

## positional number representation



Base determines:
Maximum digit (base -1 ). Minimum digit is 0 .
Weight of each position.
Each position holds a digit.
Represented value = sum of all position values
Position value $=$ digit value $\times$ base $^{\text {position }}$

Powers of 2:
memorize up to $\geq \mathbf{2 ~}^{\mathbf{1 0}}$ (in base ten)
conversion and arithmetic

$$
19_{10}=?_{2} \quad 1001_{2}=?_{10}
$$

$240_{10}=?_{2}$
$11010011_{2}=?_{10}$
$101_{2}+1011_{2}=?_{2}$
$1001011_{2} \times 2_{10}=?_{2}$

## byte $=\mathbf{8}$ bits

Smallest unit of data
used by a typical modern computer

Binary $00000000_{2}$-- $11111111_{2}$
Decimal $\quad 000_{10}--255_{10}$
Hexadecimal $\quad 00_{16}--\mathrm{FF}_{16}$
Byte $=2$ hex digits!
Programmer's hex notation (C, etc.):
$0 \times B 4=B 4{ }_{16}$
Octal (base 8) also useful.

| $x^{e^{t}} \rho^{e^{c i n}} \operatorname{bin}^{n^{2}} n^{r r t}$ |  |  |
| :---: | :---: | :---: |
| 0 | 0 | 0000 |
| 1 | 1 | 0001 |
| 2 | 2 | 0010 |
| 3 | 3 | 0011 |
| 4 | 4 | 0100 |
| 5 | 5 | 0101 |
| 6 | 6 | 0110 |
| 7 | 7 | 0111 |
| 8 | 8 | 1000 |
| 9 | 9 | 1001 |
| A | 10 | 1010 |
| B | 11 | 1011 |
| C | 12 | 1100 |
| D | 13 | 1101 |
| E | 14 | 1110 |
| F | 15 | 1111 |

## char: representing characters

A C-style string is represented by a series of bytes (chars).

- One-byte ASCII codes for each character.
- ASCII = American Standard Code for Information Interchange

| 32 | space | 48 | 0 |  | @ | 80 | P | 96 |  | 112 | p |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 33 | ! | 49 | 1 | 65 | A | 81 | Q | 97 | a | 113 | q |
| 34 | " | 50 | 2 | 66 | B | 82 | R | 98 | b | 114 | r |
| 35 | \# | 51 | 3 | 67 | C | 83 | S | 99 | c | 115 | s |
| 36 | \$ | 52 | 4 | 68 | D | 84 | T | 100 | d | 116 | t |
| 37 | \% | 53 | 5 | 69 | E | 85 | U | 101 | e | 117 | u |
| 38 | \& | 54 | 6 | 70 | F | 86 | $v$ | 102 | f | 118 | $v$ |
| 39 |  | 55 | 7 | 71 | G | 87 | w | 103 | g | 119 | w |
| 40 | 1 | 56 | 8 | 72 | H | 88 | X | 104 | h | 120 | x |
| 41 | ) | 57 | 9 | 73 | , | 89 | Y | 105 | 1 | 121 | $y$ |
| 42 | * | 58 | : | 74 | J | 90 | z | 106 | j | 122 | z |
| 43 | + | 59 | ; | 75 | K | 91 | [ | 107 | , | 123 | \{ |
| 44 | , | 60 | < | 76 | L | 92 | $\backslash$ | 108 | 1 | 124 | । |
| 45 | - | 61 | $=$ | 77 | M | 93 | ] | 109 | m | 125 | \} |
| 46 |  | 62 | $>$ | 78 | N | 94 | $\wedge$ | 110 | n | 126 | $\sim$ |
| 47 | 1 | 63 | ? | 79 | 0 | 95 |  | 111 | 0 | 127 | del |



## bitwise operators

Bitwise operators on fixed-width bit vectors.
AND \& OR | XOR ^ NOT ~

| 01101001 | 01101001 | 01101001 |  |
| :---: | :---: | :---: | :---: |
| \& 01010101 | \| 01010101 | $\wedge 01010101$ | $\sim 01010101$ |

01010101
Laws of Boolean algebra apply bitwise.
$\wedge 01010101$
e.g., DeMorgan's Law: ~(A | B) = ~A \& ~B
fixed-size data representations

| Java Data Type | C Data Type | (size in bytes) |  |
| :---: | :---: | :---: | :---: |
|  |  | [word = 32 | = 64 bits] |
| boolean |  | 1 | 1 |
| byte | char | 1 | 1 |
| char |  | 2 | 2 |
| short | short int | 2 | 2 |
| int | int | 4 | 4 |
| float | float | 4 | 4 |
|  | long int | 4 | 8 |
| double <br> long | double | 8 | 8 |
|  | long long | 8 | 8 |
|  | long double | 8 | 16 |
| Depends on word size! |  |  |  |

## Aside: sets as bit vectors

Representation: $n$-bit vector gives subset of $\{0, \ldots, n-1\}$.

$$
\begin{aligned}
\mathrm{a}_{i}=1 \equiv i \in A & \\
& 01101001 \\
76543210 & \{0,3,5,6\} \\
& \\
01010101 & \{0,2,4,6\} \\
76543210 &
\end{aligned}
$$

Set Operations?
bitwise operators in C

```
\& | ^ ~ apply to any integral data type
long, int, short, char, unsigned
```

Examples (char)
~0×41 =
$\sim 0 \times 00=$
$0 \times 69 \& 0 \times 55=$
$0 \times 69$ | $0 \times 55=$
Many bit-twiddling puzzles in upcoming assignment

## Encode playing cards.

52 cards in 4 suits
How do we encode suits, face cards?
What operations should be easy to implement?
Get and compare rank
Get and compare suit


## logical operations in C

```
&& || ! apply to any "integral" data type
                long, int, short, char, unsigned
    0 is false
                                    nonzero is true
                                    result always 0 or 1
    early termination a.k.a. short-circuit evaluation
Examples (char)
    ! 0\times41 =
    !0x00 =
    !!0\times41 =
0x69 && 0\times55 =
0x69 || 0x55 =
```


## Two possible representations

52 cards - 52 bits with bit corresponding to card set to 1

"One-hot" encoding
Hard to compare values and suits independently Not space efficient

4 bits for suit, 13 bits for card value - 17 bits with two set to


## Pair of one-hot encoded values

Easier to compare suits and values independently
Smaller, but still not space efficient

## Two better representations

Binary encoding of all 52 cards - only 6 bits needed

Number cards uniquely from 0
low-order 6 bits of a byte
Smaller than one-hot encodings.
Hard to compare value and suit

Binary encoding of suit (2 bits) and value (4 bits) separately

Number each suit uniquely
Number each value uniquely


Still small
Easy suit, value comparisons

## Compare Card Values

mask: a bit vector that, when bitwise ANDed with another bit vector $v$, turns all but the bits of interest in $v$ to 0
\#define VALUE_MASK


```
int greaterValue(char card1, char card2) {
```

\}
char hand[5]; // represents a 5-card hand
char card1, card2; // two cards to compare
if ( greaterValue (hand[0], hand[1]) ) \{ ... \}

## Compare Card Suits

mask: a bit vector that, when bitwise ANDed with another bit vector $v$, turns all but the bits of interest in $v$ to 0

```
#define SUIT_MASK 0x30
```



```
int sameSuit(char card1, char card2) {
    return !((card1 & SUIT_MASK) ^ (card2 & SUIT_MASK));
    //same as (card1 & SUIT_MASK) == (card2 & SUIT_MASK);
}
```

```
char hand[5]; // represents a 5-card hand
char card1, card2; // two cards to compare
if ( sameSuit(hand[0], hand[1]) ) { ... }
```


## Bit shifting

```
        x 10011001
    x<<2 1001100100 logical shift left 2
    lose bits on left
        fill with zeroes on right
```

            \(10011001 \quad x\)
        fill with zeroes on left
    logical shift right 2
            0010011001
                                    lose bits on right
    arithmetic shift right 2
$1110011001 x \gg 2$
fill with copies of MSB on left

## Shift gotchas

Logical or arithmetic shift right: how do we tell?
C: compiler chooses
Usually based on type: rain check!
Java: >> is arithmetic, >>> is logical

Shift an n-bit type by at least 0 and no more than n-1.
C : other shift distances are undefined.
anything could happen
Java: shift distance is used modulo number of bits in shifted type Given int $\mathrm{x}: \mathrm{x} \ll 34=\mathrm{x} \ll 2$

## Shift and mask: extract a bit field

## Write a C function that

extracts the $2^{\text {nd }}$ most significant byte
from its 32-bit integer argument.

## Example behavior:

argument: 0b 01100001011000100110001101100100 expected result: 0b 00000000000000000000000001100010

All other bits are zero. Desired bits in least significant byte.
int get2ndMSB(int $x$ ) \{

