

The Pros of cons:
Programming with Pairs and Lists



**CS251 Programming
Languages**
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Racket Values

- booleans: `#t`, `#f`
- numbers:
 - integers: `42`, `0`, `-273`
 - rationals: `2/3`, `-251/17`
 - floating point (including scientific notation):
`98.6`, `-6.125`, `3.141592653589793`, `6.023e23`
 - complex: `3+2i`, `17-23i`, `4.5-1.4142i`

Note: some are *exact*, the rest are *inexact*. See docs.
- strings: `"cat"`, `"CS251"`, `"αβγ"`,
`"To be\nor not\nto be"`
- characters: `#\a`, `#\A`, `#\5`, `#\space`, `#\tab`, `#\newline`
- anonymous functions: `(lambda (a b) (+ a (* b c)))`

What about compound data?

cons Glues Two Values into a Pair

A new kind of value:

- pairs (a.k.a. cons cells): (cons **v1** **v2**)

e.g.,

– (cons 17 42)

– (cons 3.14159 #t)

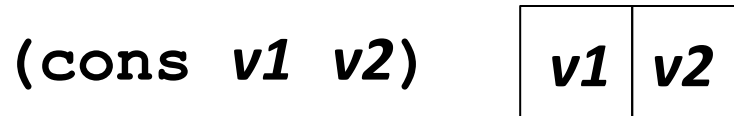
– (cons "CS251" (λ (x) (* 2 x)))

– (cons (cons 3 4.5) (cons #f #\a))

In Racket,
type Command-\
to get λ char

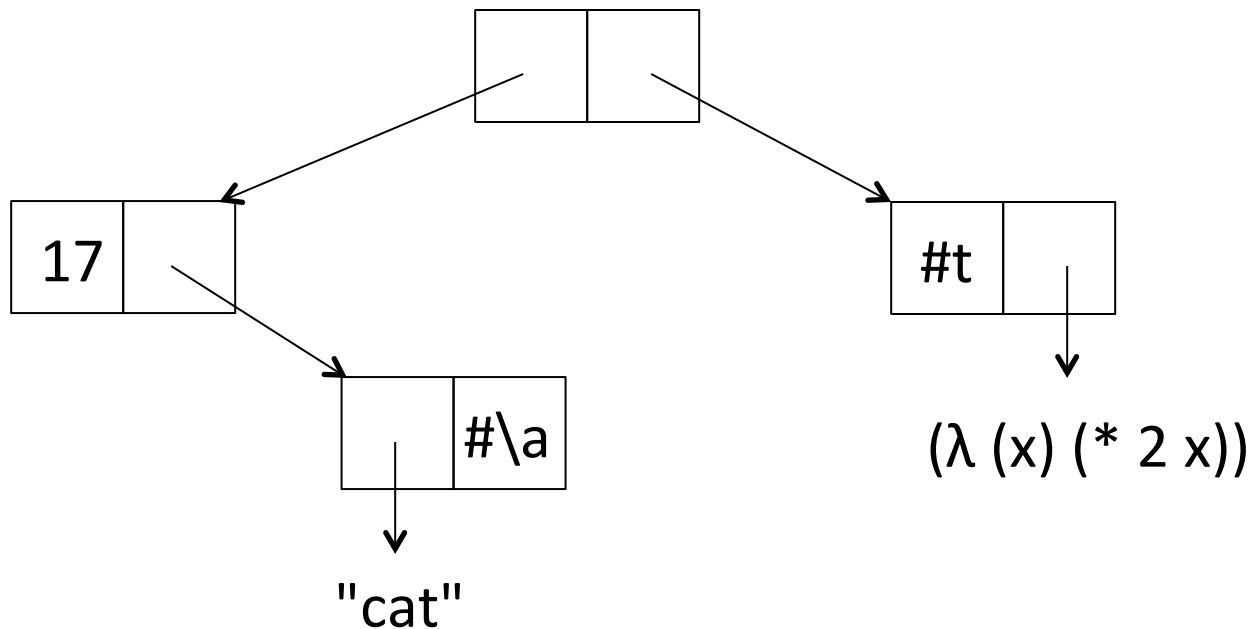
Can glue any number of values into a cons tree!

Box-and-pointer diagrams for cons trees



Convention: put “small” values (numbers, booleans, characters) inside a box, and draw a pointers to “large” values (functions, strings, pairs) outside a box.

```
(cons (cons 17 (cons "cat" #\a))  
      (cons #t (λ (x) (* 2 x))))
```



Evaluation Rules for cons

Big step semantics:

$$\frac{\begin{array}{l} \mathbf{e1} \downarrow \mathbf{v1} \\ \mathbf{e2} \downarrow \mathbf{v2} \end{array}}{\mathbf{(cons\ e1\ e2)} \downarrow \mathbf{(cons\ v1\ v2)}} \text{ (cons)}$$

Small-step semantics:

$\mathbf{(cons\ e1\ e2)}$

\Rightarrow^* $\mathbf{(cons\ v1\ e2)}$; first evaluate $\mathbf{e1}$ to $\mathbf{v1}$ step-by-step

\Rightarrow^* $\mathbf{(cons\ v1\ v2)}$; then evaluate $\mathbf{e2}$ to $\mathbf{v2}$ step-by-step

cons evaluation example

```
(cons (cons (+ 1 2) (< 3 4))  
      (cons (> 5 6) (* 7 8)))
```

```
⇒ (cons (cons 3 (< 3 4))  
        (cons (> 5 6) (* 7 8)))
```

```
⇒ (cons (cons 3 #t) (cons (> 5 6) (* 7 8)))
```

```
⇒ (cons (cons 3 #t) (cons #f (* 7 8)))
```

```
⇒ (cons (cons 3 #t) (cons #f 56))
```

car and cdr

- `car` extracts the left value of a pair

`(car (cons 7 4)) ⇒ 7`

- `cdr` extract the right value of a pair

`(cdr (cons 7 4)) ⇒ 4`

Why these names?

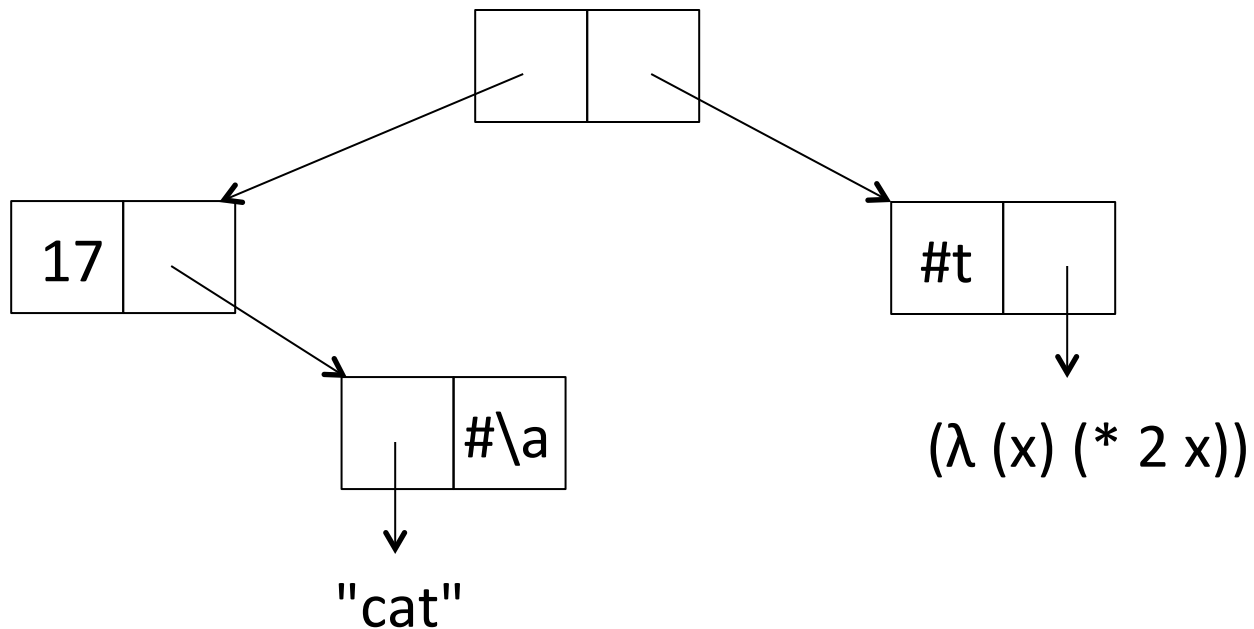
- `car` from “contents of address register”
- `cdr` from “contents of decrement register”



Practice with `car` and `cdr`

Write expressions using `car`, `cdr`, and `tr` that extract the five leaves of this tree:

```
(define tr
  (cons (cons 17 (cons "cat" #\a))
        (cons #t (lambda (x) (* 2 x)))))
```



cadr and friends

- `(caar e)` means `(car (car e))`
- `(cadr e)` means `(car (cdr e))`
- `(cdar e)` means `(cdr (car e))`
- `(cddr e)` means `(cdr (cdr e))`
- `(caaar e)` means `(car (car (car e)))`
- \vdots
- `(cddddr e)` means `(cdr (cdr (cdr (cdr (cdr e)))))`

Evaluation Rules for `car` and `cdr`

Big-step semantics:

$$\frac{e \downarrow (\text{cons } v1 \ v2)}{(\text{car } e) \downarrow v1} \quad (\text{car})$$

$$\frac{e \downarrow (\text{cons } v1 \ v2)}{(\text{cdr } e) \downarrow v2} \quad (\text{cdr})$$

Small-step semantics:

`(car e)`

\Rightarrow^* `(car (cons v1 v2))`; first evaluate `e` to pair step-by-step

\Rightarrow `v1`; then extract left value of pair

`(cdr e)`

\Rightarrow^* `(cdr (cons v1 v2))`; first evaluate `e` to pair step-by-step

\Rightarrow `v2`; then extract right value of pair

Semantics Puzzle

According to the rules on the previous page, what is the result of evaluating this expression?

```
(car (cons (+ 2 3) (* 5 #t)))
```

Note: there are two “natural” answers. Racket gives one, but there are languages that give the other one!

Printed Representations in Racket Interpreter

```
> (lambda (x) (* x 2))  
#<procedure>  
  
> (cons (+ 1 2) (* 3 4))  
'(3 . 12)  
  
> (cons (cons 5 6) (cons 7 8))  
'((5 . 6) 7 . 8)  
  
> (cons 1 (cons 2 (cons 3 4)))  
'(1 2 3 . 4)
```

What's going on here?

Display Notation and Dotted Pairs

- The **display notation** for `(cons v1 v2)` is `(dn1 . dn2)`, where ***dn1*** and ***dn2*** are the display notations for ***v1*** and ***v2***
- In display notation, a dot “eats” a paren pair that follows it directly:

`((5 . 6) . (7 . 8))`

becomes `(5 . 6) 7 . 8)`

`(1 . (2 . (3 . 4)))`

becomes `(1 . (2 3 . 4))`

becomes `(1 2 3 . 4)`

Why? Because we'll see this makes lists print prettily.

- The Racket interpreter puts a single quote mark before the display notation of a top-level pair value. (We'll say more about quotation later.)

display vs. print in Racket

```
> (display (cons 1 (cons 2 null)))  
(1 2)
```

```
> (display (cons (cons 5 6) (cons 7 8)))  
((5 . 6) 7 . 8)
```

```
> (display (cons 1 (cons 2 (cons 3 4))))  
(1 2 3 . 4)
```

```
> (print(cons 1 (cons 2 null)))  
'(1 2)
```

```
> (print(cons (cons 5 6) (cons 7 8)))  
'((5 . 6) 7 . 8)
```

```
> (print(cons 1 (cons 2 (cons 3 4))))  
'(1 2 3 . 4)
```

Functions Can Take and Return Pairs

```
(define (swap-pair pair)
  (cons (cdr pair) (car pair)))
```

```
(define (sort-pair pair)
  (if (< (car pair) (cdr pair))
      pair
      (swap pair)))
```

What are the values of these expressions?

- (swap-pair (cons 1 2))
- (sort-pair (cons 4 7))
- (sort-pair (cons 8 5))

Lists

In Racket, a **list** is just a recursive pattern of pairs.

A list is either

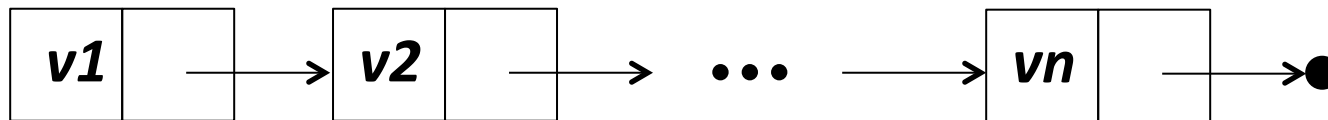
- The empty list `null`, whose display notation is `()`
- A nonempty list `(cons v_first v_rest)` whose
 - first element is *v_first*
 - and the rest of whose elements are the sublist *v_rest*

E.g., a list of the 3 numbers 7, 2, 4 is written

```
(cons 7 (cons 2 (cons 4 null)))
```

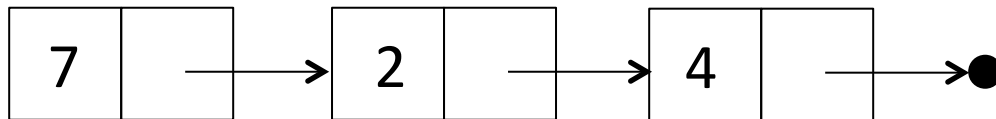

Box-and-pointer notation for lists

A list of n values is drawn like this:



Notation for null
in box-and-pointer
diagrams

For example:



list sugar

Treat `list` as syntactic sugar:

- `(list)` desugars to `null`
- `(list e1 ...)` desugars to `(cons e1 (list ...))`

For example:

```
(list (+ 1 2) (* 3 4) (< 5 6))
```

```
desugars to (cons (+ 1 2) (list (* 3 4) (< 5 6)))
```

```
desugars to (cons (+ 1 2) (cons (* 3 4) (list (< 5 6))))
```

```
desugars to (cons (+ 1 2) (cons (* 3 4) (cons (< 5 6) (list))))
```

```
desugars to (cons (+ 1 2) (cons (* 3 4) (cons (< 5 6) null)))
```

* This is a white lie, but we can pretend it's true for now

Display Notation for Lists

The “dot eats parens” rule makes lists display nicely:

```
(list 7 2 4)
```

desugars to `(cons 7 (cons 2 (cons 4 null)))`

displays as (before rule) `(7 . (2 . (4 . ())))`

displays as (after rule) `(7 2 4)`

prints as `'(7 2 4)`

In Racket:

```
> (display (list 7 2 4))  
(7 2 4)
```

```
> (display (cons 7 (cons 2 (cons 4 null))))  
(7 2 4)
```

list and small-step evaluation

It is sometimes helpful to both desugar and resugar with `list`:

```
(list (+ 1 2) (* 3 4) (< 5 6))
```

```
desugars to (cons (+ 1 2) (cons (* 3 4) (cons (< 5 6) null)))
```

```
⇒ (cons 3 (cons (* 3 4) (cons (< 5 6) null)))
```

```
⇒ (cons 3 (cons 12 (cons (< 5 6) null)))
```

```
⇒ (cons 3 (cons 12 (cons #t null)))
```

```
resugars to (list 3 12 #t)
```

Heck, let's informally write this as:

```
(list (+ 1 2) (* 3 4) (< 5 6))
```

```
⇒ (list 3 (* 3 4) (< 5 6))
```

```
⇒ (list 3 12 (cons (< 5 6)))
```

```
⇒ (list 3 12 #t)
```

first, rest, and friends

- `first` returns the first element of a list:

```
(first (list 7 2 4)) ⇒ 7
```

(`first` is almost a synonym for `car`, but requires its argument to be a list)

- `rest` returns the sublist of a list containing every element but the first:

```
(rest (list 7 2 4)) ⇒ (list 2 4)
```

(`rest` is almost a synonym for `cdr`, but requires its argument to be a list)

- Also have `second`, `third`, ..., `ninth`, `tenth`

Recursive List Functions

Because lists are defined recursively, it's natural to process them recursively.

Typically (but not always) a recursive function `recf` on a list argument `L` has two cases:

- **base case:** what does `recf` return when `L` is empty?
(Use `null?` to test for an empty list)
- **recursive case:** if `L` is the nonempty list `(cons v_first v_rest)` how are `v_first` and `(recf v_rest)` combined to give the result for `(recf L)`?

Note that we ``blindly'' apply `recf` to `v_rest`!

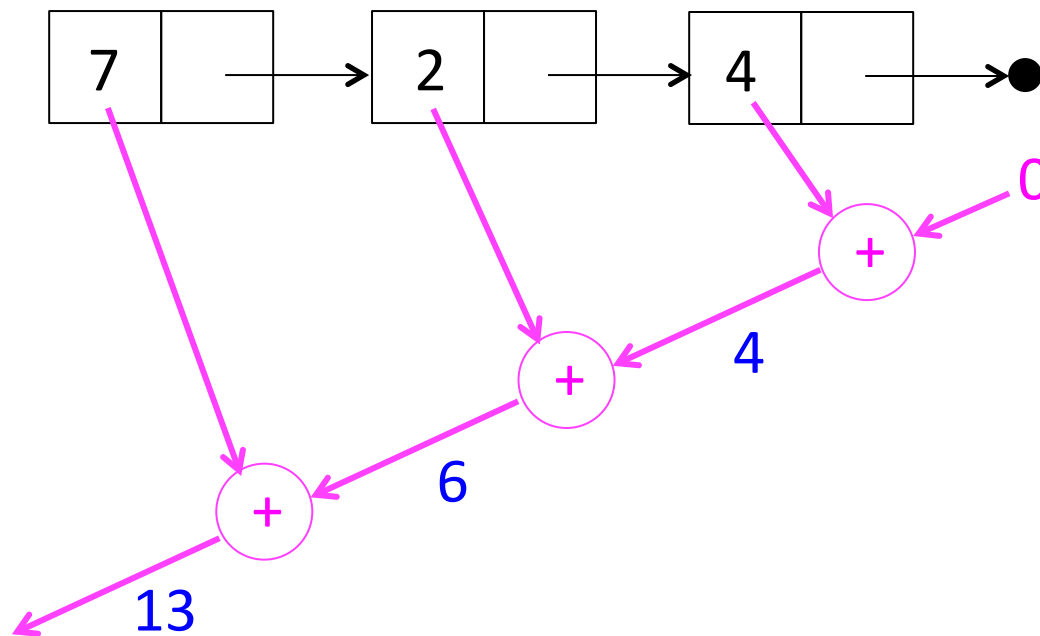
Example: sum

`(sum ns)` returns the sum of the numbers in the list `ns`

```
(define (sum ns)
  (if (null? ns)
      0
      (+ (first ns)
         (sum (rest ns)))))
```

Understanding `sum`: Approach #1

```
(sum (list 7 2 4))
```



We'll call this the **recursive accumulation** pattern

Understanding `sum`: Approach #2

In `(sum (list 7 2 4))`, the list argument to `sum` is

```
(cons 7 (cons 2 (cons 4 null)))
```

Replace `cons` by `+` and `null` by `0` and simplify:

```
(+ 7 (+ 2 (+ 4 0)))
```

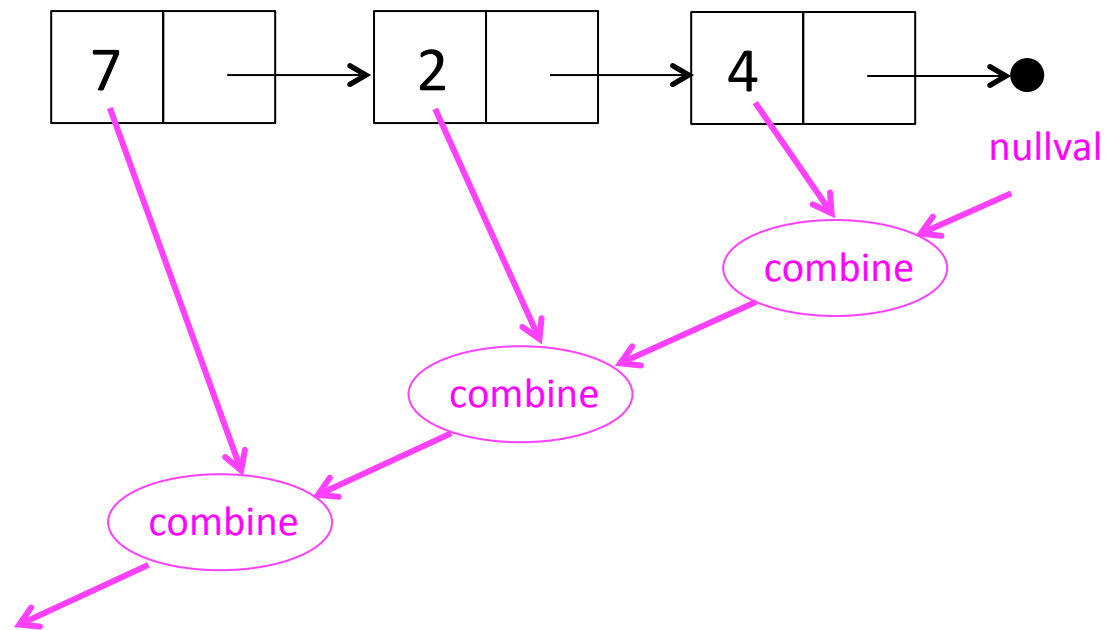
```
⇒ (+ 7 (+ 2 4))
```

```
⇒ (+ 7 6)
```

```
⇒ 13
```

Generalizing sum: Approach #1

```
(recf (list 7 2 4))
```



Generalizing `sum`: Approach #2

In `(recf (list 7 2 4))`, the list argument to `recf` is

```
(cons 7 (cons 2 (cons 4 null)))
```

Replace `cons` by **`combine`** and `null` by **`nullval`** and simplify:

```
(combine 7 (combine 2 (combine 4 nullval)))
```

Generalizing the `sum` definition

```
(define (recf ns)
  (if (null? ns)
      nullval
      (combine (first ns)
                (recf (rest ns)))))
```

Your turn

`(product ns)` returns the product of the numbers in `ns`

`(min-list ns)` returns the minimum of the numbers in `ns`

Hint: use `min` and `+inf.0` (positive infinity)

`(max-list ns)` returns the maximum of the numbers in `ns`

Hint: use `max` and `-inf.0` (negative infinity)

`(all-true? bs)` returns `#t` if all the elements in `bs` are truthy; otherwise returns `#f`. *Hint: use `and`*

`(some-true? bs)` returns a truthy value if at least one element in `bs` is truthy; otherwise returns `#f`. *Hint: use `or`*

`(my-length xs)` returns the length of the list `xs`

Recursive Accumulation Pattern Summary

	combine	nullval
sum	+	0
product	*	1
min-list	min	+inf.0
max-list	max	-inf.0
all-true?	and	#t
some-true?	or	#f
my-length	(λ (fst subres) (+ 1 subres))	0

Mapping Example: `map-double`

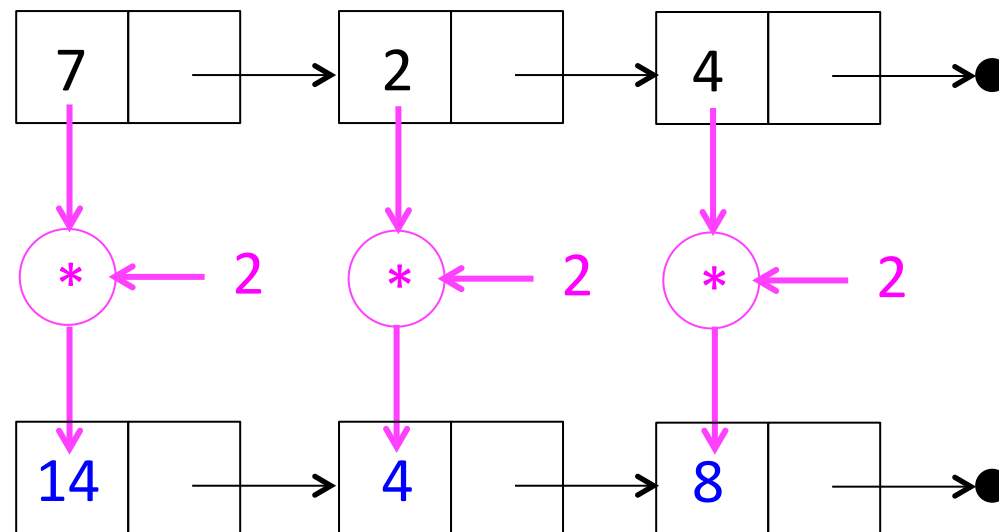
`(map-double ns)` returns a new list the same length as `ns` in which each element is the double of the corresponding element in `ns`.

```
> (map-double (list 7 2 4))  
'(14 4 8)
```

```
(define (map-double ns)  
  (if (null? ns)  
      ; Flesh out base case  
  
      ; Flesh out recursive case  
  
      ))
```

Understanding map-double

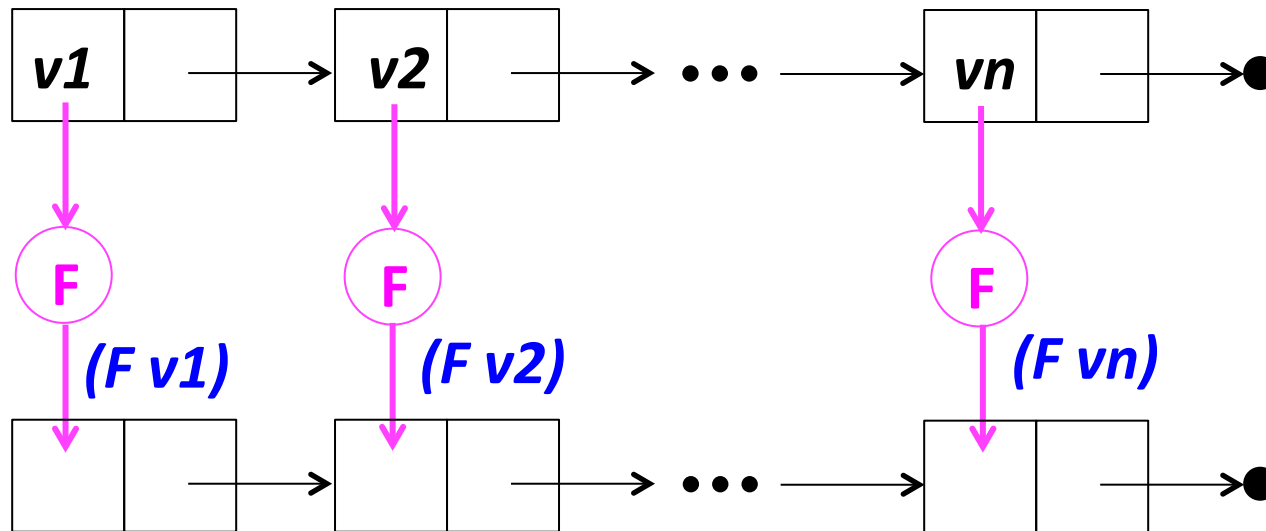
```
(map-double (list 7 2 4))
```



We'll call this the **mapping** pattern

Generalizing map-double

(map **F** (list **v1** **v2** ... **vn**))



```
(define (map F xs)
  (if (null? xs)
      null
      (cons (F (first xs))
            (map F (rest xs)))))
```

Expressing `mapF` as an accumulation

```
(define (mapF xs)
  (if (null? xs)
      null
      ((λ (fst subres)
         ) ; Flesh this out
       (first xs)
       (mapF (rest xs)))))
```

Some Recursive Listfuns Need Extra Args

```
(define (map-scale factor ns)
  (if (null? ns)
      null
      (cons (* factor (first ns))
            (map-scale factor (rest ns)))))
```

Filtering Example: `filter-positive`

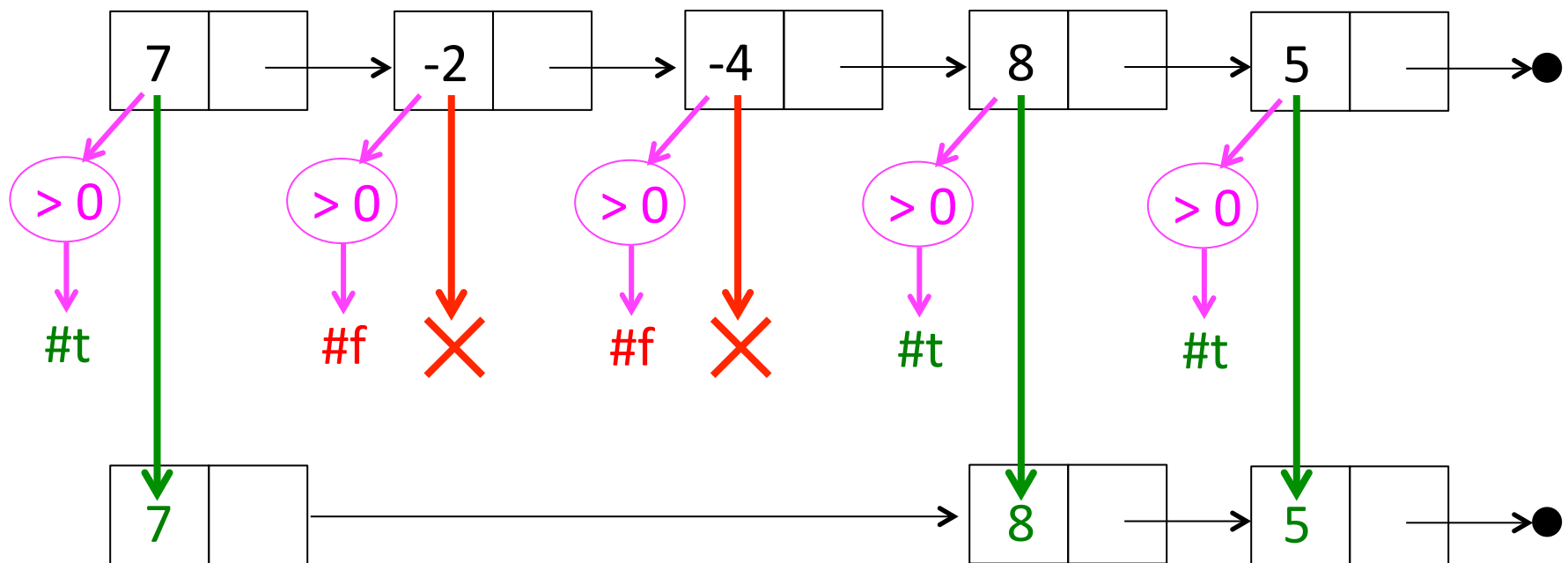
`(filter-positive ns)` returns a new list that contains only the positive elements in the list of numbers `ns`, in the same relative order as in `ns`.

```
> (filter-positive (list 7 -2 -4 8 5))  
' (7 8 5)
```

```
(define (filter-positive ns)  
  (if (null? ns)  
      ; Flesh out base case  
  
      ; Flesh out recursive case  
  
      ))
```

Understanding filter-positive

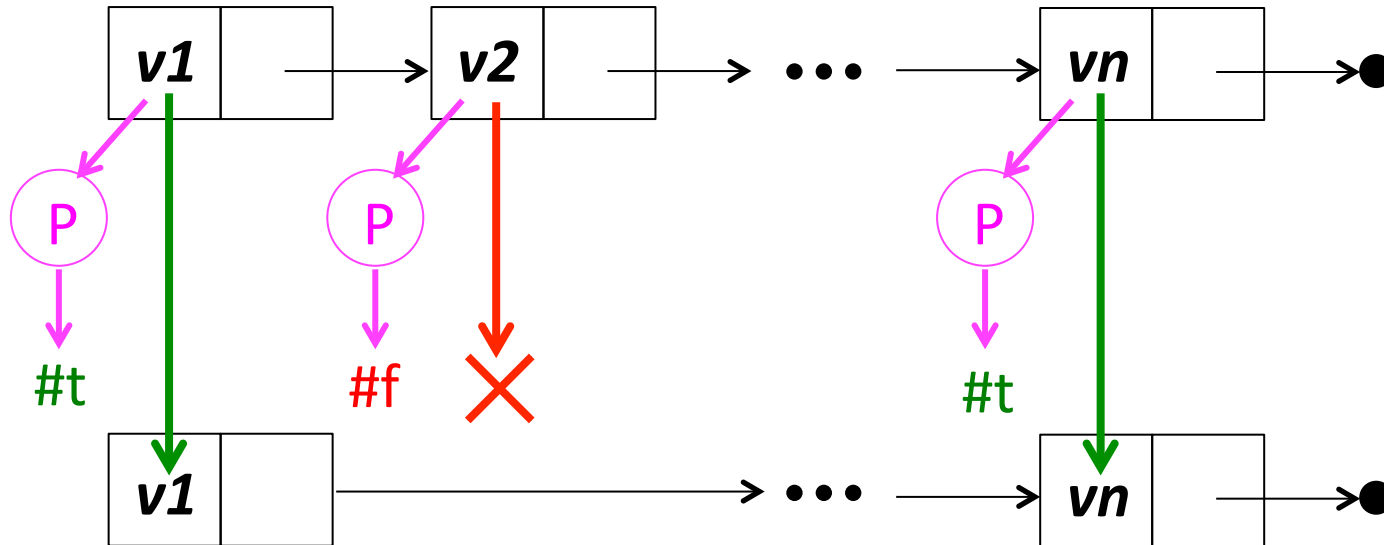
```
(filter-positive (list 7 -2 -4 8 5))
```



We'll call this the **filtering** pattern

Generalizing filter-positive

(filter **P** (list **v1 v2 ... vn**))



```
(define (filter P xs)
  (if (null? xs)
      null
      (if (P (first xs))
          (cons (first xs) (filter P (rest xs)))
          (filter P (rest xs)))))
```

Expressing `filterP` as an accumulation

```
(define (filterP xs)
  (if (null? xs)
      null
      ((lambda (fst subres)
         ) ; Flesh this out
        (first xs)
        (filterP (rest xs))))))
```

More examples

- snoc/postpend
- append
- append-all
- sorted?
- merge