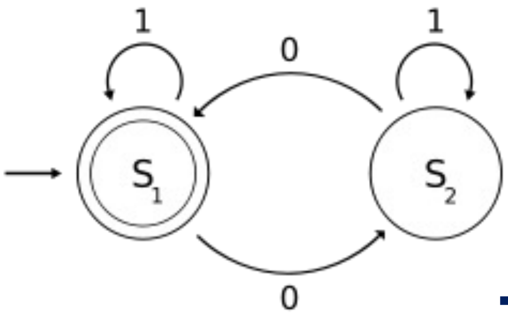
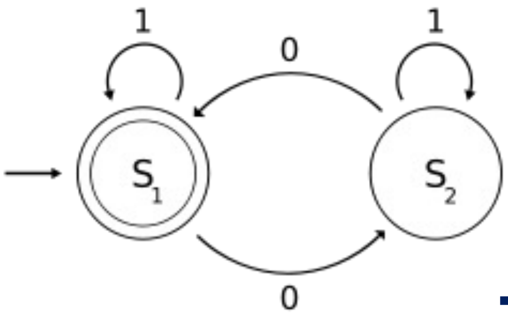


Nonregular Languages

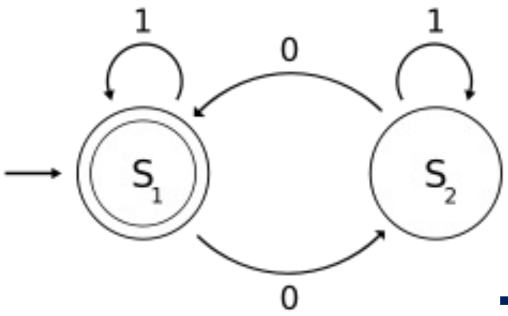
Sipser: Section 1.4 pages 77 - 82



How Do We Know What We Are Missing?

Regular languages may be specified either by regular expressions or by deterministic or nondeterministic finite automata.

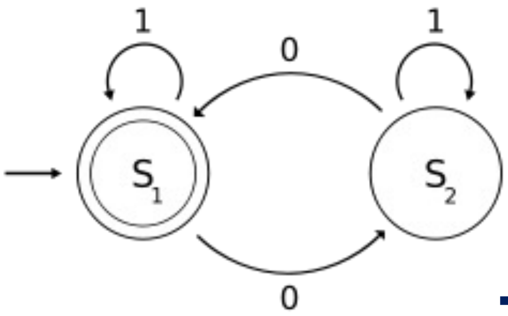
How do we show a language is not regular?



Bounded Memory

Since finite automata are not allowed to back up, the amount of memory required to determine whether or not a string is in the language must be bounded.

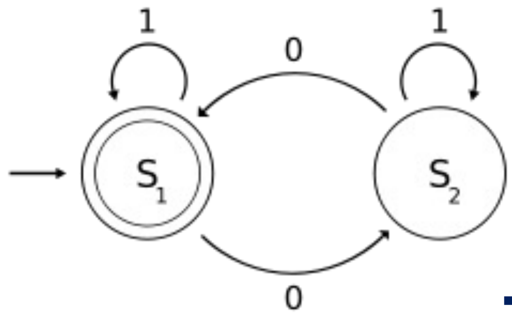
For example, consider $L = \{0^n 1^n : n \geq 0\}$.



That is a Thought, Not a Proof!

$C = \{w \mid w$ has an equal number of 0s and 1s$\}$.

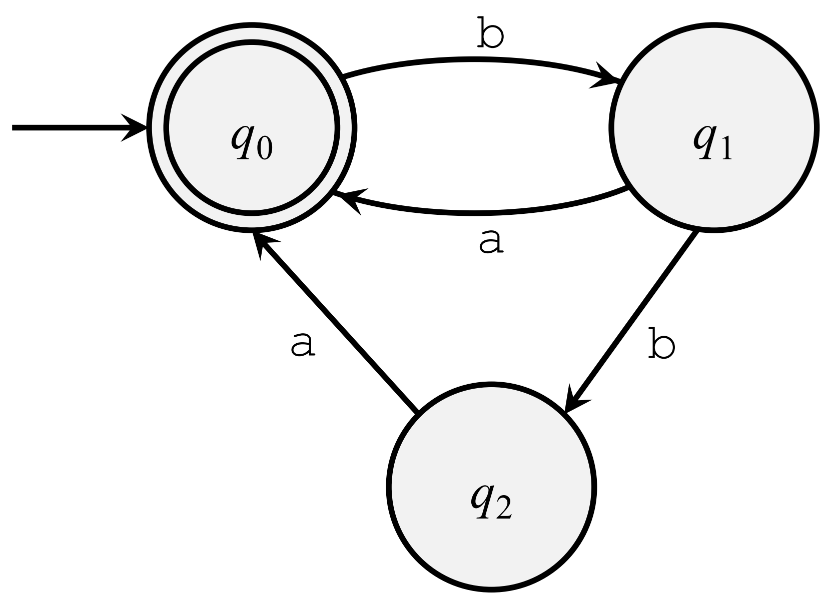
$D = \{w \mid w$ has an equal number of occurrences of 01 and 10 as substrings$\}$.



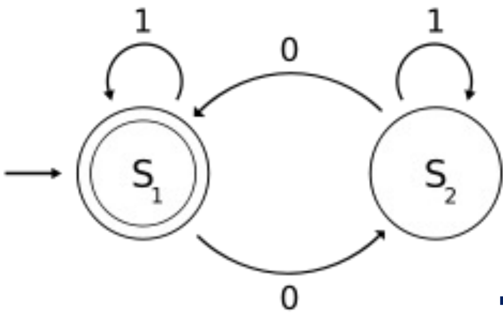
Look at it Another Way

From the specification point of view, a regular language is infinite if and only if its corresponding regular expression contains a Kleene star.

Kleene stars correspond to loops in finite automata.



Both Kleene stars and loops give rise to simple repetitive patterns in the language.



The Pigeonhole Principle

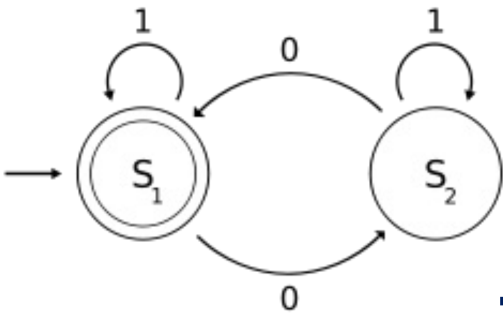
Let $M = (Q, \Sigma, \delta, q_0, F)$ be a DFA accepting an infinite language L. Suppose $|Q| = p$.

Let $w = w_1 w_2 \cdots w_p \in L$.

w = w_1 w_2 w_3 w_4 w_5 w_6 w_7 … w_p

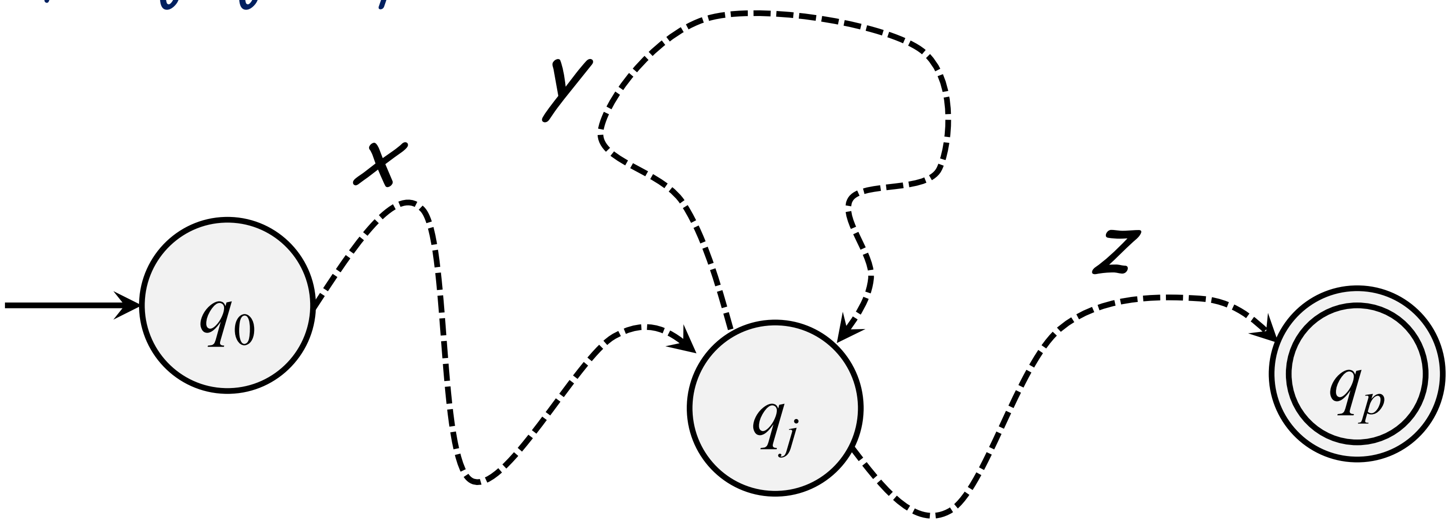
↑ ↑ ↑ ↑ ↑ ↑ ↑ ↑ ↑ ↑

q_0 q_1 q_2 q_3 q_4 q_5 q_6 q_7 q_{p-1} q_p



Machine Loops

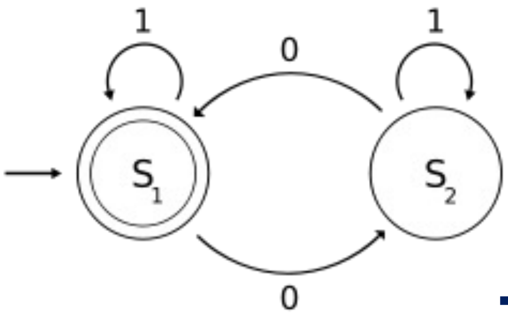
Let q_j be the first repeated state, that is, $q_j = q_{j+k}$ for some k, $0 \le j < j+k \le p$.



Where $w = w_1 w_2 \cdots w_p = xyz$,

$x = w_1 w_2 \cdots w_j$ $\quad y = w_{j+1} \cdots w_{j+k}$ $\quad z = w_{j+k+1} \cdots w_p$

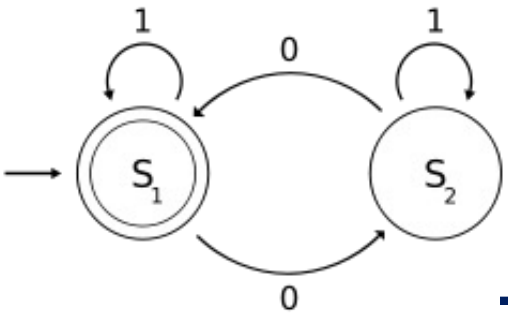
We conclude that $xy^i z \in L$ for all $i \ge 0$.



The Pumping Lemma

Theorem. If A is a regular language, then there is a number p (the pumping length) where if w is a string of length at least p, then $w = xyz$, such that

1. for each $i \geq 0$, $xy^iz \in A$,
2. $|y| > 0$, and
3. $|xy| \leq p$.



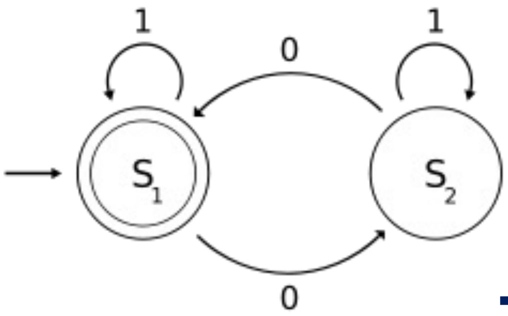
Deciding Regularity

Is $L = \{ 0^n 1^n : n \geq 0 \}$ regular? If so, then there are strings x, y, and z such that $xy^iz \in L$ for all $i \geq 0$. What does y look like?

String y consists entirely of 0s?

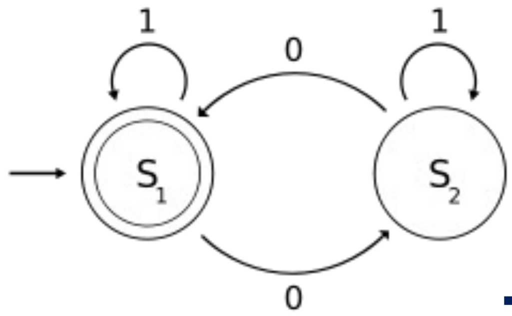
String y consists entirely of 1s?

String y consists of both 0s and 1s?



Combining Results

Is C = { w | w has an equal number of 0s and 1s } a regular language?



Picking the Right Substring to Pump

Is *PAL* = { $w \in \{0, 1\}^*$: w is a palindrome } a regular language?