
```

h [10] (fast-expt base power)

The fast exponentiation procedure fast-expt can be defined recursively as follows:

```
;; Assume power is a non-negative integer.
(define fast-expt
  (lambda (base power)
    (if (= power 0)
        1
        (if (even? power)
            (square (fast-expt base (/ power 2)))
            (* base (square (fast-expt base (quotient power 2)))))))

(define even? (lambda (n) (= 0 (remainder n 2))))
(define square (lambda (x) (* x x)))
```

Give a non-recursive definition of fast-expt using the higher-order list operators. *Hint:* use bits from above.

```
> (fast-expt 2 123)
10633823966279326983230456482242756608

> (fast-expt 2 1234)
29581122460809862906004469571610359078633968713537299223955620705065735079623892426105383
72483780501864436477590709559931208208993303817609370272124828409449413621106654437751834
95726811929203861182015218323892077355983393191208928867652655993602487903113708549402668
62452110061179427034023276609931709804888749380902312739825386061877261903500988327294112
9544640111837184

> (fast-expt 2 12345)
16417101068825821635602074166390650141012723553073588127211610308792509417139014428015903
45364394577348704191271404016671955103310856571853327210892364011930444934571162997688443
44303479235489462436380672117015123283299131391904179287678259173308536738761981139958654
88085223490844833881728901416677416986925133937982859974849291877543786473903221777805133
38829900741162462812693649337248923421345047024910400166375574298108937807651974185894775
84716543480995722533317862352141459217781316266211186486157019262080414077670264642736018
42699811352344573268085614432987697227330070339258499772920719797108394570034549409240014
71869973070120694540684895890356769794481698480608369249458241977064933061082585119360303
41393221586423523264452449403781993352421885094664052270795527632721896121424813173522474
67439588615509220340403673074847478171071574544613546809813983182408325964791917527350368
15611726846242833844385047765030004322416045504543741163208222271919113221234840850639263
50606342197146407841178028071147192533942517270553513988142925976090769695456221159699052
58353301133165207934709309817308697548353927446402335745648446548292747956943732036859222
27602781703060767334388010983707976757112746710549707114421589305616843431357741187415945
06702833147396758825015850042983343690345185995956235143825771620543546030664562647854656
43130264457411987382021559571861862448523242200657555000706888373424145468636885673449626
53859088094039724946851377411228668967196780539372858184097516703201405018430392240407358
70096889596273419106389103662095318937990625980136711988237421962315266686856089505981438
44085063806758932114175949901702383959685845554819200014008514229416698706349902479268133
48431597909363213519198597586695692005415076120997809097051989021760262198722017154220960
903436862729843514415945695067780410626632667993427938563138015409598158457885847590843248
82824856158645027117277724097179565608200184811581526093052166316748017388606401911857277
8281516735157795558881677870644325585954108439874464978816662884232331700604130259246299
50477303342180149398926073618582715358742250388958231281694757980523791263699450732952325
72766420994778606398256177532763850451691857010131939169841238860760374248441474826838966
91291180268789697357822868411168426564105746476075244189007203280453779933862798087689903
762894247573510523693939771378719981191688984930379387566356215576231384044459266598837784
22932579983878202606048149686556175703183900225709180287694924839274417566911224208843988
32483363105970012573859807769615293511988777471935310549568818083321779467514040382287185
67911769630971553915410012677600002457982207465176670752102117002773980548089696530972476
43969459988128181297321726585388472790653547974585408533885110514458548199415620649743674
58999448777325314125412790143003245948906239411455098569409828637698344300481205629667979
07114102689879364945689860493474954538422367719507882513166051007352994068319251450666676
64836820056432938299875887576041425965400497726130998826731980635485605178455399093661063
47333759841590287223786149844502553863155856319454033500021429104931902548256107074005899
76364985748467955131077971641882672895854571236368282811336220769174784720113331269084746
52420412426347505411284163093358616619503611569646968607560048042056355756761683563325262
23271728110021463927544450511821698052846302597035426339551261795201130596299142298336885
35925729676778028406897316106101038469119090984567152591962365415039646394591503830797626
33924698605707775861141366491416874537526678629814117149657394161438774412584368567706361
97829187598231060210540377578577615874722408350405804473605440290649304125699431697292381
02162312218687930203068055400275795180972382856696655279408212344832
```

i [7] (repeated *fun* *n*)

Return the *n* -fold composition of the function *fun*.

```
> ((repeated (lambda (x) (+ x 1)) 5) 0)
5
```

```
> ((repeated (lambda (x) (* 2 x)) 3) 1)
8
```

j [7] (inner-product *nums1* *nums2*)

Assume that *nums1* is the list of numbers (*a1 a2 ... an*) and *nums2* is the list of numbers (*b1 b2 ... bn*). (Note that both lists are assumed to have length *n*.) Return the sum of the products of the corresponding elements of the two lists – i.e., the value $(a1*b1) + (a2*b2) + \dots + (an*bn)$.

```
> (inner-product '(1 2 3) '(4 5 6))
32 ; 4 + 10 + 18
```

```
> (inner-product '() '())
0
```

k [10] (splits *lst*)

Returns a list of all duples (*a b*) such that appending *a* and *b* gives *lst*. The order of the duples does not matter.

```
> (splits '(1 2 3))
(((1) (2 3)) ((1) (2 3)) ((1 2) (3)) ((1 2 3) ()))
```

```
> (splits '(2 3))
(((2) (3)) ((2) (3)) ((2 3) ()))
```

```
> (splits '(2))
(((2) (2)) ((2) ()))
```

```
> (splits '())
(((2) (2)) ((2) ()))
```

l [10] (insert-all elt lst)

Assume that *lst* is a list of elements not containing *elt*. Returns a list of all the distinct ways that *elt* can be inserted into *lst*, maintaining the relative order of the elements of *lst*. The order of elements in the result does not matter.

```
> (insert-all 1 '(2 3 4))
((1 2 3 4) (2 1 3 4) (2 3 1 4) (2 3 4 1))

> (insert-all 1 '(3 4))
((1 3 4) (3 1 4) (3 4 1))

> (insert-all 1 '(4))
((1 4) (4 1))

> (insert-all 1 '())
((1))
```

Hint: use splits from above.

m [10] (permutations lst)

Assume that *lst* is a list of distinct elements (i.e., no duplicates). Returns a list of all the permutations of the elements of *lst*. The order of the permutations does not matter.

```
> (permutations '())
(())

> (permutations '(1))
((1))

> (permutations '(1 2))
((1 2) (2 1)) ; Order doesn't matter

> (permutations '(1 2 3))
((1 2 3) (1 3 2) (2 1 3) (2 3 1) (3 1 2) (3 2 1)) ; Order doesn't matter
```

Hint: use insert-all from above.

Problem 3 [20]: Lexical Ordering

Consider the following Scheme sorting function (which can be found in the CS251 download folder in `ps2/lexord.scm`):

```
(define insertion-sort
  (lambda (less-than? elts)
    (if (null? elts)
        '()
        (insert less-than?
                 (car elts)
                 (insertion-sort less-than? (cdr elts))))))

(define insert
  (lambda (less-than? elt lst)
    (if (null? lst)
        (list elt)
        (if (less-than? elt (car lst))
            (cons elt lst)
            (cons (car lst) (insert less-than? elt (cdr lst)))))))
```

The `insertion-sort` procedure takes a binary `less-than?` predicate and a list of elements and returns a list of all the elements in sorted order from smallest to largest according to the `less-than?` predicate. For example:

```
> (insertion-sort < '(6 1 8 10 3 5))
(1 3 5 6 8 10)

> (insertion-sort > '(6 1 8 10 3 5))
(10 8 6 5 3 1)

> (insertion-sort
   ;; Sort elements by their distance from 5
   (lambda (a b)
     (< (abs (- a 5))
        (abs (- b 5)))))
   '(6 1 8 10 3 5))
(5 6 3 8 1 10)
```

For any binary predicate P of two elements of type T , it is possible to "lift" P to an ordering P' that compares two lists of elements of type T . The lifted ordering P' is called a **lexical ordering predicate**. Informally, P' compares two lists $L1$ and $L2$ elementwise from left to right using P . It returns true if it discovers an element of $L1$ less than the corresponding element of $L2$ or if $L1$ is a proper prefix of $L2$; it returns false if it discovers an element of $L1$ greater than the corresponding element of $L2$ or if $L2$ is a prefix of $L1$.

More formally, suppose list $L1$ has the m elements a_1, a_2, \dots, a_m and list $L2$ has the n elements b_1, b_2, \dots, b_n . Then $(P' L1 L2)$ is true if and only if one of the following two conditions holds:

1. There is a k in the range $[1..m]$ such that a_i and b_i are equal for all i in the range $[1..k-1]$ and $(P a_k b_k)$ is true.
2. $m < n$ and a_i and b_i are equal for all i in the range $[1..m]$

Lexical ordering is the standard way that alphabetic comparisons on characters is extended to compare two strings. It is often called **dictionary ordering**, since it determines how words are arranged in a dictionary.

In this problem, you are to write a Scheme function `lexord` that returns the lexical ordering predicate P' when given the predicate P . To test two elements for equality, you should **not** use `=`, `eqv?`, `eq?`, or `equal?`. Rather, observe that ai and bi are equal if both $(P ai bi)$ and $(P bi ai)$ are false.

Here are some examples of `lexord` in action:

```
> (insertion-sort (lexord <))
'((2 1) (1 3) (1) () (1 2 3) (2 2) (1 1) (3 2 1) (1 2) (2)))
((() (1) (1 1) (1 2) (1 2 3) (1 3) (2) (2 1) (2 2) (3 2 1)))

> (insertion-sort (lexord >))
'((2 1) (1 3) (1) () (1 2 3) (2 2) (1 1) (3 2 1) (1 2) (2)))
((() (3 2 1) (2) (2 2) (2 1) (1) (1 3) (1 2) (1 2 3) (1 1)))

> (insertion-sort (lexord (lexord <)))
'(((1 2) (2 1)) ((2 1) (2 2)) ((1 1) (2 2))
  ((1 2) (1 1)) ((2 1) (1 2))))
(((1 1) (2 2)) ((1 2) (1 1)) ((1 2) (2 1)) ((2 1) (1 2)) ((2 1) (2 2)))
```

You can use the `test-lexord` procedure in `lexord.scm` to run the above three test cases.

Problem 4 [40]: Functional Representation of Sets

In CS230, you learned the the extremely important notion of an abstract data type (ADT). In short, a data type can be defined by an interface of routines that manipulate elements of that type, independent of the details of how those routines are implemented.

ADTs are realizable in almost any programming language. For example, here is the Scheme interface to an ADT for a set of numbers:

```
(set-empty)
Return an empty set.

(set-singleton x)
Return a set whose single element is x.

(list->set lst)
Return a set whose elements are the elements of the list lst.

(set-member? x s)
Return #t if x is in set s and #f otherwise.

(set-union s1 s2)
Return a set whose elements are those that are in either s1 or s2.

(set-intersection s1 s2)
Return a set whose elements are those that are in both s1 and s2.

(set-difference s1 s2)
Return a set whose elements are those in s1 that are not in s2.
```

As in Java, we can implement this set ADT in Scheme in terms of familiar data structures like arrays, lists, or trees. However, unlike Java, Scheme also allows abstract data types to be implemented as functions. Intuitively, functions are just another kind of data structure. In fact, we shall see that functions are often more flexible data structures than conventional arrays, lists, and trees.

As a concrete example of this approach, we will explore how to implement sets as functions. In particular, we will represent a set as the membership predicate that determines whether a given element is in the set. For instance, the set $\{2, 3, 5\}$ can be represented as the function

```
(lambda (x)
  (or (= x 2) (= x 3) (= x 5)))
```

This function returns #t for the numbers 2, 3, and 5, but returns #f for all other numbers. The empty set can be represented as the function that returns #f for all numbers:

```
(lambda (x) #f)
```

This functional representation has numerous advantages over the array/list/tree versions. In particular, it is easy to specify sets that have infinite numbers of elements! For example, the set of all even numbers can be represented by the function

```
(lambda (x) (= (remainder x 2) 0)).
```

This predicate is true of even integers, but is false for all other numbers. The set of all real numbers between 5 and 7 (inclusive) can be represented by the function:

```
(lambda (x) (and (>= x 5) (<= x 7)))
```

The set of all integers can be represented by the standard Scheme predicate `integer?`, while the set of all numbers can be represented by

```
(lambda (x) #t)
```

(This assumes that the predicates are only being applied to numbers. If we extended the notion of set to include any Scheme value, then the set of all numbers would be represented as the predicate `number?`)

a. Representing sets as membership predicates, implement the seven functions in the set ADT presented above. You should test out your implementation on suitable test cases, but only need to turn in your seven definitions.

b. Below are some other routines we could add to the interface to the set ADT. For each such routine, indicate whether or not it is possible to implement the routine (1) when sets are represented as lists (2) when sets are represented as membership predicates. Justify your answers.

i. `(set-empty? set)`

Return `#t` if the `set` is empty, and false otherwise.

ii. `(predicate->set pred)`

Given a membership predicate `pred`, return a set of the elements for which `pred` is true.

iii. `(set->list set)`

Return a list of all the elements in `set`.

iv. `(set-complement set)`

Return the set of all numbers not in `set`.

v. `(subset? set1 set2)`

Return `#t` if all of the elements of `set1` are also elements of `set2`, and `#f` otherwise.

Problem 5 [20] Church Numerals

The First-Class Functions handout discusses how n -fold composition functions (so-called Church numerals) can be viewed as the basis of a system for arithmetic. Write Scheme definitions for the functions `plus`, `times`, and `raise` that are described near the end of the First-Class Function handout.

Notes:

- The file `ps2/church.scm` in the CS251 download folder contains the code from the function composition section of the First-Class Functions handout, including `int->church` and `church->int`.
- For ideas on how to implement these three functions, carefully study the examples involving `twice` and `thrice` in the function composition section of the First-Class Functions handout.
- Each of your function definitions should be *extremely* short. In fact, it's possible to implement each definition as a "one-liner". (But if you obey Scheme pretty-printing conventions, your definitions will be several lines long.)

Extra Credit Problems

These problems are optional. You should only attempt them after completing the rest of the problems. (Note that extra credit problems need not be turned in by the due date; they can be handed in any time during the semester. However, experience shows that students rarely turn them in after the problem set is due.)

EC1 [20]: Partitioning

Given an equality predicate $eqpred$ (a so-called **equivalence relation**) and a list of elements, it is possible to **partition** the list into sublists (so-called **equivalence classes**) such that (1) every pair of elements from a given equivalence class are equal according to $eqpred$; and (2) no two elements from distinct equivalence classes are equal according to $eqpred$. Define a Scheme function `(partition $eqpred$ lst)` that partitions lst into equivalence classes according to $eqpred$. The order of elements within an equivalence class does not matter, nor does the order of equivalence classes within a partition. For example:

```
> (partition (lambda (a b)
              (= (remainder a 3) (remainder b 3)))
  '(17 42 6 11 16 57 51 1 23 47))
((17 11 23 47) (16 1) (42 6 57 51))

> (partition (lambda (a b)
              (= (quotient a 10) (quotient b 10)))
  '(17 42 6 11 16 57 51 1 23 47))
((17 57 47) (23) (11 51 1) (6 16) (42) )

> (partition (lambda (lst1 lst2)
              (= (length lst1) (length lst2)))
  '((1 2) () (2 2) (1 2 3) (1 3) (6) (2 1) (1 2 1) (3 1 2) (5 1)))
(((5 1) (2 1) (1 3) (2 2) (1 2)) ()) ((3 1 2) (1 2 1) (1 2 3)) ((6)) )

> (partition (lambda (lst1 lst2)
              (= (sum lst1) (sum lst2)))
  '((1 2) () (2 2) (1 2 3) (1 3) (6) (2 1) (1 2 1) (3 1 2) (5 1)))
(((2 1) (1 2)) ()) ((1 2 1) (1 3) (2 2)) ((5 1) (3 1 2) (6) (1 2 3)) )

> (partition (lambda (lst1 lst2)
              (equal? (insertion-sort < lst1)
                      (insertion-sort < lst2)))
  '((1 2) () (2 2) (1 2 3) (1 3) (6) (2 1) (1 2 1) (3 1 2) (5 1)))
(((2 1) (1 2)) ()) ((3 1 2) (1 2 3)) ((1 3)) ((6)) ((1 2 1)) ((5 1)) )
```

EC2 [20]: Improved Partitioning

Assume that $eqpred$ from EC1 has unit cost. It can be shown that the best-case running time of the `partition` function from EC1 has a worst-case quadratic asymptotic running time. This running time can be improved to $n \log n$ if the partitioning function is supplied with a less-than-or-equal-to predicate $leqpred$ (also assumed to have unit cost), and the equality predicate for partitioning is derived from $leqpred$. Define a Scheme partitioning function `(partition-leq $leqpred$ lst)` that partitions lst according to $leqpred$ and runs in $n \log n$ worst-case asymptotic time.

Note: you need not solve EC1 in order to solve EC2.

EC3 [20]: Predecessor

The predecessor function `pred` on Church numerals has the following behavior:

- if c is a Church numeral representing a non-zero integer n , then $(\text{pred } c)$ is a Church numeral representing the integer $n - 1$.
- if c is the Church numeral representing 0, then $(\text{pred } c)$ is the Church numeral 0.

For example:

```
(church->int (pred (int->church 5)))  
4
```

```
(church->int (pred (int->church 1)))  
0
```

```
(church->int (pred (int->church 0)))  
0
```

Write the `pred` function in Scheme. *Hint:* iterate over a pair of Church numerals.

Appendix A: Higher-Order List Operations

```
(define generate
  (lambda (seed next done?)
    (if (done? seed)
        '()
        (cons seed (generate (next seed) next done?))))))

(define map
  (lambda (f lst)
    (if (null? lst)
        '()
        (cons (f (car lst))
              (map f (cdr lst))))))

(define filter
  (lambda (pred lst)
    (if (null? lst)
        '()
        (if (pred (car lst))
            (cons (car lst) (filter pred (cdr lst)))
            (filter pred (cdr lst))))))

(define foldr
  (lambda (op init lst)
    (if (null? lst)
        init
        (op (car lst) (foldr op init (cdr lst))))))

(define foldl
  (lambda (op init lst)
    (if (null? lst)
        init
        (foldl op (op (car lst) init) (cdr lst))))))

(define forall?
  (lambda (pred lst)
    (if (null? lst)
        #t
        (and (pred (car lst))
              (forall pred (cdr lst))))))

(define exists?
  (lambda (pred lst)
    (if (null? lst)
        #f
        (or (pred (car lst))
            (exists? pred (cdr lst))))))

(define some
  (lambda (pred lst)
    (if (null? lst)
        #f
        (if (pred (car lst))
            (car lst)
            (some pred (cdr lst))))))
```

*Problem Set Header Page:
Please make this the first page of your submission.*

CS251 Problem Set 2 **Due Friday, February 25, 2000**

Name:

Date & Time Submitted (*only if late*):

Collaborators (*anyone you collaborated with in the process of doing the problem set*):

*In the **Time** column, please estimate the time you spent on the parts of this problem set. Please try to be as accurate as possible; this information will help me to design future problem sets. I will fill out the **Score** column when grading your problem set.*

Part	Time	Score	Part	Time	Score
General Reading			Problem 2j [7]		
Problem 1 [20]			Problem 2k [10]		
Problem 2a [5]			Problem 2l [10]		
Problem 2b [5]			Problem 2m [10]		
Problem 2c [5]			Problem 3 [20]		
Problem 2d [6]			Problem 4 [40]		
Problem 2e [7]			Problem 5 [20]		
Problem 2f [8]			Problem EC1 [20]		
Problem 2g [10]			Problem EC2 [20]		
Problem 2h [10]			Problem EC3 [20]		
Problem 2i [7]			Subtotal		
Subtotal			Total		