



x86: Procedures and the Call Stack

The call stack discipline
x86 procedure call and return instructions
x86 calling conventions
x86 register-saving conventions

Why procedures?

Why functions? Why methods?

```
int contains_char(char* haystack, char needle) {
    while (*haystack != '\0') {
        if (*haystack == needle) return 1;
        haystack++;
    }
    return 0;
}
```

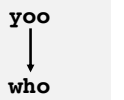
Procedural Abstraction

Implementing procedures

1. How does a caller pass arguments to a procedure? ✓
2. How does a caller receive a return value from a procedure? ✓
3. Where does a procedure store local variables? ✓?
4. How does a procedure know where to return (what code to execute next when done)? ??
5. How do procedures share limited registers and memory? ??

Procedure call/return: Jump?

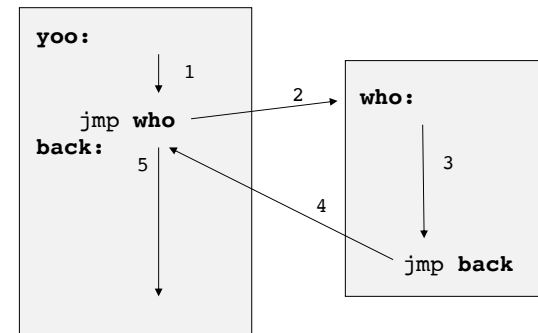
Call Chain



```
yoo(...) {
    . . .
    who();
    . . .
}
```

```
who(...) {
    . . .
    . . .
}
```

```
ru(...) {
    . . .
}
```



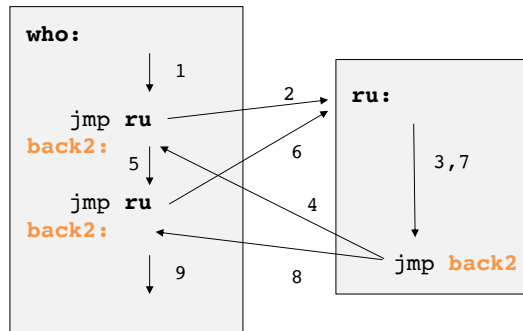
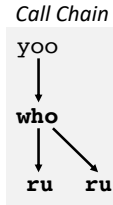
But what if we want to call a function from multiple places in the code?

Procedure call/return: Jump? **Broken!**

```
yoo(...) {
  . . .
  who();
  . . .
}
```

```
who(...) {
  . . .
  ru();
  . . .
  ru();
  . . .
}
```

```
ru(...) {
  . . .
}
```



But what if we want to call a function from multiple places in the code?
Broken: needs to track context.

Implementing procedures

requires **separate storage *per call!***
 (not just per procedure)

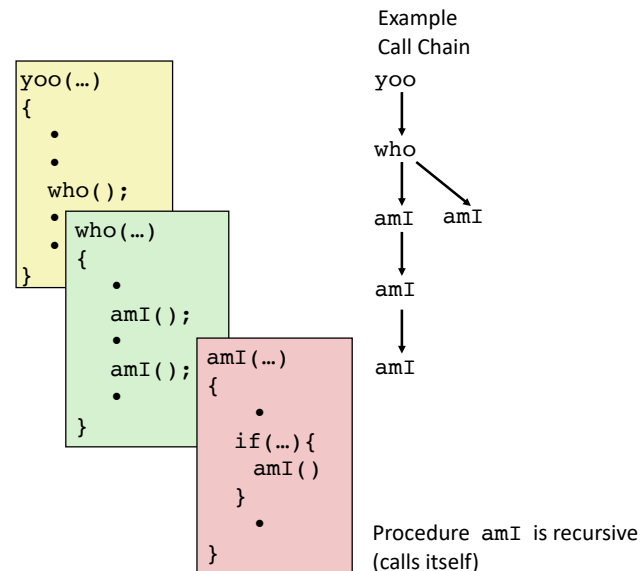
1. How does a caller pass arguments to a procedure? ✓
2. How does a caller receive a return value from a procedure? ✓
3. Where does a procedure store local variables? ✓?
4. How does a procedure know where to return (what code to execute next when done)? ??
5. How do procedures share limited registers and memory? ??

Memory Layout

reminder

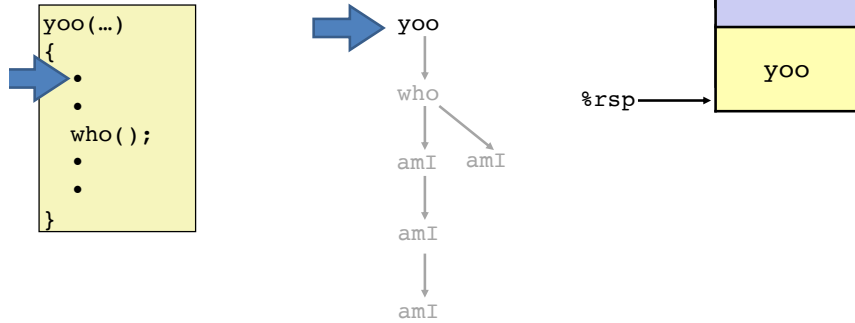
Addr	Perm	Contents	Managed by	Initialized
2^N-1 ↑	RW	Procedure context	Compiler	Run-time
Stack ↓				
↑	RW	Dynamic data structures	Programmer, malloc/free, new/GC	Run-time
Heap				
↓	RW	Global variables/ static data structures	Compiler/ Assembler/Linker	Startup
Statics				
↑	R	String literals	Compiler/ Assembler/Linker	Startup
Literals				
↓	X	Instructions	Compiler/ Assembler/Linker	Startup
Text				
0				

Call stack tracks context

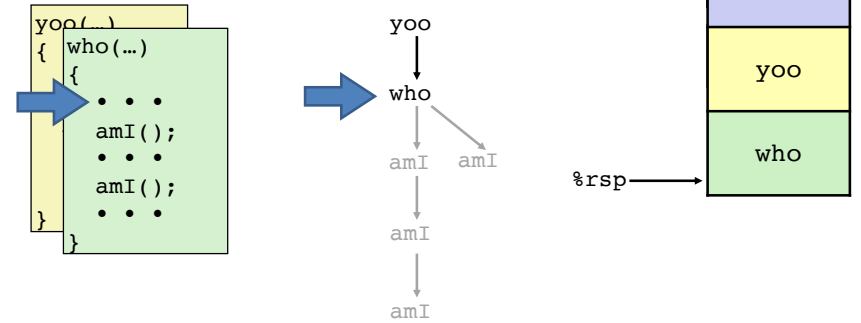


Procedure amI is recursive (calls itself)

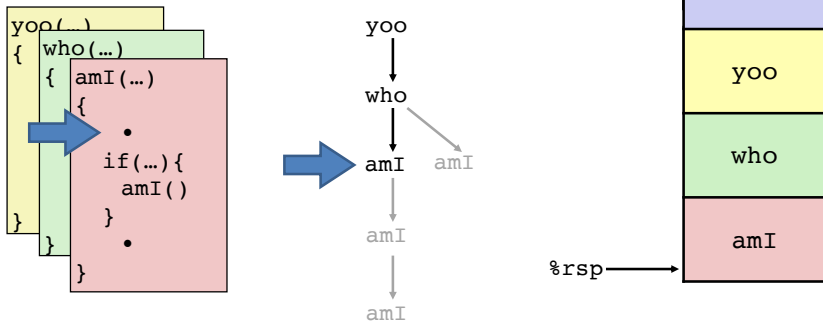
Call stack tracks context



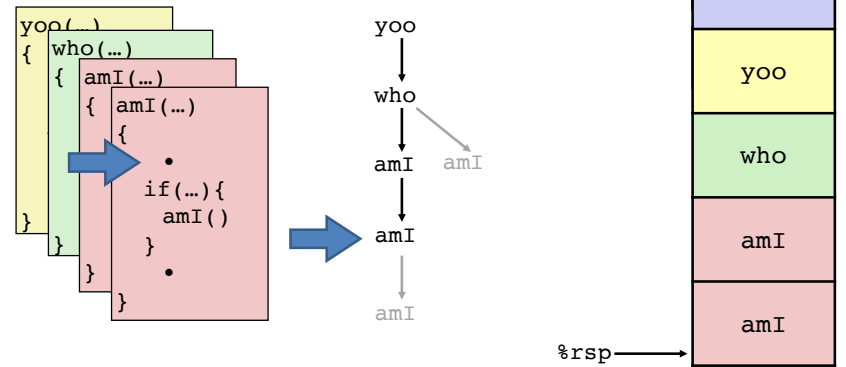
Call stack tracks context



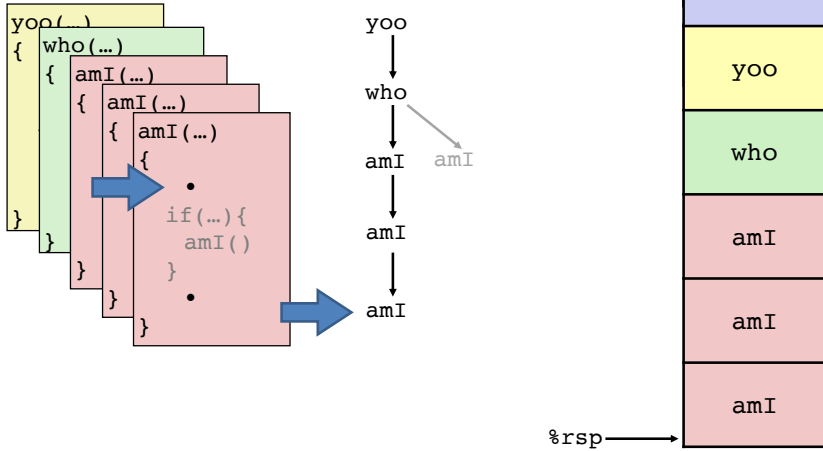
Call stack tracks context



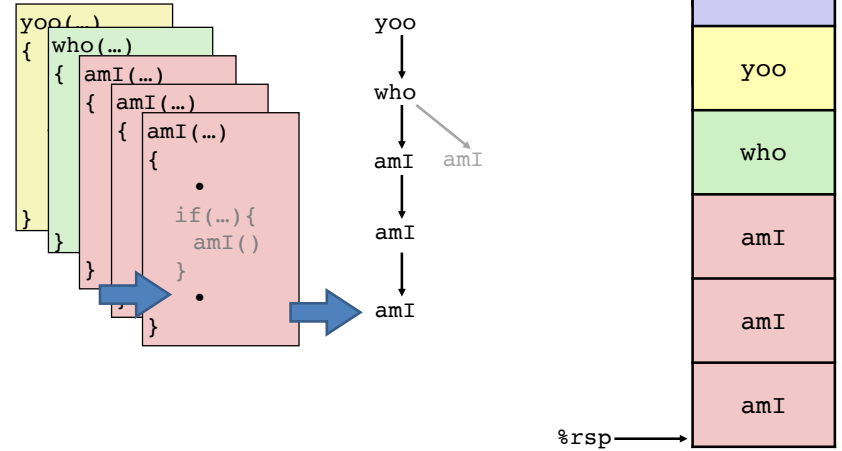
Call stack tracks context



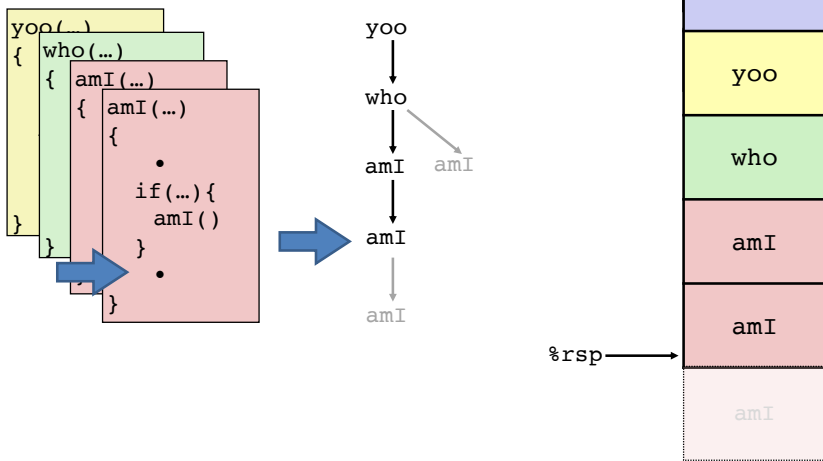
Call stack tracks context



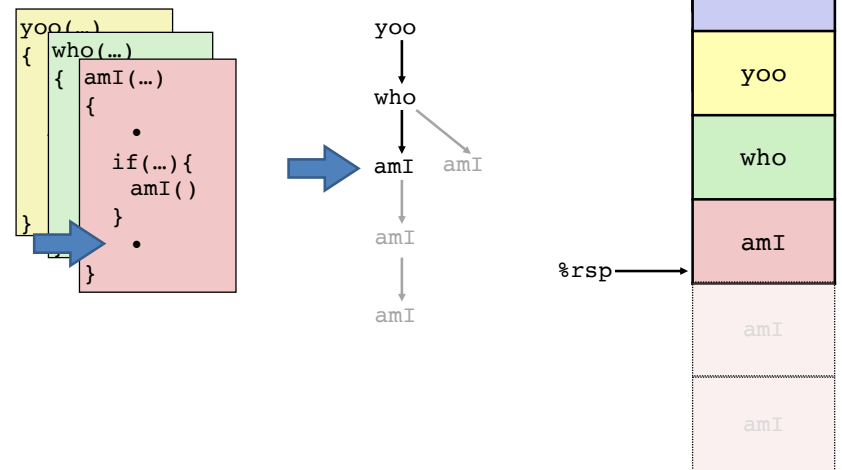
Call stack tracks context



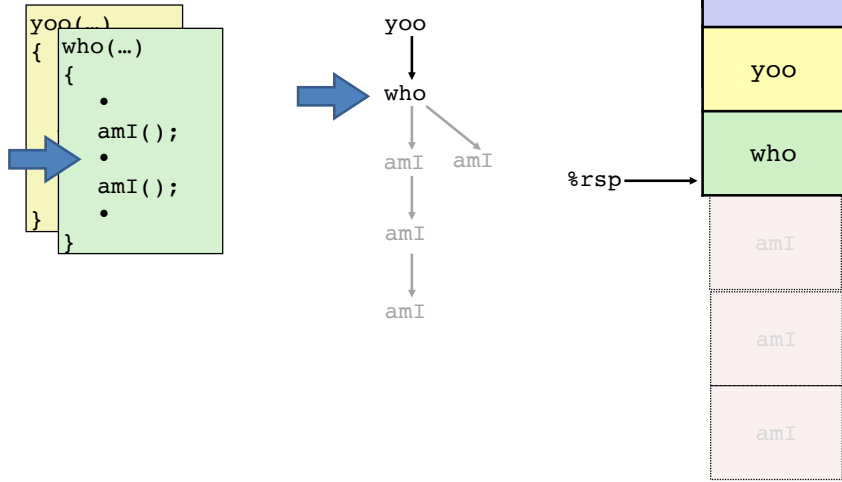
Call stack tracks context



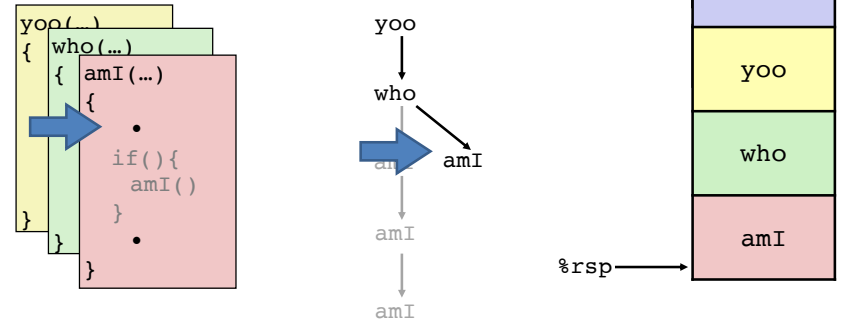
Call stack tracks context



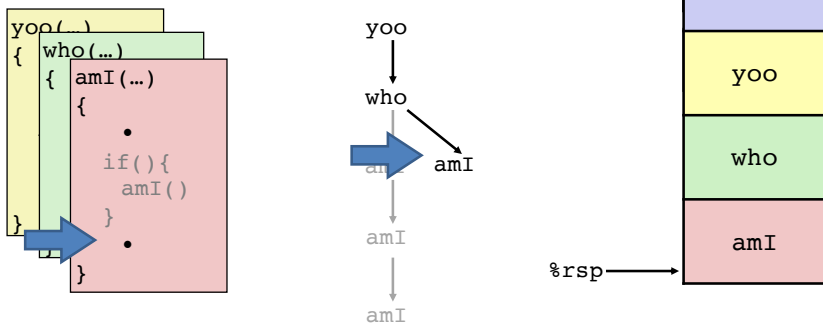
Call stack tracks context



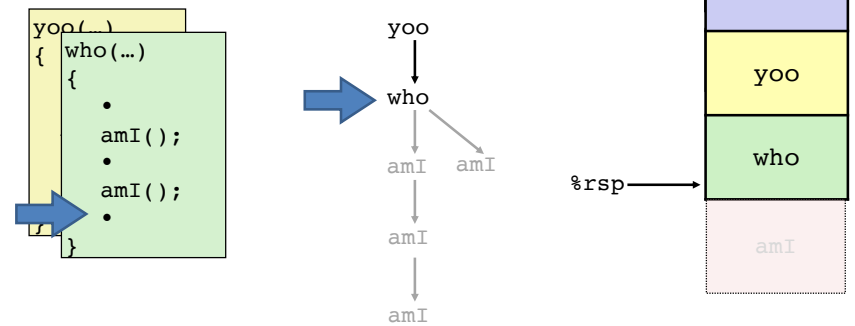
Call stack tracks context



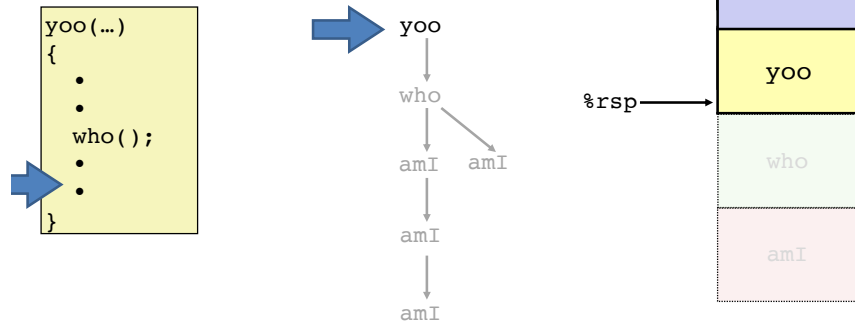
Call stack tracks context



Call stack tracks context



Call stack tracks context

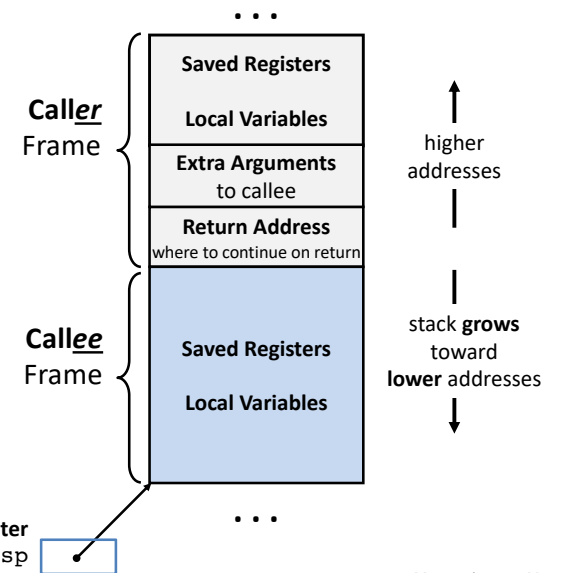


The call stack supports procedures.

Stack frame: section of stack used by one procedure *call* to store context while running.

Procedure code manages stack frames explicitly.

- **Setup:** allocate space at start of procedure.
- **Cleanup:** deallocate space before return.



Procedure control flow instructions

Procedure call: `callq target`

1. Push return address on stack
2. Jump to *target*

Return address: Address of instruction after call.

```
400544: callq 400550 <mult2>
400549: movq %rax, (%rbx)
```

Procedure return: `retq`

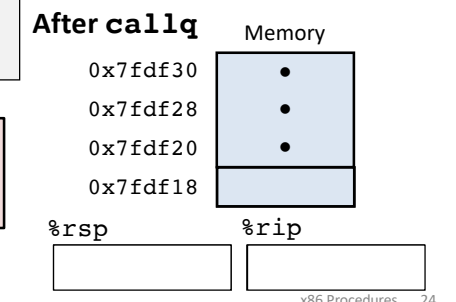
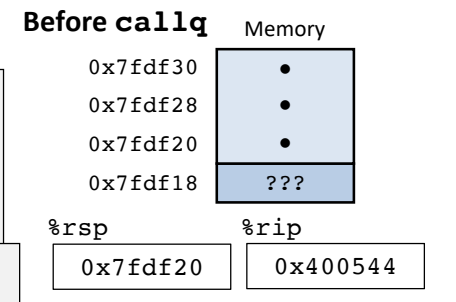
1. Pop return address from stack
2. Jump to return address

Call example

```
000000000400540 <multstore>:
.
.
400544: callq 400550 <mult2>
400549: mov  %rax, (%rbx)
.
.
000000000400550 <mult2>:
400550: mov  %rdi,%rax
.
.
400557: retq
```

callq target

1. Push return address on stack
2. Jump to *target*



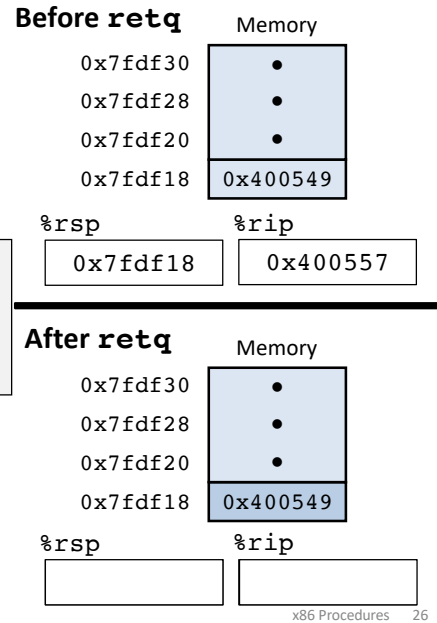
Return example

```
0000000000400540 <multstore>:
.
.
400544: callq 400550 <mult2>
400549: mov  %rax, (%rbx)
.
```

```
0000000000400550 <mult2>:
400550: mov  %rdi,%rax
.
.
400557: retq
```

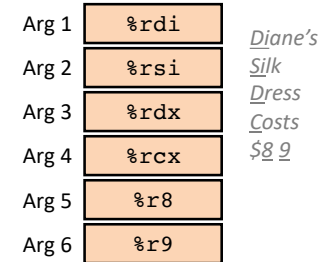
retq

1. Pop return address from stack
2. Jump to return address

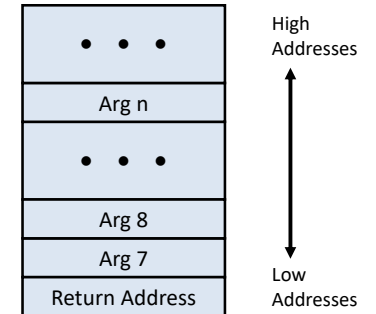


Procedure data flow conventions

First 6 arguments:
passed in **registers**



Remaining arguments:
passed on **stack** (in memory)



Return value:
passed in **%rax**

%rax

Allocate stack space for arguments only when needed.

callq puzzle

optional

```
callq next
next:
popq %rax
```

What gets stored into %rax?

Why is there no ret instruction corresponding to the call?

What does this code do? (Hint: unusual use of call.)

Procedure data flow puzzle

ex

C function body:

```
_____ huh(_____ , _____ , _____ , _____ ) {
    *p = d;
    return x - c;
}
```

Translated to x86 assembly:

```
huh:
    movsbl %dl, %edx
    movl %edx, (%rsi)
    movswl %di, %edi
    subl %edi, %ecx
    movl %ecx, %eax
    retq
```

Reverse engineer the x86 huh procedure and the body of the C huh function to fill blanks in the C huh function header with:

- the parameter types / order; and
- the return type.

movsbl = **move** sign-extending a **byte** to a long (4-byte)

movswl = **move** sign-extending a **word** (2-byte) to a long (4-byte)

Procedure data flow puzzle



C function body:

```
int huh(short c, int* p, char d, int x) {
    *p = d;
    return x - c;
}
```

Translated to x86 assembly:

```
huh:
    movsbl %dl, %edx
    movl %edx, (%rsi)
    movswl %di, %edi
    subl %edi, %ecx
    movl %ecx, %eax
    retq
```

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movsbl = **move** sign-extending a **byte** to a long (4-byte)
 movswl = **move** sign-extending a **word** (2-byte) to a long (4-byte)

Procedure call / stack frame example

```
long step_up() {
    long v1 = 240;
    long v2 = increment(&v1, 61);
    return v1+v2;
}
```

```
step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq
```

Passes address of local variable (in stack).

Uses memory through pointer.

```
long increment(long* p, long val) {
    long x = *p;
    long y = x + val;
    *p = y;
    return x;
}
```

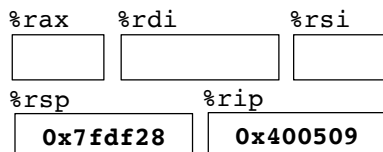
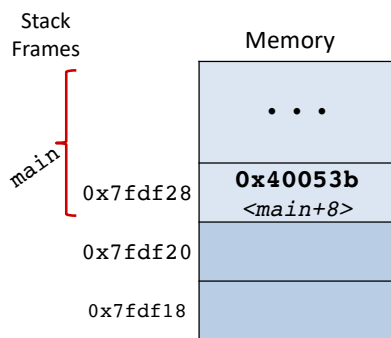
```
increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq
```

Procedure call example (step 0) main called step_up

```
long step_up() {
    long v1 = 240;
    long v2 = increment(&v1, 61);
    return v1+v2;
}
```

```
step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq
```

```
increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq
```

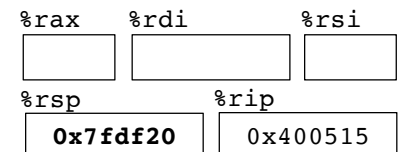
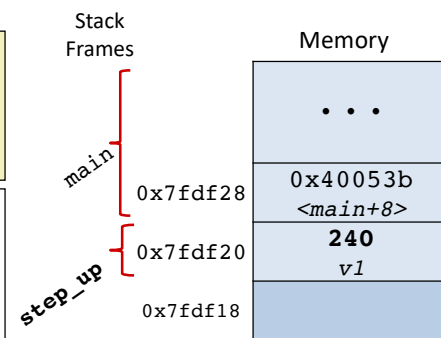


Procedure call example (step 1) Allocate space for local vars

```
long step_up() {
    long v1 = 240;
    long v2 = increment(&v1, 61);
    return v1+v2;
}
```

```
step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq
```

```
increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq
```

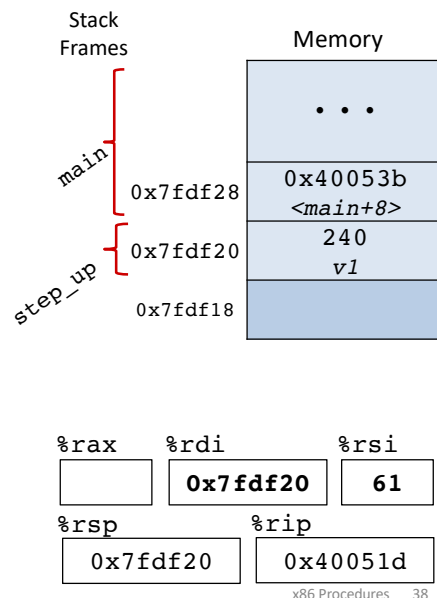


Procedure call example (step 2) Set up args for call to increment

```
long step_up() {
    long v1 = 240;
    long v2 = increment(&v1, 61);
    return v1+v2;
}
```

```
step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq
```

```
increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq
```

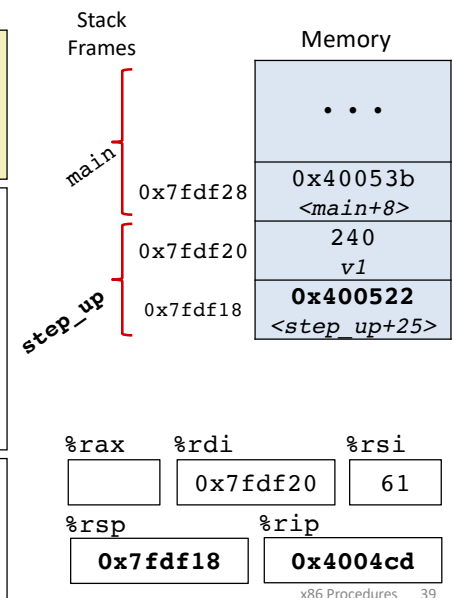


Procedure call example (step 3) Call increment

```
long step_up() {
    long v1 = 240;
    long v2 = increment(&v1, 61);
    return v1+v2;
}
```

```
step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq
```

```
increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq
```

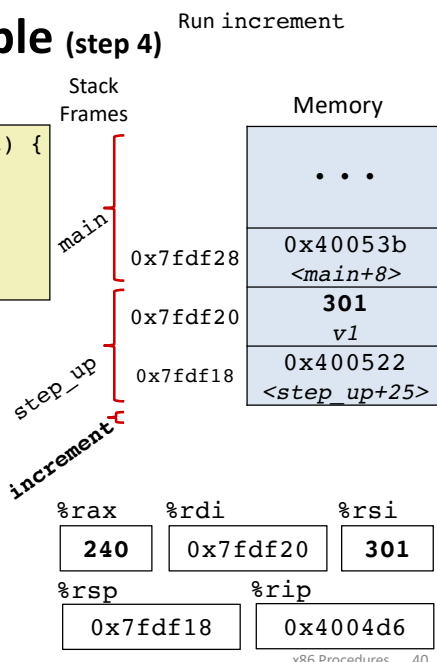


Procedure call example (step 4) Run increment

```
long step_up() {
    long increment(long* p, long val) {
        long x = *p;
        long y = x + val;
        *p = y;
        return x;
    }
}
```

```
step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq
```

```
increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq
```

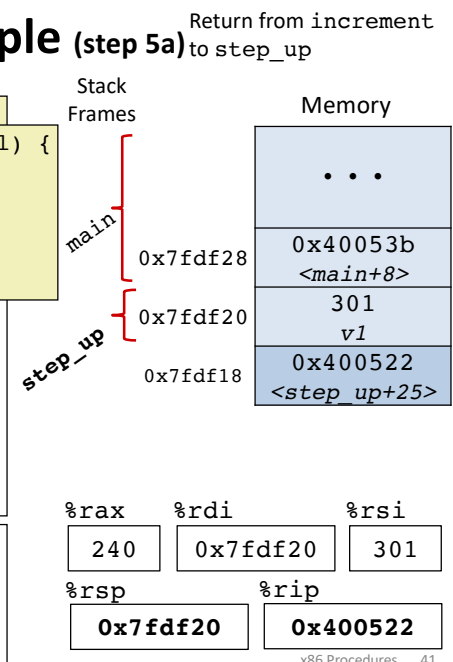


Procedure call example (step 5a) Return from increment to step_up

```
long step_up() {
    long increment(long* p, long val) {
        long x = *p;
        long y = x + val;
        *p = y;
        return x;
    }
}
```

```
step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq
```

```
increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq
```

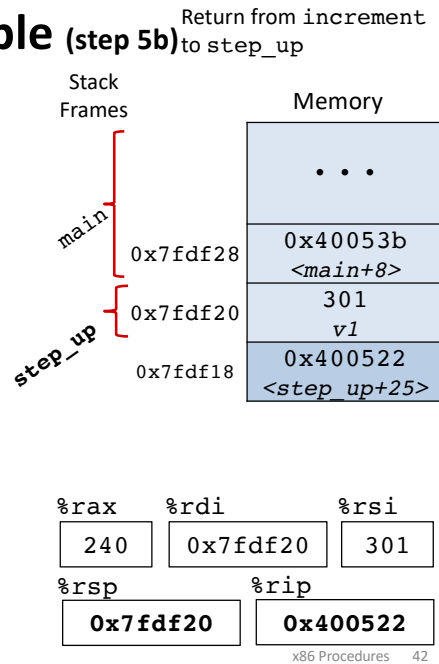


Procedure call example (step 5b) Return from increment to step_up

```
long step_up() {
    long v1 = 240;
    long v2 = increment(&v1, 61);
    return v1+v2;
}
```

```
step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq
```

```
increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq
```

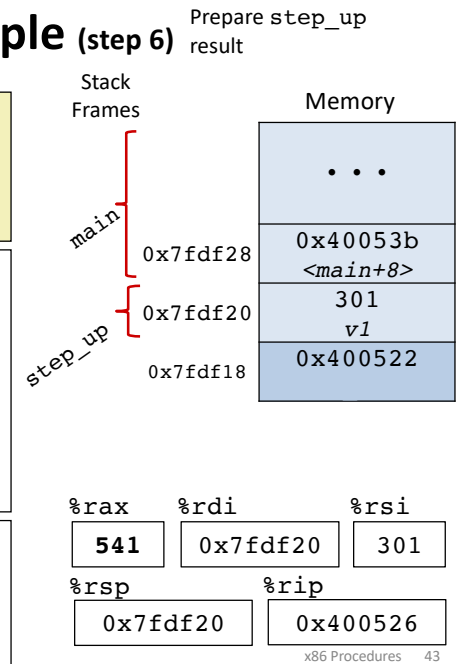


Procedure call example (step 6) Prepare step_up result

```
long step_up() {
    long v1 = 240;
    long v2 = increment(&v1, 61);
    return v1+v2;
}
```

```
step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq
```

```
increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq
```

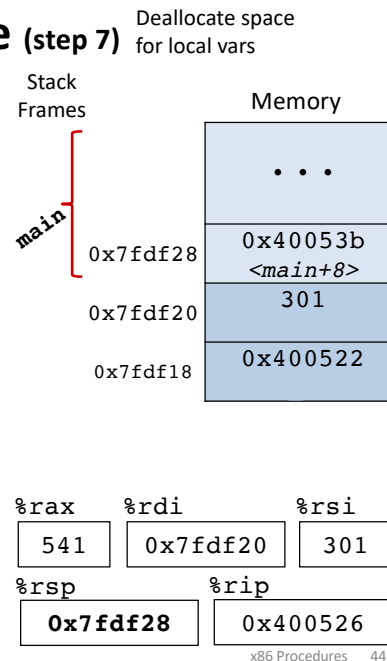


Procedure call example (step 7) Deallocate space for local vars

```
long step_up() {
    long v1 = 240;
    long v2 = increment(&v1, 61);
    return v1+v2;
}
```

```
step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq
```

```
increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq
```

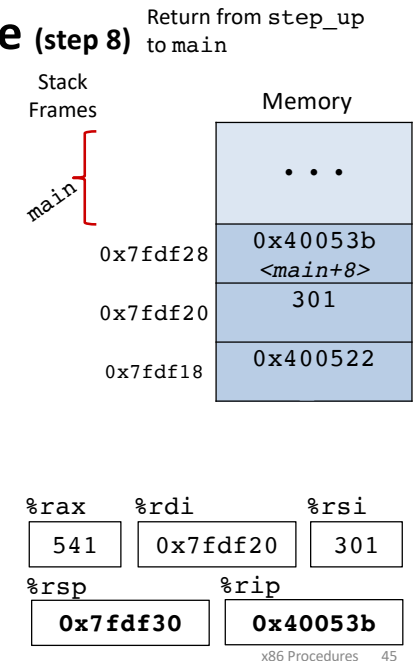


Procedure call example (step 8) Return from step_up to main

```
long step_up() {
    long v1 = 240;
    long v2 = increment(&v1, 61);
    return v1+v2;
}
```

```
step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq
```

```
increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq
```



Implementing procedures

1. How does a caller pass arguments to a procedure? ✓
2. How does a caller receive a return value from a procedure? ✓
3. Where does a procedure store local variables? ✓
4. How does a procedure know where to return (what code to execute next when done)? ✓
5. How do procedures share limited registers and memory? ??

Register saving conventions

yoo calls who:

Caller **Callee**

Will register contents still be there after a procedure call?

```
yoo:
. . .
movq $12345, %rbx
call who
addq %rbx, %rax
. . .
ret
```

```
who:
. . .
addq %rdi, %rbx
. . .
ret
```

Conventions:

Caller Save

Callee Save

x86-64 register conventions

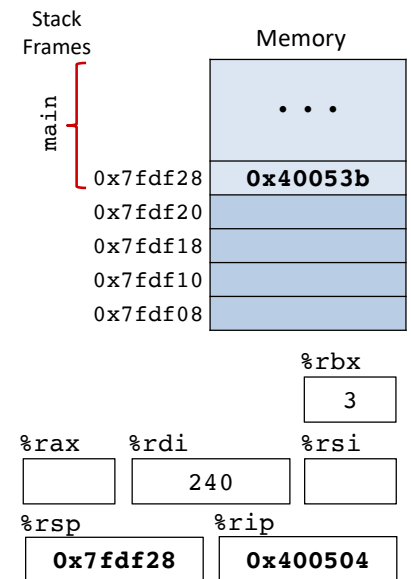
%rax	Return value – Caller saved	%r8	Argument #5 – Caller saved
%rbx	Callee saved	%r9	Argument #6 – Caller saved
%rcx	Argument #4 – Caller saved	%r10	Caller saved
%rdx	Argument #3 – Caller saved	%r11	Caller Saved
%rsi	Argument #2 – Caller saved	%r12	Callee saved
%rdi	Argument #1 – Caller saved	%r13	Callee saved
%rsp	Stack pointer	%r14	Callee saved
%rbp	Callee saved	%r15	Callee saved

Callee-save example (step 0)

main called step_by(240)

```
long step_by(long x) {
    long v1 = x;
    long v2 = increment(&v1, 61);
    return x + v2;
}
```

```
step_by:
400504: pushq %rbx
400506: movq %rdi, %rbx
400509: subq $16, %rsp
40050d: movq %rdi, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq %rbx, %rax
400525: addq $16, %rsp
400529: popq %rbx
40052b: retq
```

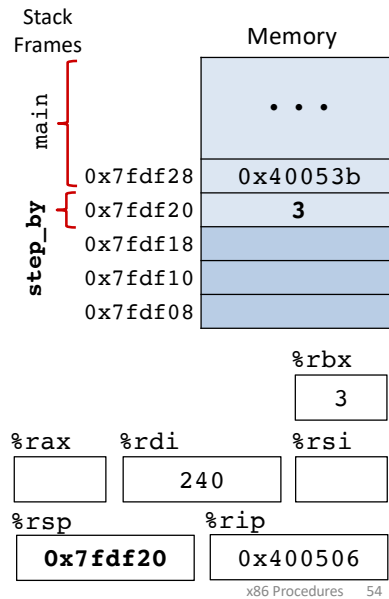


Callee-save example (step 1)

Save register %rbx

```
long step_by(long x) {
    long v1 = x;
    long v2 = increment(&v1, 61);
    return x + v2;
}
```

```
step_by:
400504: pushq %rbx
400506: movq  %rdi, %rbx
400509: subq  $16, %rsp
40050d: movq  %rdi, (%rsp)
400515: movq  %rsp, %rdi
400518: movl  $61, %esi
40051d: callq 4004cd <increment>
400522: addq  %rbx, %rax
400525: addq  $16, %rsp
400529: popq  %rbx
40052b: retq
```

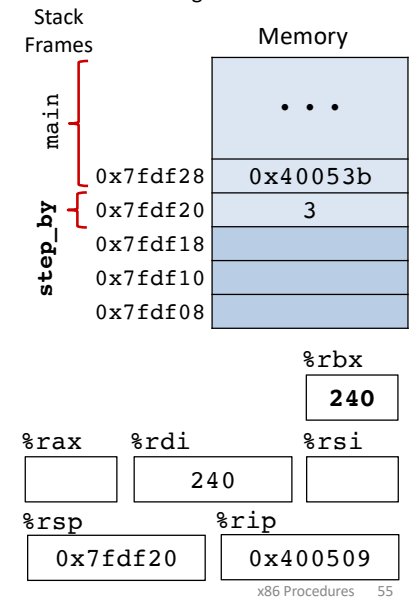


Callee-save example (step 2)

Copy argument x to %rbx for continued use after calling increment.

```
long step_by(long x) {
    long v1 = x;
    long v2 = increment(&v1, 61);
    return x + v2;
}
```

```
step_by:
400504: pushq %rbx
400506: movq %rdi, %rbx
400509: subq  $16, %rsp
40050d: movq  %rdi, (%rsp)
400515: movq  %rsp, %rdi
400518: movl  $61, %esi
40051d: callq 4004cd <increment>
400522: addq  %rbx, %rax
400525: addq  $16, %rsp
400529: popq  %rbx
40052b: retq
```

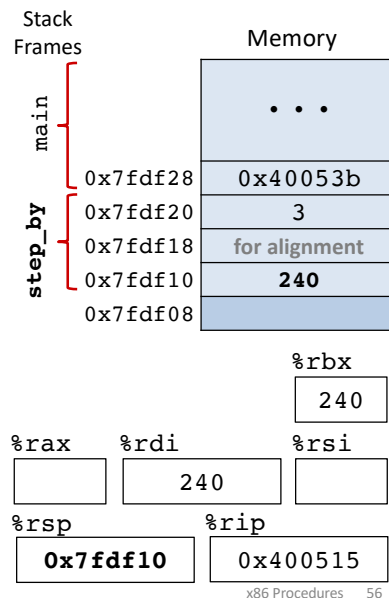


Callee-save example (step 3)

Set up stack frame Initialize v1

```
long step_by(long x) {
    long v1 = x;
    long v2 = increment(&v1, 61);
    return x + v2;
}
```

```
step_by:
400504: pushq %rbx
400506: movq  %rdi, %rbx
400509: subq $16, %rsp
40050d: movq %rdi, (%rsp)
400515: movq  %rsp, %rdi
400518: movl  $61, %esi
40051d: callq 4004cd <increment>
400522: addq  %rbx, %rax
400525: addq  $16, %rsp
400529: popq  %rbx
40052b: retq
```

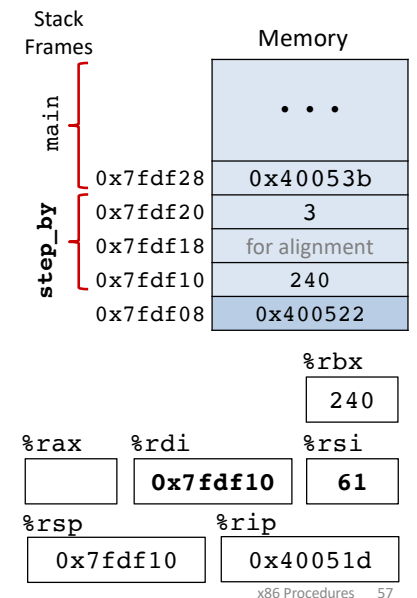


Callee-save example (step 4)

Set up arguments

```
long step_by(long x) {
    long v1 = x;
    long v2 = increment(&v1, 61);
    return x + v2;
}
```

```
step_by:
400504: pushq %rbx
400506: movq  %rdi, %rbx
400509: subq  $16, %rsp
40050d: movq  %rdi, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq  %rbx, %rax
400525: addq  $16, %rsp
400529: popq  %rbx
40052b: retq
```

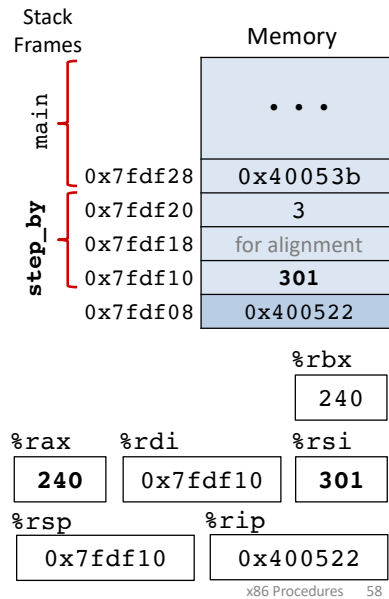


Callee-save example (step 5)

Call, execute, and return from increment

```
long step_by(long x) {
    long v1 = x;
    long v2 = increment(&v1, 61);
    return x + v2;
}
```

```
step_by:
400504: pushq %rbx
400506: movq  %rdi, %rbx
400509: subq  $16, %rsp
40050d: movq  %rdi, (%rsp)
400515: movq  %rsp, %rdi
400518: movl  $61, %esi
40051d: callq 4004cd <increment>
400522: addq  %rbx, %rax
400525: addq  $16, %rsp
400529: popq  %rbx
40052b: retq
```

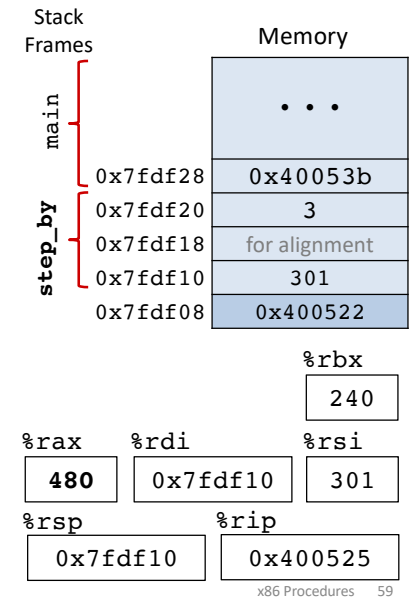


Callee-save example (step 6)

Prepare return value

```
long step_by(long x) {
    long v1 = x;
    long v2 = increment(&v1, 61);
    return x + v2;
}
```

```
step_by:
400504: pushq %rbx
400506: movq  %rdi, %rbx
400509: subq  $16, %rsp
40050d: movq  %rdi, (%rsp)
400515: movq  %rsp, %rdi
400518: movl  $61, %esi
40051d: callq 4004cd <increment>
400522: addq  %rbx, %rax
400525: addq  $16, %rsp
400529: popq  %rbx
40052b: retq
```

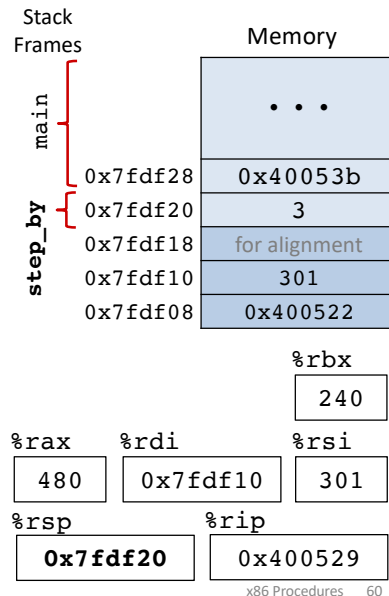


Callee-save example (step 7)

Clean up stack frame

```
long step_by(long x) {
    long v1 = x;
    long v2 = increment(&v1, 61);
    return x + v2;
}
```

```
step_by:
400504: pushq %rbx
400506: movq  %rdi, %rbx
400509: subq  $16, %rsp
40050d: movq  %rdi, (%rsp)
400515: movq  %rsp, %rdi
400518: movl  $61, %esi
40051d: callq 4004cd <increment>
400522: addq  %rbx, %rax
400525: addq  $16, %rsp
400529: popq  %rbx
40052b: retq
```

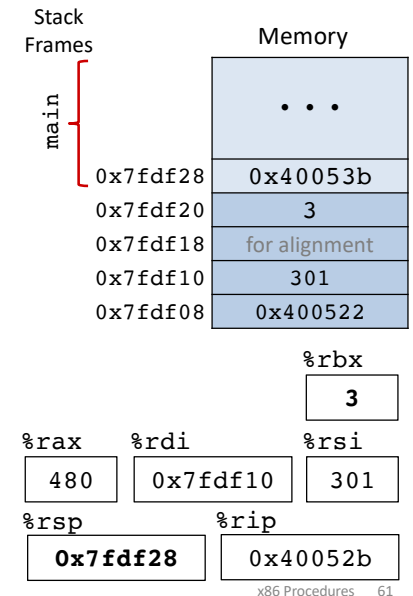


Callee-save example (step 8)

Restore register %rbx
Ready to return

```
long step_by(long x) {
    long v1 = x;
    long v2 = increment(&v1, 61);
    return x + v2;
}
```

```
step_by:
400504: pushq %rbx
400506: movq  %rdi, %rbx
400509: subq  $16, %rsp
40050d: movq  %rdi, (%rsp)
400515: movq  %rsp, %rdi
400518: movl  $61, %esi
40051d: callq 4004cd <increment>
400522: addq  %rbx, %rax
400525: addq  $16, %rsp
400529: popq  %rbx
40052b: retq
```



Recursion example: code

```
long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

pcount:

```
4005dd: movl $0, %eax
4005e2: testq %rdi, %rdi
4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
4005e8: movq %rdi, %rbx
4005eb: andl $1, %ebx
4005ee: shrq %rdi
4005f1: callq pcount
4005f6: addq %rbx, %rax
4005f9: popq %rbx
.L6:
4005fa: rep
4005fb: retq
```

Annotations:

- base case/condition: points to `testq %rdi, %rdi` and `je 4005fa <.L6>`
- recursive case: points to `pushq %rbx`, `movq %rdi, %rbx`, `andl $1, %ebx`, `shrq %rdi`, `callq pcount`, `addq %rbx, %rax`, `popq %rbx`
- x&1 in %rbx across call: points to `andl $1, %ebx`
- save/restore %rbx (callee-save): points to `pushq %rbx` and `popq %rbx`

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Recursion Example: pcount (2)

```
long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

pcount:

```
4005dd: movl $0, %eax
4005e2: testq %rdi, %rdi
4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
4005e8: movq %rdi, %rbx
4005eb: andl $1, %ebx
4005ee: shrq %rdi
4005f1: callq pcount
4005f6: addq %rbx, %rax
4005f9: popq %rbx
.L6:
4005fa: rep
4005fb: retq
```

Stack Frames:

main	0x7fdf38	0x4006ed
	0x7fdf30	
	0x7fdf28	
	0x7fdf20	
	0x7fdf18	
	0x7fdf10	
	0x7fdf08	

Registers:

%rax	0	%rdi	2	%rbx	42
%rsp	0x7fdf38	%rip	0x4005dd		

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Recursion Example: pcount (2)

```
long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

pcount:

```
4005dd: movl $0, %eax
4005e2: testq %rdi, %rdi
4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
4005e8: movq %rdi, %rbx
4005eb: andl $1, %ebx
4005ee: shrq %rdi
4005f1: callq pcount
4005f6: addq %rbx, %rax
4005f9: popq %rbx
.L6:
4005fa: rep
4005fb: retq
```

Stack Frames:

main	0x7fdf38	0x4006ed
pc(2)	0x7fdf30	
	0x7fdf28	
	0x7fdf20	
	0x7fdf18	
	0x7fdf10	
	0x7fdf08	

Registers:

%rax	0	%rdi	2	%rbx	42
%rsp	0x7fdf38	%rip	0x4005e7		

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Recursion Example: pcount (2)

```
long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

pcount:

```
4005dd: movl $0, %eax
4005e2: testq %rdi, %rdi
4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
4005e8: movq %rdi, %rbx
4005eb: andl $1, %ebx
4005ee: shrq %rdi
4005f1: callq pcount
4005f6: addq %rbx, %rax
4005f9: popq %rbx
.L6:
4005fa: rep
4005fb: retq
```

Stack Frames:

main	0x7fdf38	0x4006ed
pc(2)	0x7fdf30	42
	0x7fdf28	
	0x7fdf20	
	0x7fdf18	
	0x7fdf10	
	0x7fdf08	

Registers:

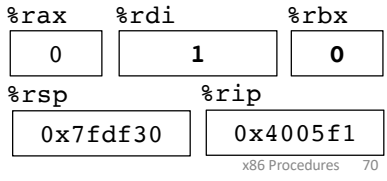
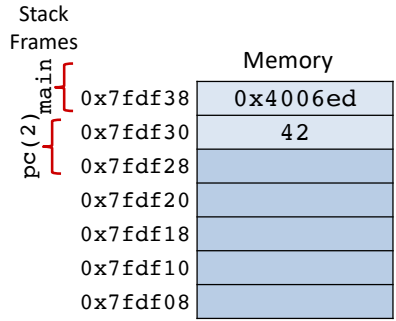
%rax	0	%rdi	2	%rbx	2
%rsp	0x7fdf30	%rip	0x4005eb		

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Recursion Example: pcount(2)

```
long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

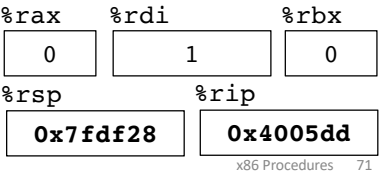
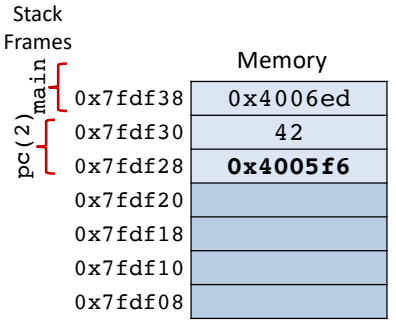
```
pcount:
4005dd: movl $0, %eax
4005e2: testq %rdi, %rdi
4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
4005e8: movq %rdi, %rbx
4005eb: andl $1, %ebx
4005ee: shrq %rdi
4005f1: callq pcount
4005f6: addq %rbx, %rax
4005f9: popq %rbx
.L6:
4005fa: rep
4005fb: retq
```



Recursion Example: pcount(2) → pcount(1)

```
long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

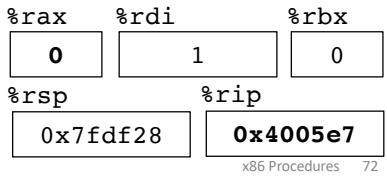
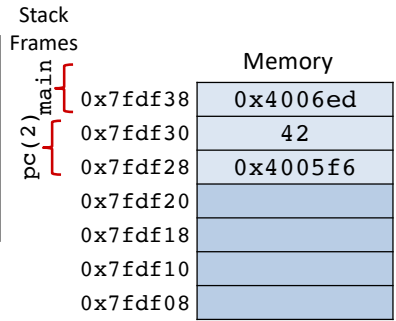
```
pcount:
4005dd: movl $0, %eax
4005e2: testq %rdi, %rdi
4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
4005e8: movq %rdi, %rbx
4005eb: andl $1, %ebx
4005ee: shrq %rdi
4005f1: callq pcount
4005f6: addq %rbx, %rax
4005f9: popq %rbx
.L6:
4005fa: rep
4005fb: retq
```



Recursion Example: pcount(2) → pcount(1)

```
long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

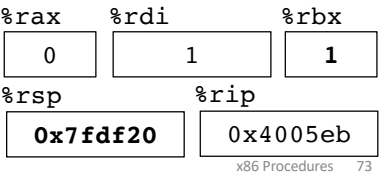
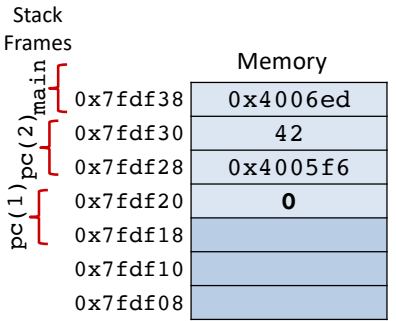
```
pcount:
4005dd: movl $0, %eax
4005e2: testq %rdi, %rdi
4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
4005e8: movq %rdi, %rbx
4005eb: andl $1, %ebx
4005ee: shrq %rdi
4005f1: callq pcount
4005f6: addq %rbx, %rax
4005f9: popq %rbx
.L6:
4005fa: rep
4005fb: retq
```



Recursion Example: pcount(2) → pcount(1)

```
long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

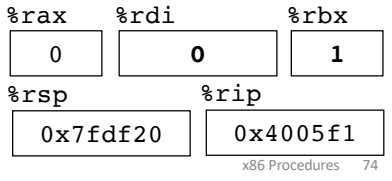
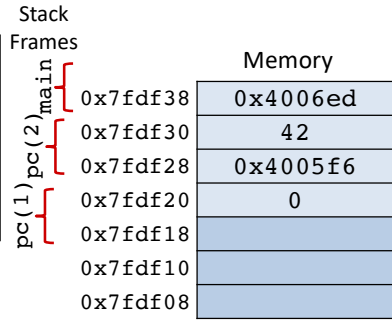
```
pcount:
4005dd: movl $0, %eax
4005e2: testq %rdi, %rdi
4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
4005e8: movq %rdi, %rbx
4005eb: andl $1, %ebx
4005ee: shrq %rdi
4005f1: callq pcount
4005f6: addq %rbx, %rax
4005f9: popq %rbx
.L6:
4005fa: rep
4005fb: retq
```



Recursion Example: pcount(2) → pcount(1)

```
long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

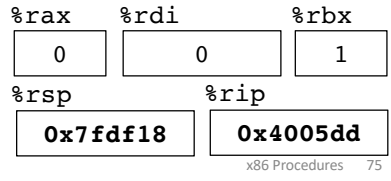
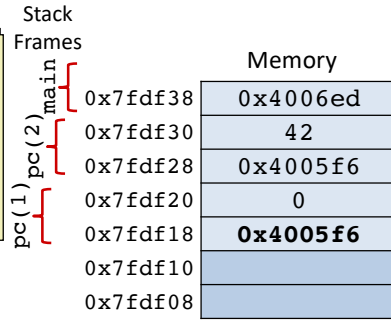
```
pcount:
4005dd: movl $0, %eax
4005e2: testq %rdi, %rdi
4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
4005e8: movq %rdi, %rbx
4005eb: andl $1, %ebx
4005ee: shrq %rdi
4005f1: callq pcount
4005f6: addq %rbx, %rax
4005f9: popq %rbx
.L6:
4005fa: rep
4005fb: retq
```



Recursion Example: pcount(2) → pcount(1) → pcount(0)

```
long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

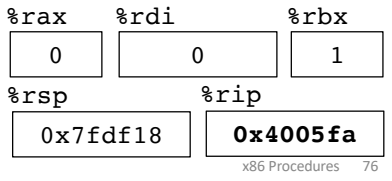
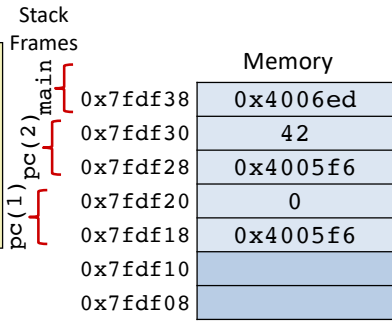
```
pcount:
4005dd: movl $0, %eax
4005e2: testq %rdi, %rdi
4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
4005e8: movq %rdi, %rbx
4005eb: andl $1, %ebx
4005ee: shrq %rdi
4005f1: callq pcount
4005f6: addq %rbx, %rax
4005f9: popq %rbx
.L6:
4005fa: rep
4005fb: retq
```



Recursion Example: pcount(2) → pcount(1) → pcount(0)

```
long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

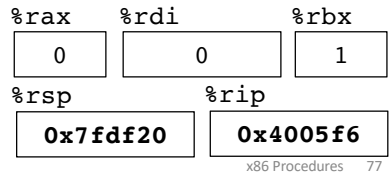
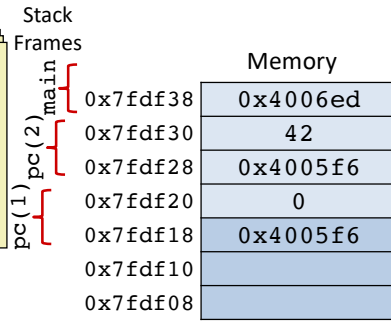
```
pcount:
4005dd: movl $0, %eax
4005e2: testq %rdi, %rdi
4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
4005e8: movq %rdi, %rbx
4005eb: andl $1, %ebx
4005ee: shrq %rdi
4005f1: callq pcount
4005f6: addq %rbx, %rax
4005f9: popq %rbx
.L6:
4005fa: rep
4005fb: retq
```



Recursion Example: pcount(2) → pcount(1) → pcount(0)

```
long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

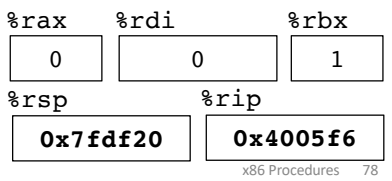
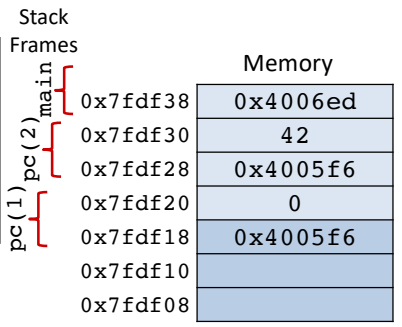
```
pcount:
4005dd: movl $0, %eax
4005e2: testq %rdi, %rdi
4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
4005e8: movq %rdi, %rbx
4005eb: andl $1, %ebx
4005ee: shrq %rdi
4005f1: callq pcount
4005f6: addq %rbx, %rax
4005f9: popq %rbx
.L6:
4005fa: rep
4005fb: retq
```



Recursion Example: pcount(2) → pcount(1) → pcount(0)

```
long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

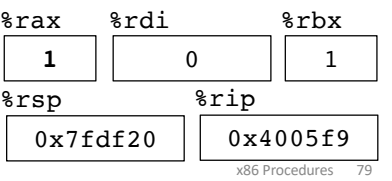
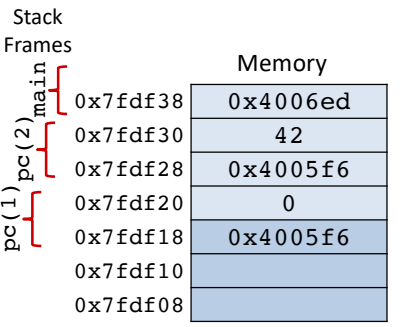
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pcount:
4005dd: movl $0, %eax
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4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
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.L6:
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4005fb: retq
```



Recursion Example: pcount(2) → pcount(1) → pcount(0)

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long pcount(unsigned long x) {
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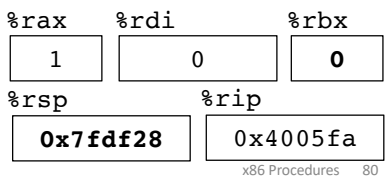
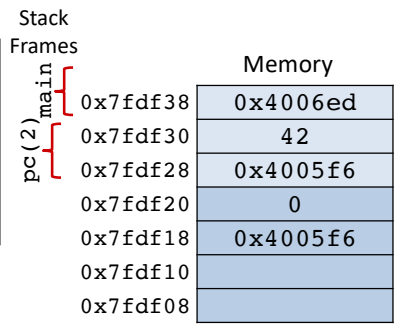
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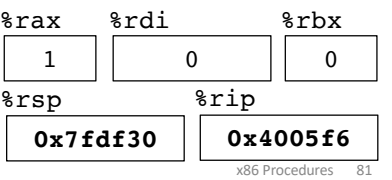
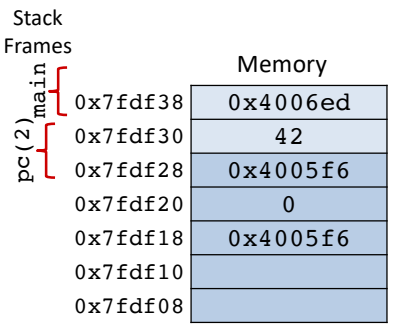
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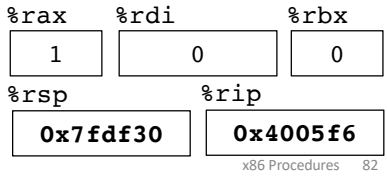
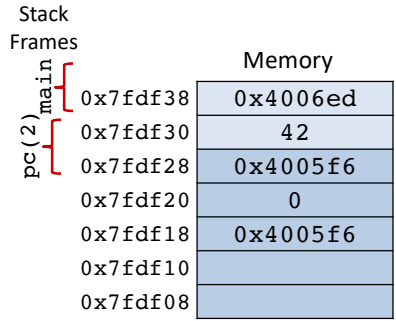
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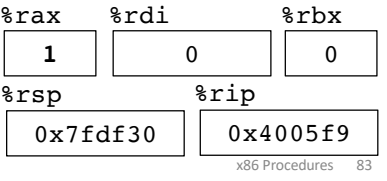
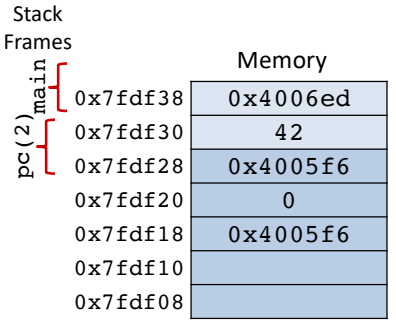
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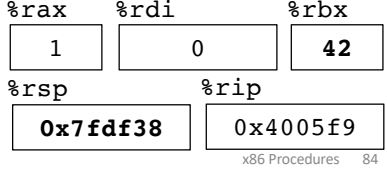
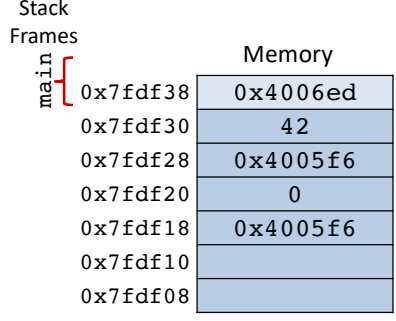
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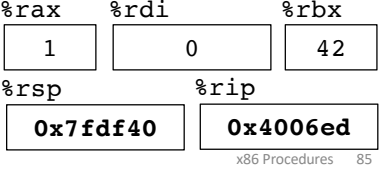
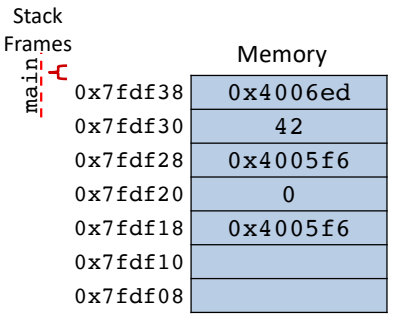
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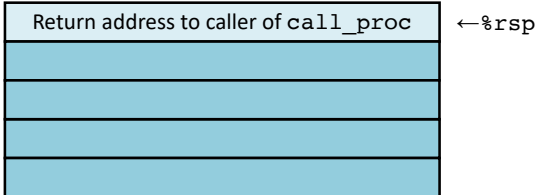


Stack storage example (1)

optional

```
long int call_proc()
{
    long x1 = 1;
    int x2 = 2;
    short x3 = 3;
    char x4 = 4;
    proc(x1, &x1, x2, &x2,
         x3, &x3, x4, &x4);
    return (x1+x2)*(x3-x4);
}
```

```
call_proc:
    subq $32,%rsp
    movq $1,16(%rsp) # x1
    movl $2,24(%rsp) # x2
    movw $3,28(%rsp) # x3
    movb $4,31(%rsp) # x4
    . . .
```



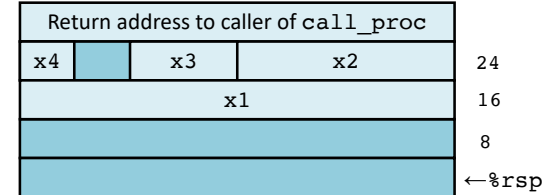
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Stack storage example (2) Allocate local vars

optional

```
long int call_proc()
{
    long x1 = 1;
    int x2 = 2;
    short x3 = 3;
    char x4 = 4;
    proc(x1, &x1, x2, &x2,
         x3, &x3, x4, &x4);
    return (x1+x2)*(x3-x4);
}
```

```
call_proc:
    subq $32,%rsp
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    . . .
```



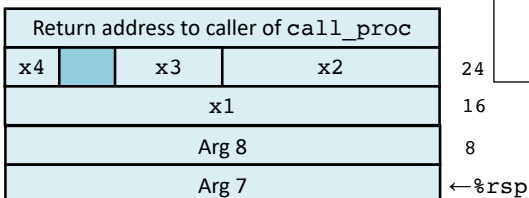
x86 Procedures 87

Stack storage example (3) setup args to proc

optional

```
long int call_proc()
{
    long x1 = 1;
    int x2 = 2;
    short x3 = 3;
    char x4 = 4;
    proc(x1, &x1, x2, &x2,
         x3, &x3, x4, &x4);
    return (x1+x2)*(x3-x4);
}
```

```
call_proc:
    . . .
    leaq 24(%rsp),%rcx # &x2
    leaq 16(%rsp),%rsi # &x1
    leaq 31(%rsp),%rax # &x4
    movq %rax,8(%rsp) # ...
    movl $4,(%rsp) # 4
    leaq 28(%rsp),%r9 # &x3
    movl $3,%r8d # 3
    movl $2,%edx # 2
    movq $1,%rdi # 1
    call proc
    . . .
```



Arguments passed in (in order):
%rdi, %rsi, %rdx, %rcx, %r8, %r9

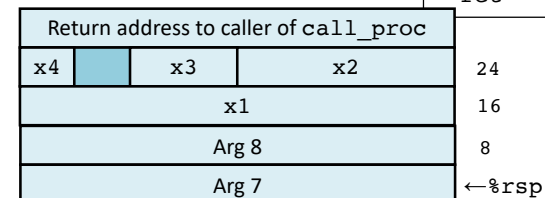
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Stack storage example (4) after call to proc

optional

```
long int call_proc()
{
    long x1 = 1;
    int x2 = 2;
    short x3 = 3;
    char x4 = 4;
    proc(x1, &x1, x2, &x2,
         x3, &x3, x4, &x4);
    return (x1+x2)*(x3-x4);
}
```

```
call_proc:
    . . .
    movswl 28(%rsp),%eax # x3
    movsbl 31(%rsp),%edx # x4
    subl %edx,%eax # x3-x4
    cltq # sign-extend %eax->%rax
    movslq 24(%rsp),%rdx # x2
    addq 16(%rsp),%rdx # x1+x2
    imulq %rdx,%rax # *
    addq $32,%rsp
    ret
```



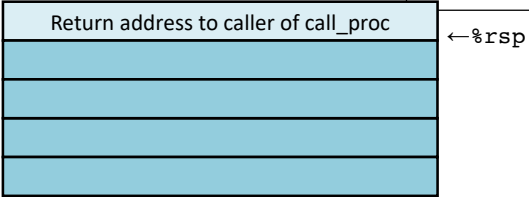
x86 Procedures 89

Stack storage example (5) deallocate local vars

optional

```
long int call_proc()
{
    long  x1 = 1;
    int   x2 = 2;
    short x3 = 3;
    char  x4 = 4;
    proc(x1, &x1, x2, &x2,
         x3, &x3, x4, &x4);
    return (x1+x2)*(x3-x4);
}
```

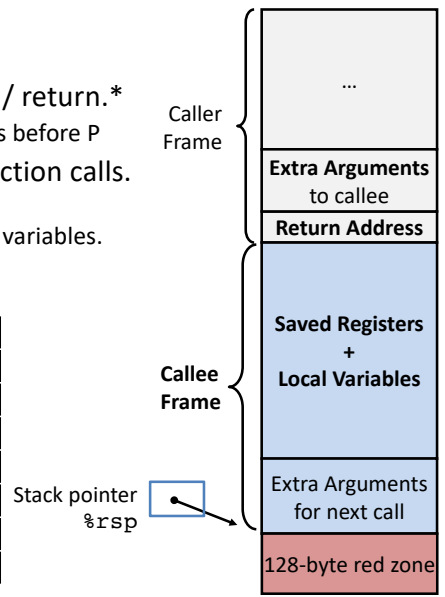
```
call_proc:
    . . .
    movswl 28(%rsp),%eax
    movsbl 31(%rsp),%edx
    subl  %edx,%eax
    cltq
    movslq 24(%rsp),%rdx
    addq  16(%rsp),%rdx
    imulq %rdx,%rax
    addq  $32,%rsp
    ret
```



Procedure Summary

call, ret, push, pop
 Stack discipline fits procedure call / return.*
 If P calls Q: Q (and calls by Q) returns before P
 Conventions support arbitrary function calls.
 Register-save conventions.
 Stack frame saves extra args or local variables.
 Result returned in %rax

%rax	Return value – Caller saved	%r8	Argument #5 – Caller saved
%rbx	Callee saved	%r9	Argument #6 – Caller saved
%rcx	Argument #4 – Caller saved	%r10	Caller saved
%rdx	Argument #3 – Caller saved	%r11	Caller saved
%rsi	Argument #2 – Caller saved	%r12	Callee saved
%rdi	Argument #1 – Caller saved	%r13	Callee saved
%rsp	Stack pointer	%r14	Callee saved
%rbp	Callee saved	%r15	Callee saved



*Take CS 251 to learn about languages where a simple stack does not suffice.