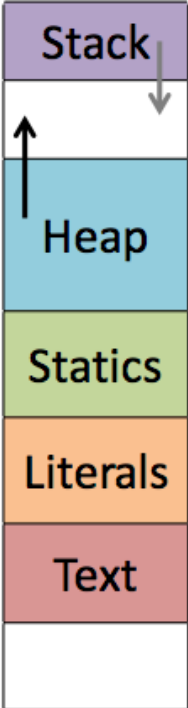


# CS240 Laboratory 6

## Disassembly and Reverse Engineering

### Memory Layout

	Addr		Perm	Contents	Managed by	Initialized
<b>Kernel</b> above 0x7fffffff	$2^N-1$ 	<b>Stack</b>	RW	Procedure context	Compiler	Run-time
<b>Stack</b> below 0x7fffffff grows down						
<b>Heap</b> above Data segment		<b>Heap</b>	RW	Dynamic data structures	Programmer, malloc/free, new/GC	Run-time
<b>Data segment</b> statics and literals		<b>Statics</b>	RW	Global variables/ static data structures	Compiler/ Assembler/Linker	Startup
		<b>Literals</b>	R	String literals	Compiler/ Assembler/Linker	Startup
<b>Text segment</b> starts at 0x400000	<b>Text</b>	X	Instructions	Compiler/ Assembler/Linker	Startup	
	0					

# Instructions

## Moving Data

*mov Src, Dest*

Note: the size of the data being referenced is often specified with an additional character:

*b* (byte)

*w* (2 bytes)

*l* (4 bytes), or

*q* (8 bytes).

## Arithmetic/Logical operations – 2 operands

*add Src, Dest*

*sub Src, Dest*

*imul Src, Dest*

*shr Src, Dest*

*sar Src, Dest*

*shl Src, Dest*

*sal Src, Dest*

*shr Src, Dest*

*xor Src, Dest*

*and Src, Dest*

*or Src, Dest*



**Memory**      -0x18(%rsp)

**Most General Form:**

$$D(Rb, Ri, S) \qquad \text{Mem}[\text{Reg}[Rb] + S * \text{Reg}[Ri] + D]$$

D:      Constant "displacement" value represented in 1, 2, or 4 bytes

Rb:      Base register: Any register

Ri:      Index register: Any except %esp (or %rsp if 64-bit); %ebp unlikely

S:      Scale: 1, 2, 4, or 8 (*why these numbers?*)

**Special Cases:** can use any combination of D, Rb, Ri and S

$$(Rb, Ri) \qquad \text{Mem}[\text{Reg}[Rb] + \text{Reg}[Ri]] \qquad (S=1, D=0)$$

$$D(Rb, Ri) \qquad \text{Mem}[\text{Reg}[Rb] + \text{Reg}[Ri] + D] \qquad (S=1)$$

$$(Rb, Ri, S) \qquad \text{Mem}[\text{Reg}[Rb] + S * \text{Reg}[Ri]] \qquad (D=0)$$

## Control Flow

Conditional jump instructions in X86 implement the following high-level constructs:

- if (condition) then {...} else {...}
- while (condition) {...}
- do {...} while (condition)
- for (initialization; condition; iterative) {...}

Unconditional jumps are used for high-level constructs such as:

- break
- continue

<b>jX</b>	<b>Condition</b>	<b>Description</b>
<b>jmp</b>	<b>1</b>	Unconditional
<b>je</b>	<b>ZF</b>	Equal / Zero
<b>jne</b>	<b>~ZF</b>	Not Equal / Not Zero
<b>ja</b>	<b>SF</b>	Negative
<b>jns</b>	<b>~SF</b>	Nonnegative
<b>jg</b>	<b>~ (SF^OF) &amp; ~ZF</b>	Greater (Signed)
<b>jge</b>	<b>~ (SF^OF)</b>	Greater or Equal (Signed)
<b>jl</b>	<b>(SF^OF)</b>	Less (Signed)
<b>jle</b>	<b>(SF^OF)   ZF</b>	Less or Equal (Signed)
<b>ja</b>	<b>~CF &amp; ~ZF</b>	Above (unsigned)
<b>jb</b>	<b>CF</b>	Below (unsigned)

Jump instructions encode the offset from next instruction to destination PC, instead of the absolute address of the destination (makes it easier to relocate the code)

# Turning C into Machine Code

## C Code

```
void sumstore(long x, long y,  
              long *dest) {  
    long t = x + y;  
    *dest = t;  
}
```

sum.c

**compiler (CS 301)**

gcc -Og -S sum.c

## Generated x86 Assembly Code

Human-readable language close to machine code.

```
sum:  
    addq %rdi,%rsi  
    movq %rsi,(%rdx)  
    retq
```

sum.s

**assembler**

## Object Code

```
01010101100010011110010110  
00101101000101000011000000  
00110100010100001000100010  
01111011000101110111000011
```

sum.o

**linker**

**Executable: sum**

Resolve references between object files,  
libraries, (re)locate data

- X86 instructions can be in different order from C code
- Some C expressions require multiple X86 instructions
- Some X86 instructions can cover multiple C expressions
- Compiler optimization can do some surprising things!
- Local or temporary variables can be stored in registers or on the stack

## Function Calling Conventions

- Arguments for functions are stored in registers, in the following order: arg1 – arg6: `%rdi,%rsi,%rdx,%rcx,%r8,%r9`
- If there are more than 6 parameters for a function, the rest of the arguments are stored on the stack before the function is called
- Return value from function is always in `%rax`

The compiler will use only part of a register if the value stored there will fit in less than 64 bits (8 bytes). This is an optimization that makes instructions a bit shorter.

So, in the code, you may see register names of the following form, all of which refer to `%rax`:

`%rax` = 8 byte value

`%eax` = 4 byte value

`%ax` = 2 byte value

`%al` = 1 byte value

# Tools

Tools can be used to examine bytes of object code (executable program) and reconstruct (reverse engineer) the assembly source.

**gdb** – disassembles an executable file into the associated assembly language representation, and provides tools for memory and register examination, single step execution, breakpoints, etc.

Object	Disassembled by GDB
0x00400536:	0x0000000000400536 <+0>: add %rdi,%rsi
0x48	0x0000000000400539 <+3>: mov %rsi,(%rdx)
0x01	0x000000000040053c <+6>: retq
0xfe	
0x48	\$ gdb sum
0x89	(gdb) disassemble sumstore
0x32	(disassemble function)
0xc3	(gdb) x/7b sum
	(examine the 13 bytes starting at sum)

## objdump

can also be used to disassemble and display information

```
$ objdump -t p
```

Prints out the program's symbol table. The symbol table includes the names of all functions and global variables, the names of all the functions the called, and their addresses.



*\$ objdump -d p*

## **Object Code**

0x401040 <sum>:

0x55

0x89

0xe5

0x8b

0x45

0x0c

0x03

0x45

0x08

0x89

0xec

0x5d

0xc3

## **Disassembled version**

00401040 <\_sum>:

0: 55        push %ebp

1: 89 e5     mov %esp,%ebp

3: 8b 45 0c  mov 0xc(%ebp),%eax

6: 03 45 08  add 0x8(%ebp),%eax

9: 89 ec     mov %ebp,%esp

b: 5d pop    %ebp

c: c3        ret

## **strings**

*\$ strings -t x p*

Displays the printable strings in your program.