CS 240 SI Worksheet Solutions Valerie Zhao Session #3 2/2/17

## **Integer Representation**

1. Using <u>8-bits</u> (which is 1 <u>byte</u> [fill in the blank]), what's **-25**<sub>10</sub> in:

## a. <u>Unsigned</u> integer representation? **Not possible**

- b. <u>Signed</u> integer representation? 10011001
- c. <u>Two's complement</u> representation? 11100111

## i. What's **25** in two's complement? **00011001**

2. <u>Without looking at your notes or any other materials</u>, fill in the following table for an <u>8-bit</u> binary integer:

Integer Representation	Minimum value (in base 10) Maximum value (in base	
Unsigned	0	255
Signed	-127	127
Two's Complement	-128	127

3. Why is <u>signed</u> integer representation flawed? (2 reasons)

- Normal addition involving negative integers doesn't produce the right results

- Two representations of 0 (+/- 0)
- a. How does two's complement remedy this?

- Addition involving negative integers is (usually) correct, i.e. if no

overflow

- Only one representation of 0

 Interpret the numbers given under "Integer in binary" according to the 3 different representations, then record the base-10 value it encodes: (for example, 0100 is 4 in all 3 encodings.)

Integer in binary	Unsigned	Signed	Two's Complement
1010	10	-2	-6
0111	7	7	7
1111	15	7	-1
0000	0	0	0
1000	8	0	-8

5. Calculate 0010 - 0111:

0010 - 0111 = 1011

Steps of the borrow algorithm: (like subtraction in base 10, but in binary)

- 0010 ("A")
- <u>-0111</u> ("B")

- In the rightmost column (least significant bit), 0 in "A" is smaller than 1 in "B", so 0 in "A" borrow from the bit to the left of it to become 10. 10 1 = 1.
- Moving leftward, in the 2nd-to-rightmost column, 1 in "A" became 0 because of the <u>borrowing from step 1</u>, which is smaller than 1 in "B", so it <u>borrows from the</u> <u>bit to the left of it</u> to become 10. 10 - 1 = 1.
- In the 3rd-to-rightmost column, 0 in "A" became 1 because of the borrowing from step 2 WHICH required this 0 to also borrow from the bit to the left of it (most significant bit). In other words, this 0 in "A" <u>borrowed from the most significant bit</u> to become 10 before step 2 borrowed from it and it became 1 (It's like subtraction in base 10, ex. 123 - 49.). 1 - 1 = 0.
- In the leftmost column, just like in step 3, 0 in "A" had to become 10 by <u>borrowing</u> from the bit to its left (imagine there was a 1 to its left--it'll maybe make sense in 4a) to become 10, then after step 3's borrowing it became 1. 1 0 = 1.
  - a. Essentially we're treating the equation as **10010 00111**. Try doing the reverse, i.e. calculating **1011 + 0111** → what do you get if you keep all bits of the result?

**Steps of the two's complement algorithm:** (even if the terms are in signed representation -- *why?*)

- 1. -0111 = +(-0111) → -x = -x+1 → 1001
- 2. 0010 + 1001 = **1011**

a. What's the answer (in base 10) if this expression was in <u>signed</u> integer representation?

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1011<sub>2</sub> = -3
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- b. In <u>two's complement</u>? **1011**<sub>2</sub> = -5
- c. How did overflow apply to what you did in parts a and b?

If you calculate 0111 + 1011 (this is the reverse of the given subtraction 0010 - 0111 = 1011 with the answer known), the complete result has 5 bits instead of 4. In signed integer representation, this indicates overflow. In two's complement, overflow was not an issue because despite the extra bit, the carry-in was equal to the carry-out, and it is not the case that the two terms being added have the same sign bit while the sum has the opposite sign bit.