Haskell and HUGS

HASKELL is a lazy, statically scoped, purely functional programming language. Like OCAML, it is statically typed, but most types are automatically deduced by type reconstruction. We will explore HASKELL in the context of HUGS, a HASKELL interpreter. HUGS is only one of many HASKELL implementations; visit www.haskell.org for more information on the language and its implementations.

1 Launching HUGS

The HUGS system is launched in Linux by executing hugs at the Linux prompt. When HUGS is launched, it displays the following herald:

```
[lyn@jaguar private] hugs
```

```
|| || || || || ||__
                             Hugs 98: Based on the Haskell 98 standard
||___|| ||__|| ||__|| __||
                             Copyright (c) 1994-2001
||---||
              ___||
                             World Wide Web: http://haskell.org/hugs
Report bugs to: hugs-bugs@haskell.org
|| Version: December 2001
                             _____
Haskell 98 mode: Restart with command line option -98 to enable extensions
Reading file "/usr/share/hugs/lib/Prelude.hs":
Hugs session for:
/usr/share/hugs/lib/Prelude.hs
Type :? for help
Prelude>
```

By default, HUGS loads a large set of HASKELL libraries, known as the "prelude". These libraries are defined in the file /usr/share/hugs/lib/Prelude.hs, which you are encouraged to skim. Lots of standard functions are defined in this file, including common list utilities like map, filter, length, foldr, foldr, zip, unzip, take, and drop. The prompt Prelude> indicates that any expressions typed in will be evaluated in the context of the Prelude module containing the standard library functions.

2 Quitting HUGS

You can quit out of HUGS in one of two ways:

- 1. Execute the :quit directive within the HUGS interpreter.
- 2. Type C-d (i.e., "Control-d").

3 Interacting with HUGS

You interact with the HUGS interpreter in one of two ways:

1. By typing a HASKELL expression to be evaluated. In this case, HUGS displays the value of the expression. See Fig. 1 for some examples. Unlike the OCAML interpreter, HUGS does not display the type of the expression.

```
Prelude> 2*(3+4)
14
Prelude> head [1,2,3,4]
Prelude> tail [1,2,3,4]
[2, 3, 4]
Prelude> map (2*) [1,2,3,4]
[2,4,6,8]
Prelude> take 2 [10,20,30,40,50]
[10,20]
Prelude> drop 2 [10,20,30,40,50]
[30, 40, 50]
Prelude> foldr (+) 0 [1,2,3,4]
10
Prelude> zip [1,2,3] [10,20,30,40]
[(1,10),(2,20),(3,30)]
Prelude> unzip [(1,10),(2,20),(3,30)]
([1,2,3],[10,20,30])
Prelude> fst (1,2)
1
Prelude> snd (1,2)
2
```

Figure 1: Sample expressions evaluated in HUGS.

2. By typing a HUGS directive, along with its arguments. All directives begin with a colon. For example, the :cd directive changes the current working directory to a given directory. For example, if you execute

```
Prelude> :cd /students/your-username/cs251/ps7
```

then HUGS will interpret all following filenames relative to this directory.

Another important directive is the :type directive, which can be abbreviated :t. This displays the type of a HASKELL expression. For example:

Prelude> :type map

```
map :: (a -> b) -> [a] -> [b]
Prelude> :type foldr
foldr :: (a -> b -> b) -> b -> [a] -> bw2
Prelude> :type zip
zip :: [a] -> [b] -> [(a,b)]
Prelude> :type unzip
unzip :: [(a,b)] -> ([a],[b])
Prelude> :type "foo"
"foo" :: String
Prelude> :type "foo" == "bar"
"foo" == "bar" :: Bool
Prelude> :type 1+2
1 + 2 :: Num a => a
Prelude> :type [1,2,3]
[1,2,3] :: Num a => [a]
Prelude > :type 1 == 2
1 == 2 :: Num a => Bool
```

In the last three :type examples, the type begins with Num $a \Rightarrow \dots$ This is a so-called **qualified type**. It turns out that HASKELL has many kinds of numeric types, and integers can have any of these types. A qualified type of the form Num $a \Rightarrow t$ specifies a type t that is parameterized over any numeric type a.

By far the most important directive is the :load directive, which can be abbreviated :l. This loads the HASKELL declarations in the specified file. For example,

Prelude> :l Test.hs

loads the declarations in the file Test.hs. Note that the filename need not be delimited by double quotes, although they are allowed. The :reload directive, abbreviated :r, re-executes the most recent :load directive. For example, if Test.hs has been loaded as shown above, then :r will load the contents of Test.hs again. The :reload directive is commonly used after editing a file to add or fix a declaration.

The :quit directive exits the HUGS interpreter. The :? displays a list of all directives.

4 HASKELL Declarations

Unlike in the OCAML and MIT-SCHEME interpreters, in HUGS it is not possible to enter a declaration directly to the interpreter. Instead, all declarations must be written in files, and the :load and :reload directives are used to communicate these declarations to the HUGS interpreter.

Fig. 2 shows some representative HASKELL declarations, which we can imagine are in the file Test.hs. We will discuss these declarations in the context of some sample expressions that will be evaluated in HUGS after executing the directive :load Test.hs. Because the file Test.hs does

not have any module declarations, the declarations in the file are interpreted relative to the default Main module.

```
Prelude> :load Test.hs
Reading file "Test.hs":
Hugs session for:
/usr/share/hugs/lib/Prelude.hs
Test.hs
Main>
```

A line comment in HASKELL is introduced via the double dashes, --, and goes until the end of the line. Various comments are sprinkled throught Test.hs in Fig. 2.

In HASKELL, a name may be attached to any value via the syntax I = E, as in a = 2 + 3. Because HASKELL is a lazy language, the definition expression E (in this case, 2 + 3 is not evaluated until the the name I is required later (if ever). If it is evaluated later, the value is memoized so that it will be computed at most once. Evaluating a variable in the HUGS interpreter forces its value to be computed in order to print the value.

Main> a 5

The HASKELL syntax for abstractions is $I_{formal} \rightarrow E_{body}$, where the slash mark $\$ was chosen because it resembles a Greek λ symbol. So $x \rightarrow x*x$ is the HASKELL notation for a squaring function. The notation

 $\setminus I_1 \ldots I_n \rightarrow E_{body}$

is sugar for the curried function

 $\setminus I_1 \rightarrow \setminus I_2 \rightarrow \ldots \setminus I_n \rightarrow E_{body}$

The declaration

 I_{name} I_1 ... I_n = E_{body}

is syntactic sugar for the curried function declaration

 $I_{name} = \langle I_1 \rightarrow \langle I_2 \rightarrow \dots \rangle I_n \rightarrow E_{body}$

For example,

avg x y = x+y/2

is syntactic sugar for

avg = $\langle x \rightarrow \langle y \rightarrow (x+y)/2 \rangle$

Function application is denoted by juxtaposition of function and argument(s). For example, sq a denotes the result of applying the squaring function to the value of a. Function application is left-associative, which is consistent with curried functions. For example, avg 3 8 is parsed as (avg 3) 8.

```
Main> sq a
25
Main> avg 3 8
5.5
```

```
a = 2 + 3 - - declare variable a
sq = \langle x \rightarrow x * x - sugared form: sq x = x * x
fact 0 = 1
             -- Recursive factorial
fact n = n * fact (n-1)
factIter n = loop n 1 -- Iterative factorial
  where loop 0 ans = ans
        loop num ans = loop (num-1) (num*ans)
isEven 0 = True -- Mutually recursive functions isEven and isOdd
isEven m = isOdd (m - 1)
isOdd 0 = False
isOdd n = isEven (n - 1)
sumList = foldr (+) 0 -- list summation function
nats = 0 : (map (1+) nats) -- infinite list of natural numbers
twos = 1 : (map (2*) twos) -- infinite list of powers of two
fibs = 0 : 1 : (zipWith (+) fibs (tail fibs)) -- infinite list of Fibonacci numbers
data Tree a = Leaf | Node (Tree a, a, Tree a) -- tree datatype declaration
  deriving (Show, Eq) -- show and equality (==) functions on trees
testTree = Node(Node(Leaf, 4, Leaf),
                     1,
                     Node(Node(Leaf,5,Leaf),
                          2,
                          Leaf)),
                6,
                Node(Leaf, 3, Node(Leaf, 7, Leaf)))
value (Node(_,v,_)) = v -- accessor functions for tree nodes
left (Node(1,_,_)) = 1
right (Node(\_,\_,r)) = r
height Leaf = 0
height (Node(l, r)) = 1 + max (height 1) (height r)
treeSum Leaf = 0
treeSum (Node(1,v,r)) = (treeSum 1) + v + (treeSum r)
treeMap f Leaf = Leaf
treeMap f (Node(1,v,r)) = Node(treeMap f 1, f v, treeMap f r)
-- infinite tree of integers in which every node has
-- its binary address as its value.
intTree = makeTree 0
  where makeTree n = Node(makeTree (2*n), n, makeTree ((2*n)+1))
```

Figure 2: Sample HUGS declarations in the file Test.hs.

Like OCAML, HASKELL has a case-based pattern-matching construct, which has the form:

case $E_{discriminant}$ of $P_1 \rightarrow E_1$ \vdots $P_n \rightarrow E_n$

As an example, here is the definition of a swapList function that swaps the first two elements of a list with at least two elements:

```
swapList = \ xs -> case xs of
    [] -> []
    [x] -> [x]
    x:y:zs -> y:x:zs
```

Note that the clauses of the **case** construct are not separated by any sort of syntax (like the vertical bar that separates **match** clauses in OCAML). This is because HASKELL, unlike almost every other modern language, actually uses indedentation and whitespace to as a disambiguation aid in parsing.

It is rare to see explicit **case** constructs in HASKELL programs, because they are usually written in a sugared form as a sequence of function definitions with different patterns in the parameter position(s). For example, the sugared form of the above **swapList** function is:

```
swapList [] = []
swapList [x] = [x]
swapList (x:y:zs) = (y:x:zs)
```

All names declared in a file are defined in a single recursive scope. In Fig. 2, fact is an example of a recursive function definition, isEven and isOdd are mutually recursive functions, and nats, twos, and fibs are recursively defined infinite lists. Mutually recursive definitions – especially of non-function values – are much easier to handle in a lazy language than in a strict one. Local recursive bindings are introduced in HASKELL via the where clause, which appears in the factIter and intTree declarations in Fig. 2. The where clause is HASKELL's version of Scheme's letrec, OCAML's let rec, and HOILIC's bindrec. Interestingly, the concrete syntax of where has the local declarations following the body rather then preceding it.

```
Main> fact 5
120
Main> factIter 6
720
Main> isEven 10
True
Main> isOdd 10
False
Main> sumList [1,2,3, 4]
10
Main> take 20 nats
[0,1,2,3,4,5,6,7,8,9,10,11,12,13,14,15,16,17,18,19]
Main> take 15 fibs
[0,1,1,2,3,5,8,13,21,34,55,89,144,233,377]
```

The portions after this point are still under construction!

HASKELL supports OCAML-like sum-of-product data types via the data declaration. Here is a binary tree data type:

```
data Tree a = Leaf | Node (Tree a, a, Tree a)
  deriving (Show, Eq)
```

There are two tree constructors: the nullary Leaf constructor, and the unary Node constructor, which takes a triple of the left subtree, the root value, and the right subtree. The declaration deriving (Show, Eq) tells HASKELL that string representations of trees (via the show function) and equality on trees (via the == function) should be automatically defined in a structural way. Here is a sample tree:

Tree operations can be defined via pattern matching. For example, we can define functions that access the three parts of a tree node:

```
value (Node(_,v,_)) = v -- accessor functions for tree nodes
left (Node(1,_,_)) = l
right (Node(_,_,r)) = r
```

For example:

```
Main> testTree
Node (Node (Leaf,4,Leaf),1,Node (Node (Leaf,5,Leaf),2,Leaf)),6,Node (Leaf,3,Node (Leaf,7,Leaf
Main> value testTree
6
Main> left testTree
Node (Node (Leaf,4,Leaf),1,Node (Node (Leaf,5,Leaf),2,Leaf))
Main> right testTree
Node (Leaf, 3, Node (Leaf, 7, Leaf))
Main> height testTree
Ο
Main> :r
Reading file "Test.hs":
Hugs session for:
/usr/share/hugs/lib/Prelude.hs
Test.hs
Main> height testTree
4
Main> treeSum testTree
28
```

Main> treeSum (treeMap (2*) testTree) 56 Main>