CS251 Jeopardy
Spring 2002
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What data structure is commonly used in interpreters to associate names with values?
Data 2

What feature in Scheme, ML, and Java is responsible for reclaiming storage used by values that are no longer accessible from the program?
ML’s datatype and Haskell’s data construct are examples of "sum-of-product" data type declarations. What are traditional names for "sum" and "product " in programming languages?
What is the value of the following ML program?

```
let val yourMom = [[1,2], [3,4,5,6,7], [8]]
in map (foldr (fn (_,x) => 1+x) 0) yourMom
end
```
What problem does invoking the following C function lead to?

```c
int* elts (int c, int n) {
    int a[n];
    int i;
    for (i = 0; i < n; i++) {
        a[i] = c*i;
    }
    return a;
}
```

Extra: How can the problem be fixed?
Naming 1

List all of the free variables of the following HOFL expression:

\[(\text{abs } (a) \\
   (a \ b \ (\text{abs } (b) \ (+ \ b \ c))))\]
Naming 2

List *all* of the following languages that are block structured:

- Pascal
- C
- Java
- Scheme
- ML
The following is a legal program in both FOBS and HOFL. In FOBS, it denotes the factorial function, while in HOFL it does not. What programming language feature accounts for the difference in meaning between the two languages?

```
(program (n)
  (funrec ((fact (fact)
    (if (= fact 0)
      1
      (* fact (fact (- fact 1))))))
  (fact n)))
```
Give the value of the following expression in both lexically scoped and dynamically scoped versions of Scheme:

(let ((a 1)
       (b 2))
  (let ((f (let ((a 10))
              (lambda () (+ a b))))
       (let ((b 20))
         (f))))

Back
Naming 5

Give the value of the following Scheme expression under all four parameter passing mechanisms: call-by-value, call-by-name, call-by-need, call-by-reference. Assume procedure arguments are evaluated in left-to-right order.

(let ((a 1))
  (let ((b a))
    (let ((c (begin (set! a (* a 2) a))))
      (begin (set! b 10)
              (+ a (+ c c))))))

Back
Laziness 1

Which one of the following does not belong:

- lazy data
- call-by-value
- memoization
- call-by-need.

Back
In his paper "Why Functional Programming Matters", John Hughes argues that laziness is important because it enhances something? What?
Below are two definitions of an if0 construct: the first defined by desugaring, the second defined as a function:

(1) \(\text{if0 } E_{\text{num}} \ E_{\text{zero}} \ E_{\text{nonzero}}\) desugars to \(\text{if } (= E_{\text{num}} 0) \ E_{\text{zero}} \ E_{\text{nonzero}}\)

(2) (define if0
      (lambda (Enum Ezero Enonzero)
        (if (= Enum 0) Ezero Enonzero)))

List all of the following parameter-passing mechanisms under which the two definitions are equivalent:
call-by-value call-by-name call-by-need
What are the elements of the list returned by evaluating the following Haskell expression?

```haskell
take 5 (scanl (+) 0 elts)
where elts = 1 : (map (2 *) elts)
```
What is the value of the following statically-scoped call-by-value Scheme expression? Assume left-to-right operand evaluation.

(let ((n 0))
  (let ((inc! (lambda (x)
                (begin (set! n (+ n x)) n))))
   (let ((inc1 (lambda () (inc! 1))))
     (let ((inc2 (delay (inc! 2))))
       (+ (* (inc1) (force inc2))
          (* (inc1) (force inc2)))))))

Extra: What if the operand evaluation order is right-to-left?
Xforms 1

What common program transformation have we studied that Alan Perlis once quipped could cause "cancer of the semi-colon"?

Back
What is the name of a transformation that can transform an ML function of type

\[ \text{int} \times \text{char} \rightarrow \text{bool} \]

to a function of type

\[ \text{int} \rightarrow \text{char} \rightarrow \text{bool} \]
Consider the following program transformation:

\[(+ E E) \implies (* 2 E)\]

For each of the following programming paradigms, indicate whether the above transformation is safe - that is, it preserves the meaning of the expression for all possible expressions \(E\).

- purely functional
- imperative
- object-oriented
Consider the following transformation in an imperative version of Scheme:

```
((lambda (x) 3) E) => 3
```

List all of the following parameter passing mechanisms for which the above transformation is safe - that is, it preserves the meaning of the expression for all possible expressions $E$.

- call-by-value
- call-by-name
- call-by-need
- call-by-reference
In Scheme, the special form \((\text{or } E_1 \ E_2)\) first evaluates \(E_1\) to a value \(V_1\). If \(V_1\) is not false, it is returned without evaluating \(E_2\). If \(V_1\) is false, the value of \(E_2\) is returned. Bud Lojack suggests the following desugaring rule for \((\text{or } E_1 \ E_2)\):\n
\[
(\text{or } E_1 \ E_2)
\]

\textit{desugars to} \n
\[
(\text{let } ((x \ E_1)) (\text{if } x \ x \ E_2))
\]

Unfortunately, this desugaring has a bug. Give a concrete expression in which Bud’s desugaring fails to have the right meaning.
Imperative 1

List all of the following languages in which a variable is always bound to an explicit mutable cell.

- Scheme
- ML
- Java
- Haskell
- C

Back
Imperative 2

What programming language property corresponds to the mathematical notion of "substituting equals for equals" (Functional languages have it; imperative languages don’t.)
What is the value of executing $f(5)$, where $f$ is the following C function?

```c
int f (int n) {
    int ans = 1;
    while (n > 0) {
        n = n - 1;
        ans = n * ans;
    }
    return ans;
}
```
What is the value of executing $g(1, 2)$ in the context of the following C definitions?

```c
void h (int x, int* y) {
    x = x + *y;
    *y = *y + x;
}

int g (int a, int b) {
    h(a, &b);
    return a * b;
}
```
What is the value of the following Scheme program? Assume operands are evaluated from left to right. (Hint: draw environments!)

```
(let ((f (let ((a 0))
               (lambda ()
                 (begin (set! a (+ a 1))
                        (let ((b 0))
                            (lambda ()
                              (begin (set! b (+ a b))
                                     b)))))))))

(let ((p (f)))
 (+ (p)
    (let ((q (f))
           (+ (q)
              (+ (p) (q))))))
```
Control 1

Name the property that allows Scheme to perform iterations in constant space without explicit looping constructs.

Back
Which one of the following most closely models Pascal’s `goto` construct?

- Scheme’s `error` construct
- Scheme’s `call-with-current-continuation` construct
- ML’s `raise` construct
- Java’s `try/catch` construct
- Java’s `break` construct
What is the value of the following expression in a version of Scheme supporting raise and handle?

\[
\text{(handle err (lambda (y) (+ y 200)))}
\]
\[
\text{(let ((f (lambda (x) (+ (raise err x) 1000))))}
\]
\[
\text{(handle err (lambda (z) (+ z 50)))}
\]
\[
\text{(f 4))}
\]

*Extra:* what if the handles are replaced by traps?
Consider the following procedure in a version of Scheme supporting `label` and `jump`:

```
(define test
  (lambda (x)
    (+ 1 (label a
          (+ 20 (label b
                  (+ 300 (jump a
                           (label c
                             (if (> x 0)
                               (+ 4000 (jump c x))
                             (jump b x))))))))))
```

What is the value of the expression

```
(+ (test 0) (test 5))
```

Assume operands are evaluated left-to-right.
What is the value of the following expression in a version of Scheme supporting label and jump?

\[
\text{(let ((twice (\lambda (f) (\lambda (x) (f (f x))))))}
\text{(let ((g (label a (\lambda (z) (jump a z))))))}
\text{(((g twice) 1+) 0))}
\]
Name a "real-world" statically-typed language that does not require explicit types.

Back
What feature is lacking in Java’s type system that makes it impossible to write a general Scheme or ML style `map` function in Java?
What type would the ML type reconstructor infer for the following function definition?

fun some pred [] = NONE  
| some pred (x::xs) =  
  if (pred x)  
  then SOME(x)  
  else some pred xs
Types 4

Write an explicitly typed HOFLEMT expression that has the following type. The function must actually use each of its arguments.

\[
(\rightarrow (\text{int} \ (\rightarrow (\text{int}) \ \text{bool})) \ \text{bool})
\]
Types 5

Translate the following (implicitly typed) HOFLIPT expression into an explicitly typed HOFLEPT expression with the most general possible type.

\[(\text{abs } (f \ x) (f \ x \ x))\]

Extra: What is the type of your HOFLEPT expression?

Back
Complete the following Guy Steele poem by filling in the ???:

A one slot cons is called a ???
A two-slot cons makes lists as well
And I would bet a coin of bronze
There isn’t any three-slot cons.
Who was the inventor of the lambda calculus, a formal system upon which functional programming is based?
Is it possible to write an interpreter for an imperative language in a purely functional language?
Fill in the ??? in the following Norman Adams quote: “Objects are a poor man’s ???”.
List five properties that values must have in order to be considered "first-class".