CS 240 Stage 2!

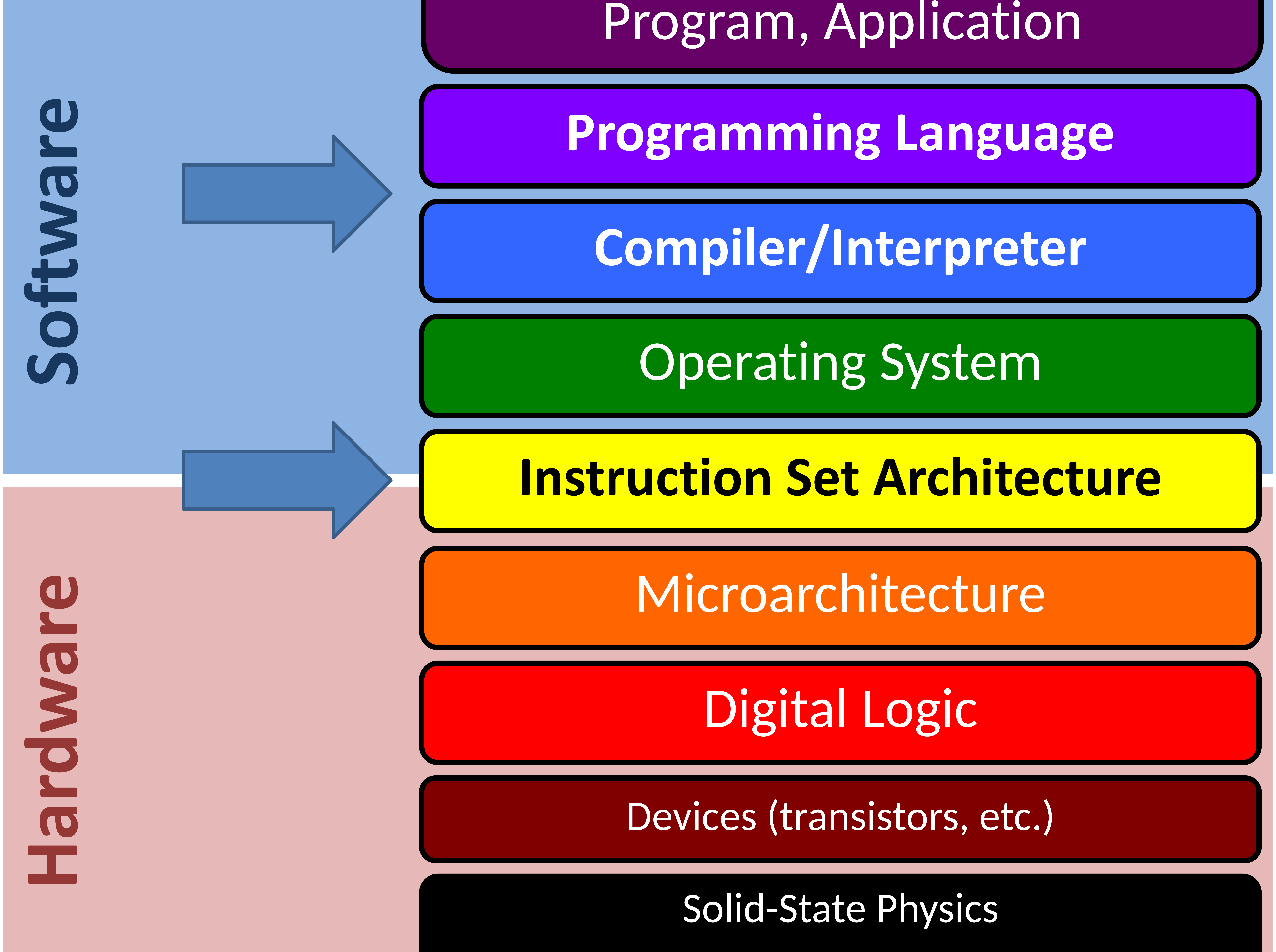
Hardware-Software Interface

Memory addressing, C language, pointers
Assertions, debugging
Machine code, assembly language, program translation
Control flow
Procedures, stacks
Data layout, security, linking and loading

Programming with Memory

the memory model

pointers and arrays in C



Software

Hardware

Program, Application

Programming Language

Compiler/Interpreter

Operating System

Instruction Set Architecture

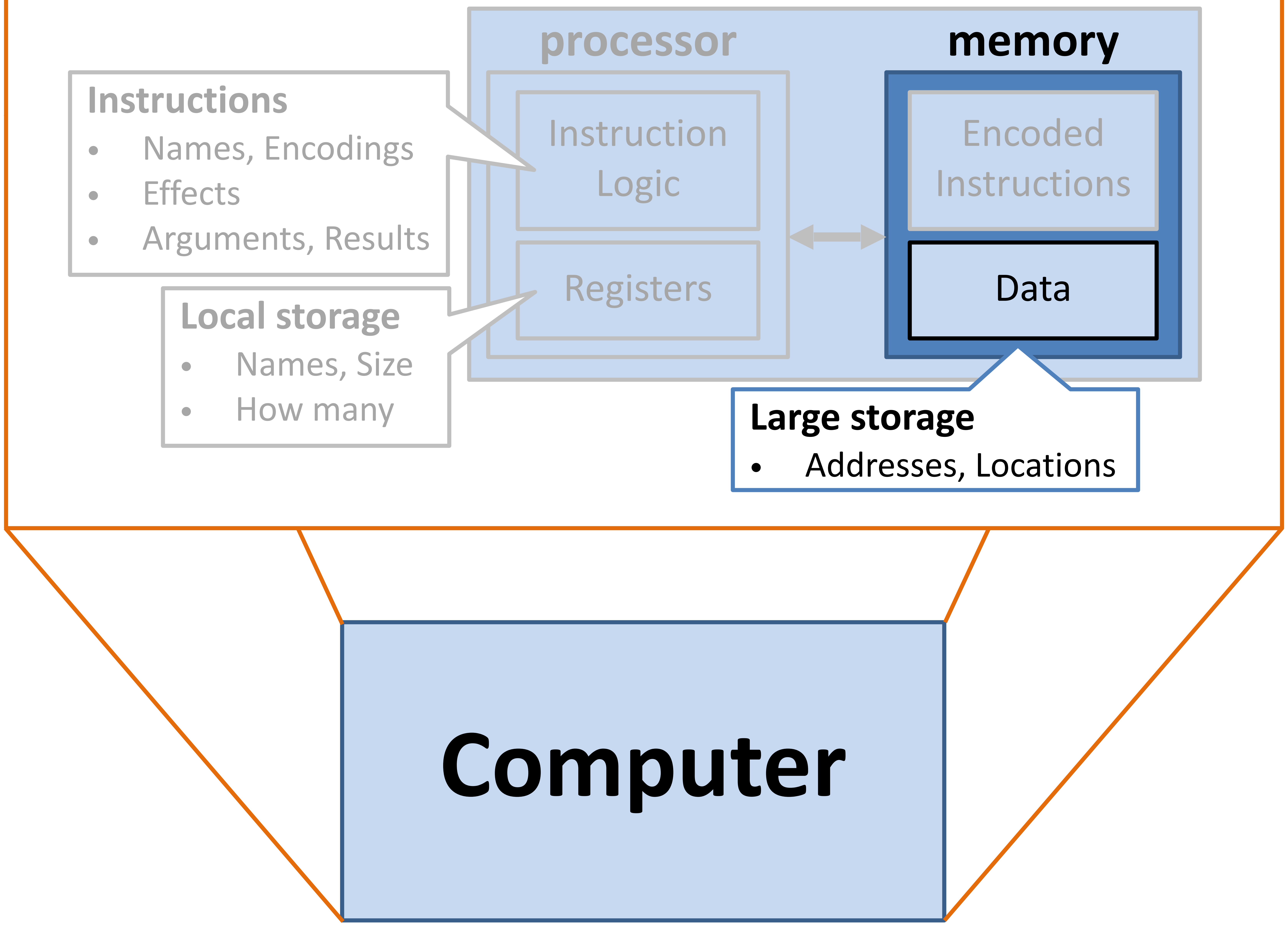
Microarchitecture

Digital Logic

Devices (transistors, etc.)

Solid-State Physics

Instruction Set Architecture (HW/SW Interface)



processor

memory

Instructions

- Names, Encodings
- Effects
- Arguments, Results

Instruction Logic

Encoded Instructions

Local storage

- Names, Size
- How many

Registers

Data

Large storage

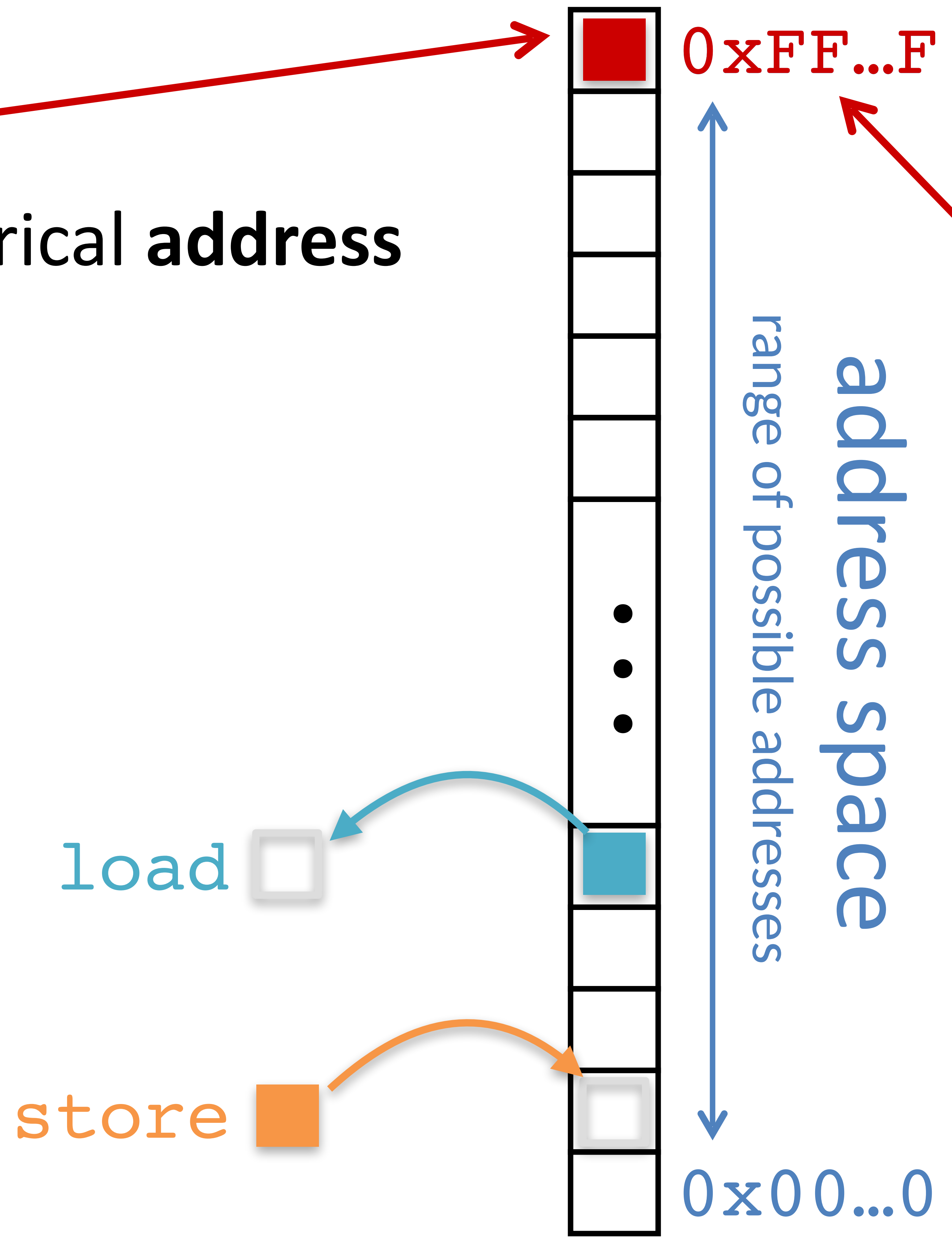
- Addresses, Locations

Computer

Byte-addressable memory = mutable byte array

Location / cell = element

- Identified by unique numerical **address**
- Holds one byte



Address = index

- Unsigned number
- Represented by one word
- Computable and storable as a value

Operations:

- **Load:** read contents at given address
- **Store:** write contents at given address

Multi-byte values in memory

Store across contiguous byte locations.

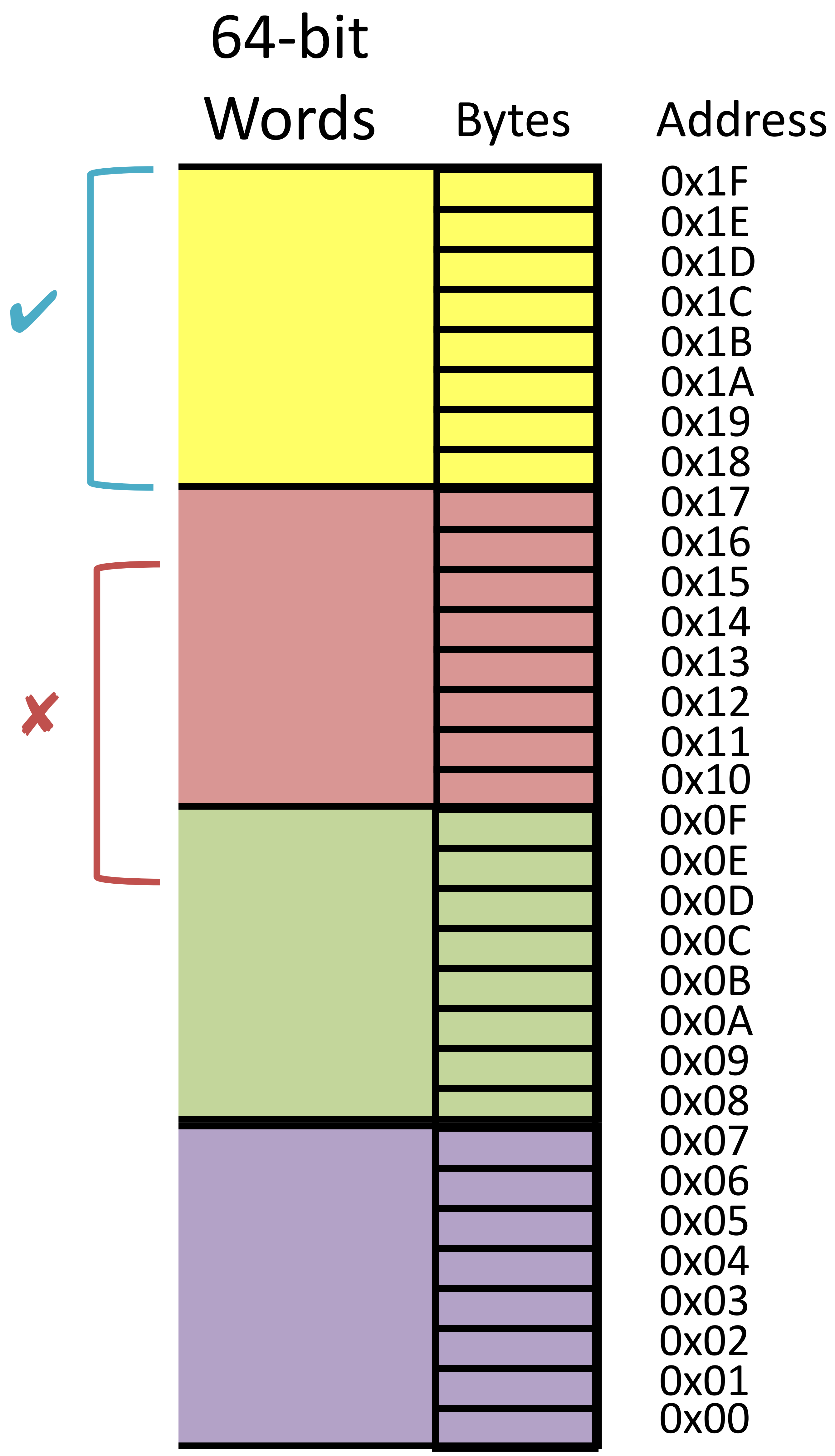
Example: 8 byte (64 bit) values

Alignment

Multi-byte values start at addresses that are multiples of their size

Bit order within byte always same.

Byte ordering within larger value?



Is an `int` stored at address 0x00000002 aligned?

✓ 0

Yes

No

Maybe

Start the presentation to see live content. For screen share software, share the entire screen. Get help at **pollev.com/app**

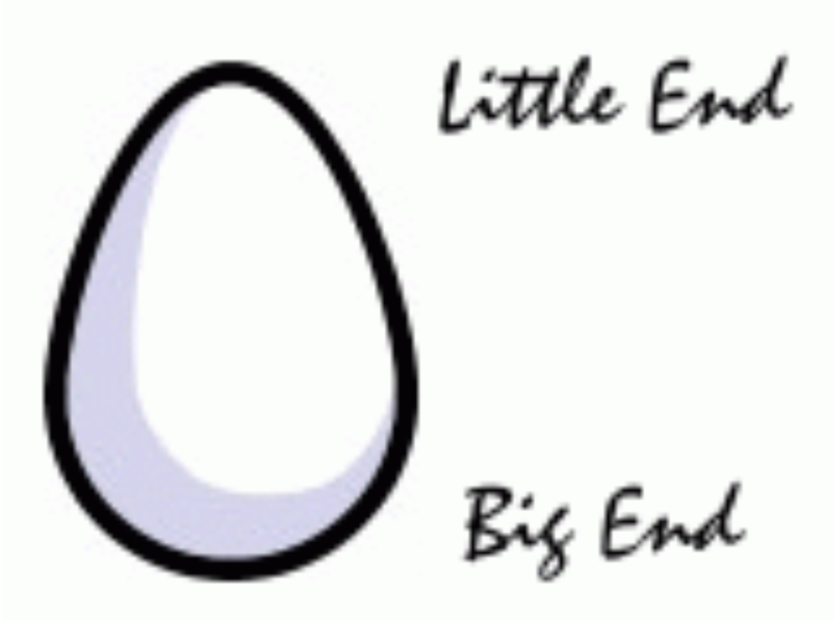
Endianness

In what order are the individual bytes of a multi-byte value stored in memory?

most significant byte

least significant byte

31 30 29 28 27 26 25 24	23 22 21 20 19 18 17 16	15 14 13 12 11 10 9 8	7 6 5 4 3 2 1 0
2A	B6	00	0B



Address	Contents
03	**2A**
02	B6
01	00
00	**0B**

Little Endian: least significant byte first

- low order byte at low address
- high order byte at high address
- used by **x86**, …

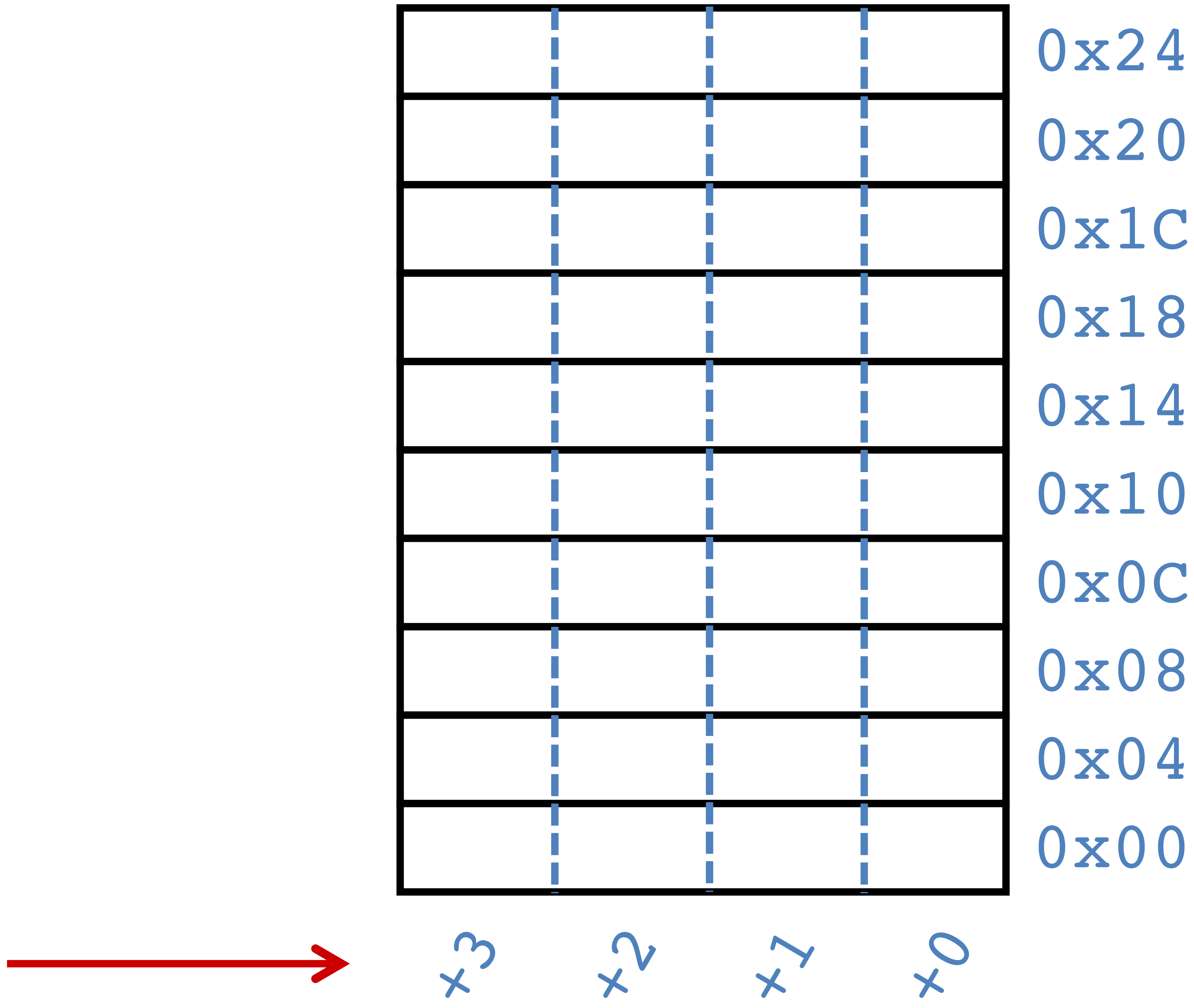
Address	Contents
03	**0B**
02	00
01	B6
00	**2A**

Big Endian: most significant byte first

- high order byte at low address
- low order byte at high address
- used by networks, SPARC, …

Data, addresses, and pointers

For these slides, we'll draw the bytes in this reverse order so that multi-byte values can be read directly



memory drawn as 32-bit values, little endian order

Data, addresses, and pointers

address = index of a location in memory

pointer = a reference to a location in memory,
represented as an address stored as data

Let's store the number 240 at address `0x20`.

240_{10} = $F0_{16}$ = `0x00 00 00 F0`

At address `0x08` we store a pointer to the contents at address `0x20`.

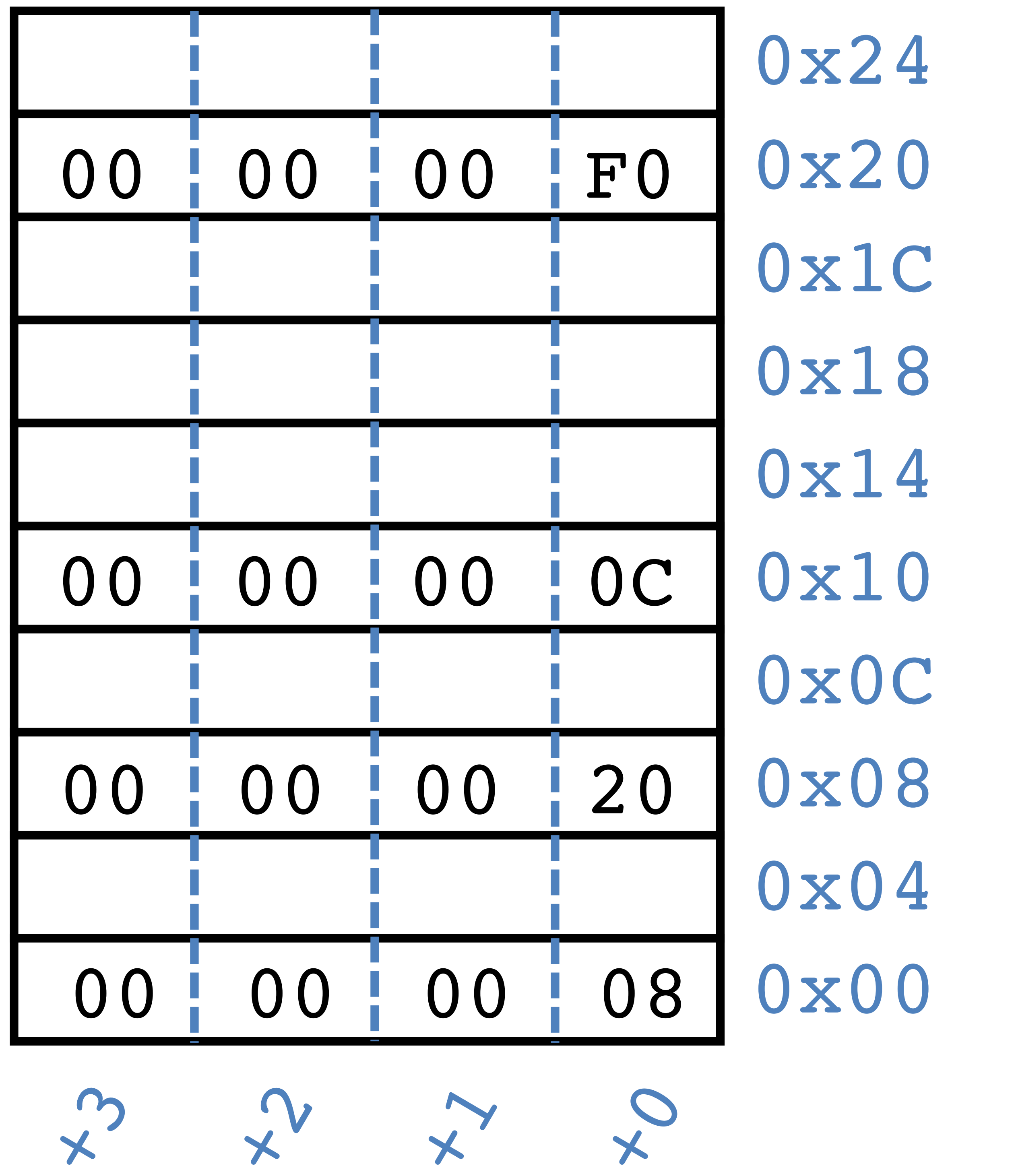
At address `0x00`, we store a pointer to a pointer.

The number 12 is stored at address `0x10`.

Is it a pointer?

How do we know if values are pointers or not?

How do we manage use of memory?



+3	+2	+1	+0	
				0x24
00	00	00	F0	0x20
				0x1C
				0x18
				0x14
00	00	00	0C	0x10
				0x0C
00	00	00	20	0x08
				0x04
00	00	00	08	0x00

memory drawn as 32-bit values,
little endian order

C: Variables are locations

The compiler creates a map from variable name → location.

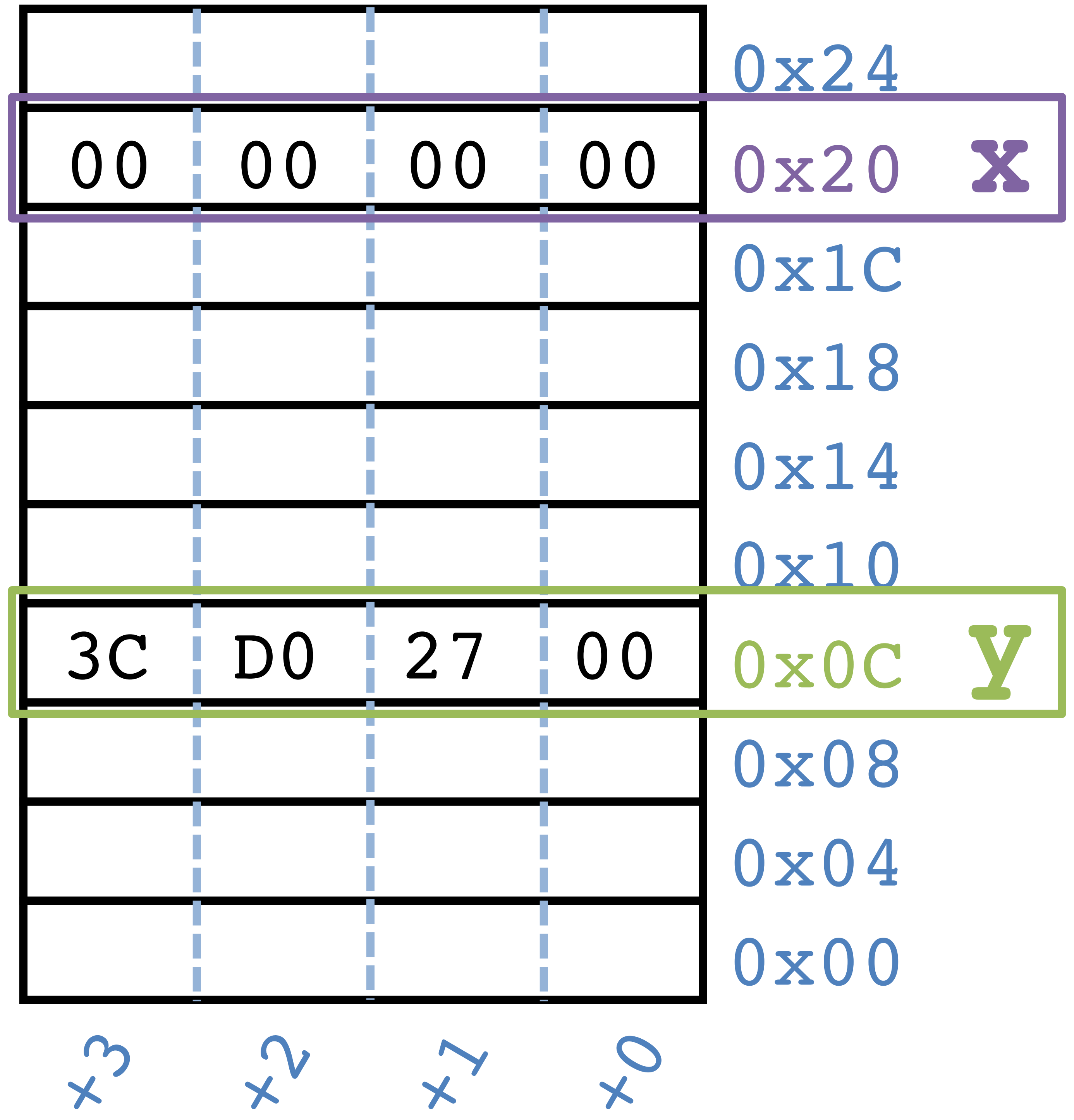
Declarations do not initialize!

```
int x; // x @ 0x20
int y; // y @ 0x0C

x = 0; // store 0 @ 0x20

// store 0x3CD02700 @ 0x0C
y = 0x3CD02700;

// 1. load the contents @ 0x0C
// 2. add 3
// 3. store sum @ 0x20
x = y + 3;
```



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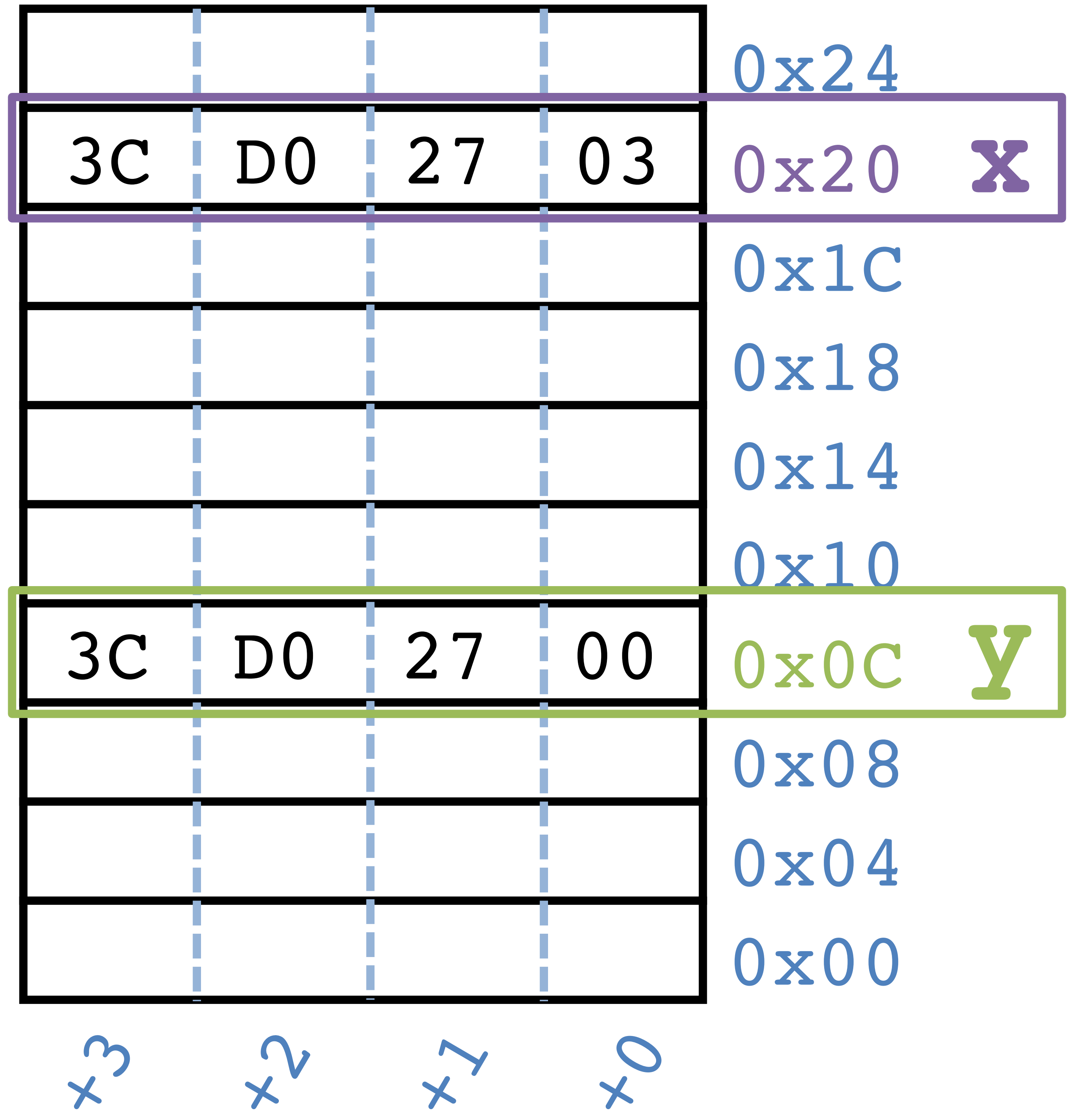


C: Pointer operations and types

address = index of a location in memory

pointer = a *reference* to a location in memory, an address stored as data

Expressions using addresses and pointers:

&___ **address of** the memory location representing ___

a.k.a. "reference to ___"

*___ **contents at** the memory address given by ___

a.k.a. "dereference ___"

Pointer types:

___* address of a memory location holding a ___

a.k.a. "a reference to a ___"

C: Types determine sizes

Sizes of data types (in bytes)

Java Data Type	C Data Type	32-bit word	64-bit word
boolean	*bool*	1	1
byte	char	1	1
char		2	2
short	short int	2	2
int	int	4	4
float	float	4	4
	long int	4	8
double	double	8	8
long	long long	8	8
	long double	8	16
(reference)	(pointer) *	4	8

address size = word size

C: Pointer example

& = address of

** = contents at*

```
int* p;
```

Declare a variable, p

that will hold the address of a memory location holding an int

```
int x = 5;
int y = 2;
```

Declare two variables, x and y, that hold ints, and store 5 and 2 in them, respectively.

```
p = &x;
```

Take the address of the memory

representing x

... and store it in the memory location representing p.
Now, *"p points to x."*

```
y = 1 + *p;
```

Add 1 to

the contents of memory at the address

given by the contents of the memory location representing p

... and store it in the memory location representing y.

C: Pointer example

& = address of

** = contents at*

location

C assignment:

***Left-hand-side* = *right-hand-side*;**

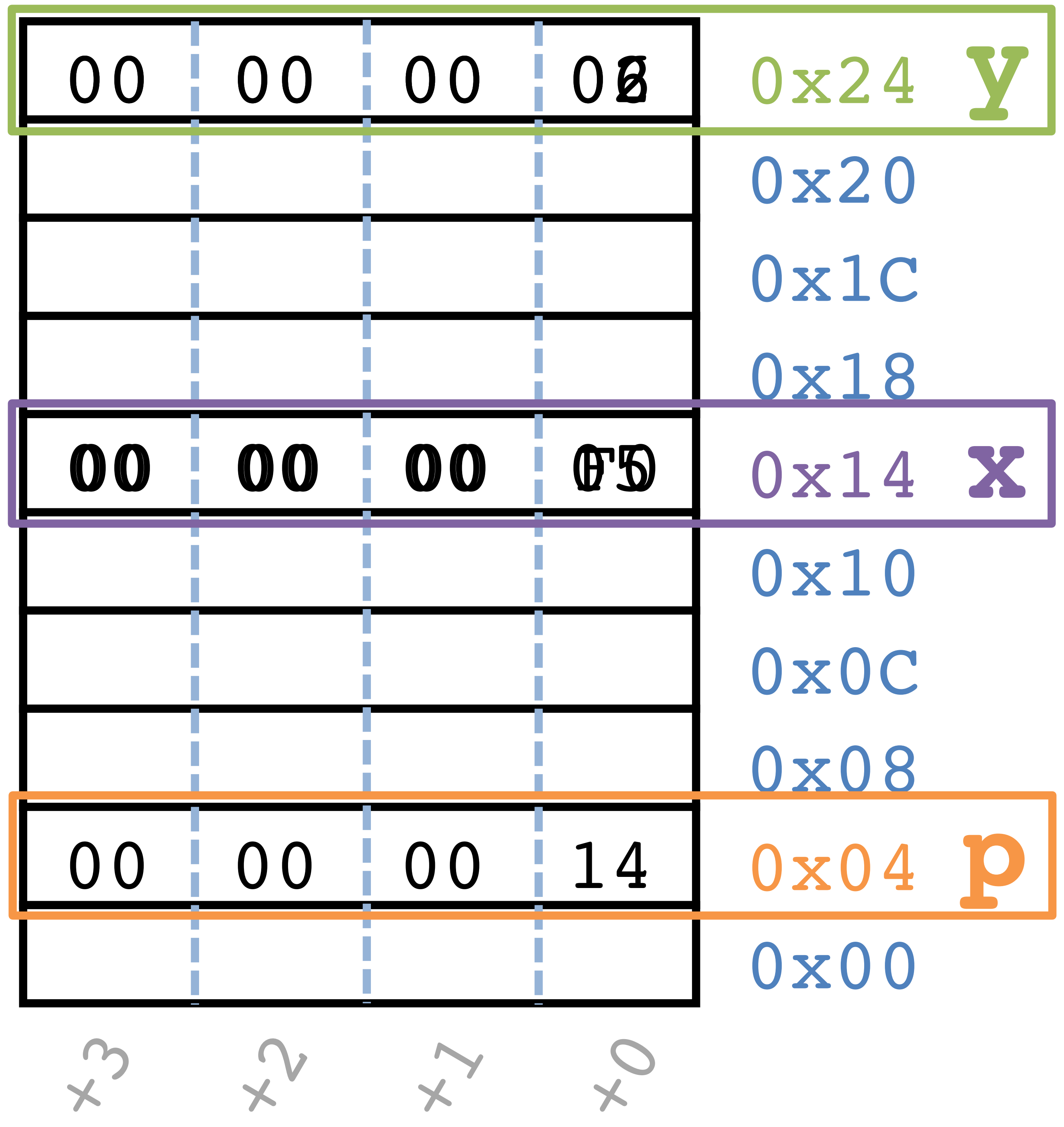
value

What is the type of *p?

What is the type of &x?

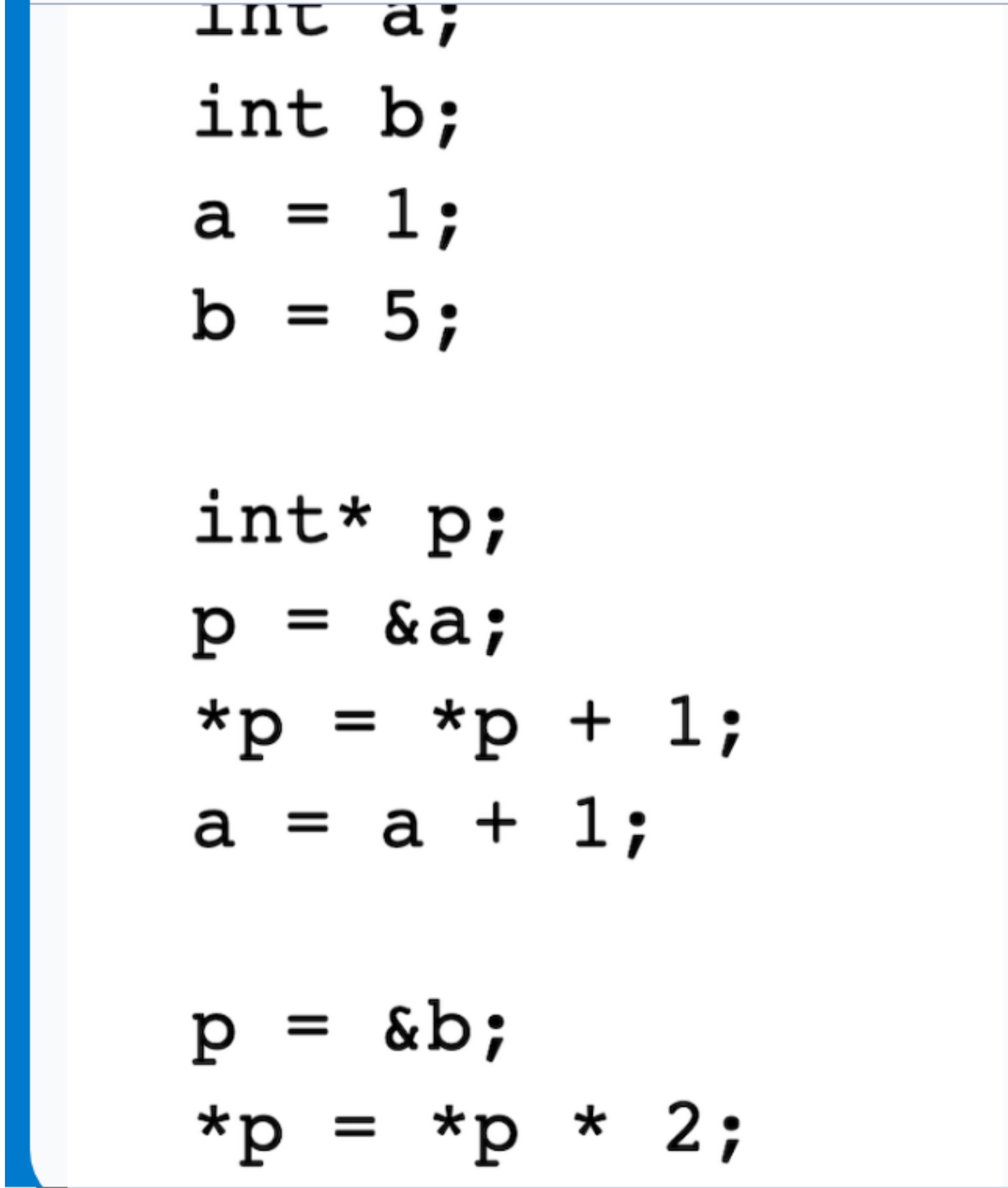
What is `*(&y)` ?

```
int* p;     // p @ 0x04
int x = 5;  // x @ 0x14, store 5 @ 0x14
int y = 2;  // y @ 0x24, store 2 @ 0x24
p = &x;     // store 0x14 @ 0x04

// 1. load the contents @ 0x04  (=0x14)
// 2. load the contents @ 0x14  (=0x5)
// 3. add 1
// 4. store sum as contents @ 0x24
y = 1 + *p;

// 1. load the contents @ 0x04  (=0x14)
// 2. store 0xF0 as contents @ 0x14
*p = 240;
```



+3	+2	+1	+0		
00	00	00	06	0x24	y
				0x20	
				0x1C	
				0x18	
00	00	00	F0	0x14	x
				0x10	
				0x0C	
				0x08	
00	00	00	14	0x04	p
				0x00	

What is the result of printing the decimal values of `a` and `b` at the end of this code?



```
int a;
int b;
a = 1;
b = 5;

int* p;
p = &a;
*p = *p + 1;
a = a + 1;

p = &b;
*p = *p * 2;
```

- 2, 10
- 3, 5
- 3, 10
- 6, 5
- None of the above

C: Pointer type syntax

Spaces between base type, *, and variable name mostly do not matter.

The following are **equivalent**:

```
int* ptr;
```

I see: "The variable **ptr** holds an **address of an int** in memory."

```
int * ptr;
```

```
int *ptr;
```

more common C style

Looks like: "Dereferencing the variable **ptr** will yield an **int**."

Or "The **memory location** where the variable **ptr** points holds an **int**."

Caveat: do not declare multiple variables unless using the last form.

`int* a, b;` means `int *a, b;` means `int* a; int b;`

C: Arrays

Arrays are adjacent memory locations storing the same type of data.

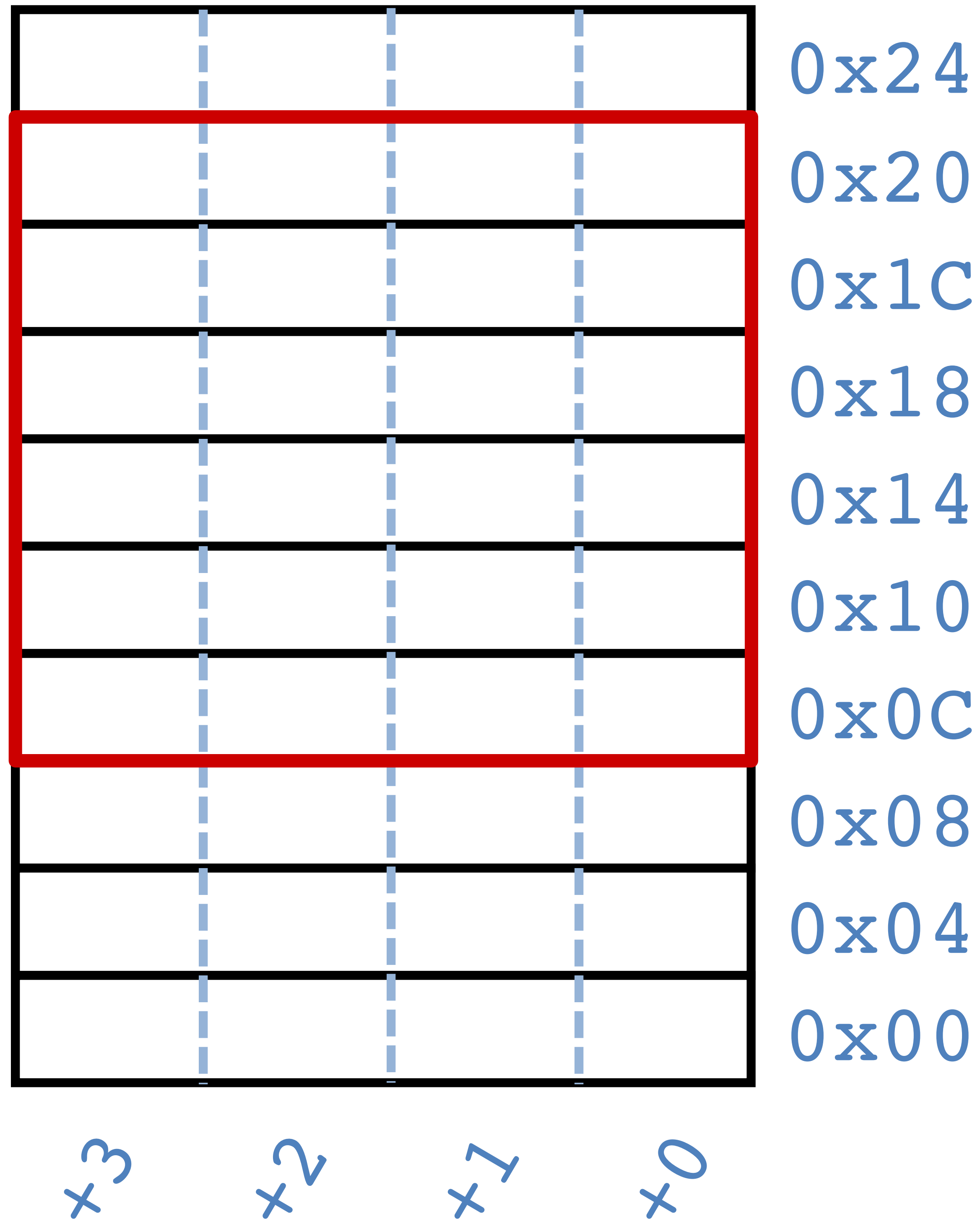
a is a name for the array's base address, can be used as an *immutable* pointer.

Declaration: `int a[6];`

element type

name

number of elements



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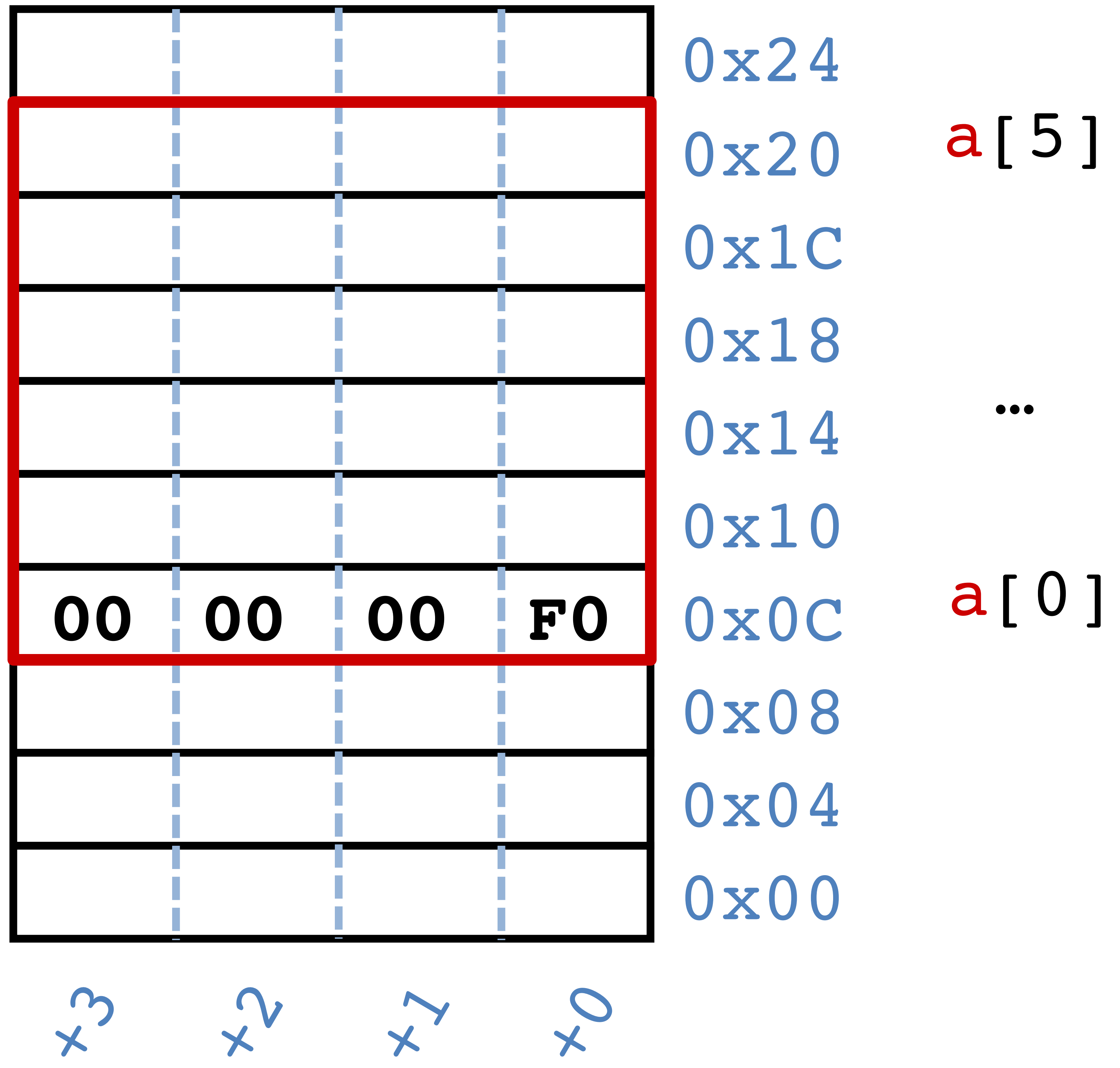
Address of `a[i]` is base address a plus `i` times element size in bytes.

Declaration: `int a[6];`

Indexing: `a[0] = 0xf0;`



C: Arrays

Declaration:

```
int a[6];
```

Indexing:

```
a[0] = 0xf0;
a[5] = a[0];
```

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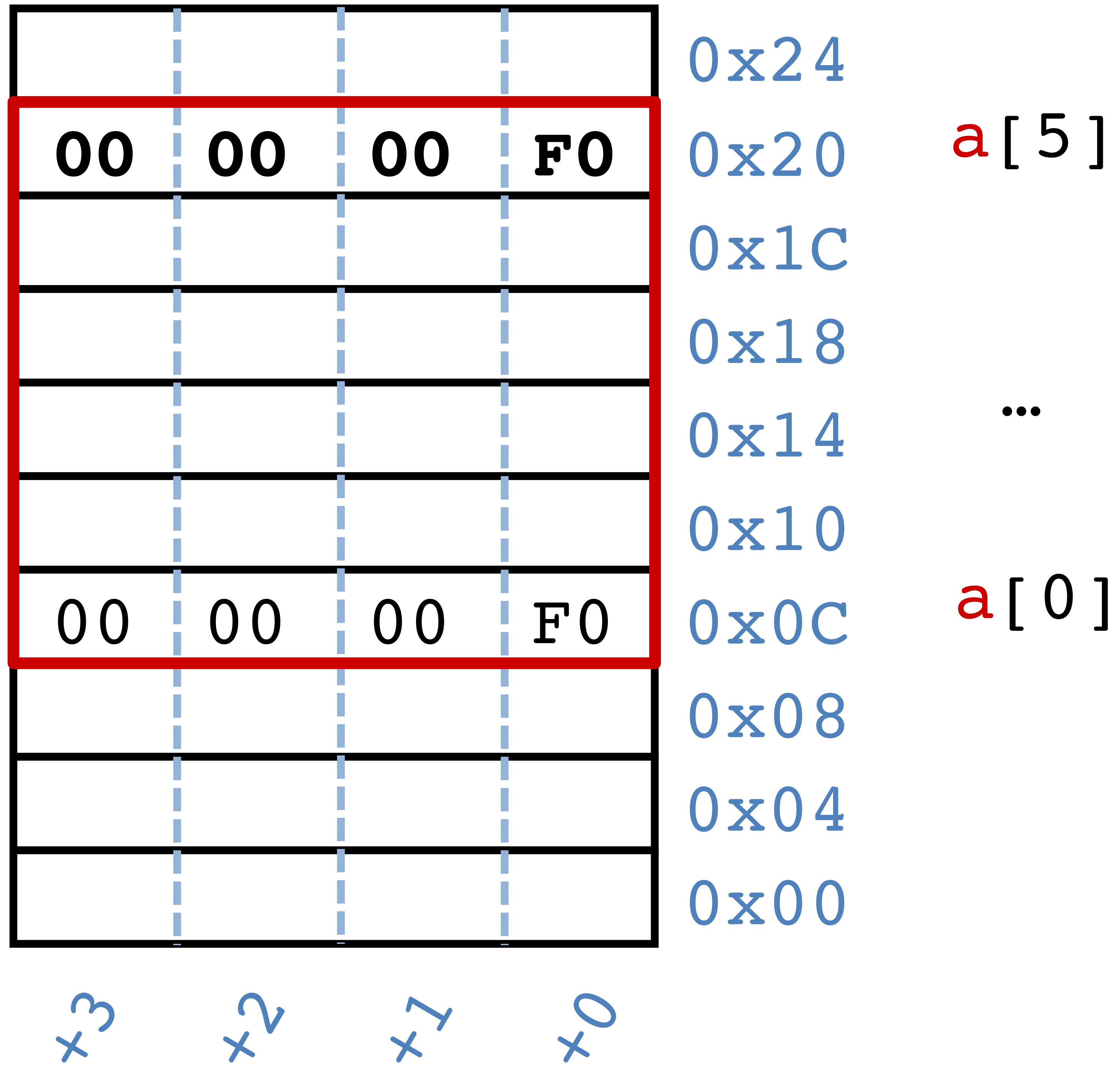

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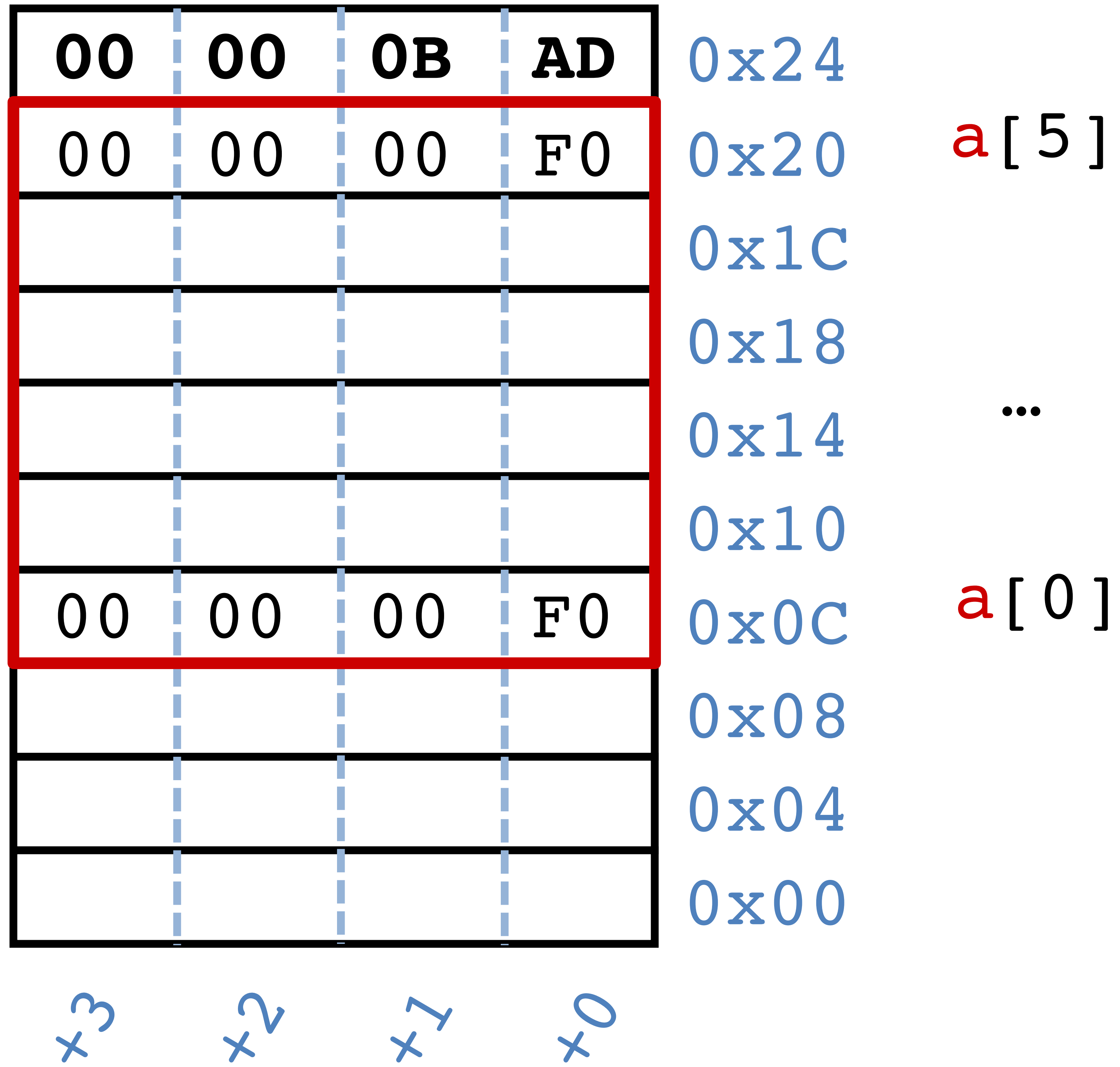
No bounds check: `a[6] = 0xBAD;`



	+3	+2	+1	+0	
0x24	00	00	0B	AD	
0x20	00	00	00	F0	a[5]
0x1C					
0x18					
0x14					…
0x10					
0x0C	00	00	00	F0	a[0]
0x08					
0x04					
0x00					

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```

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```
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No bounds check:

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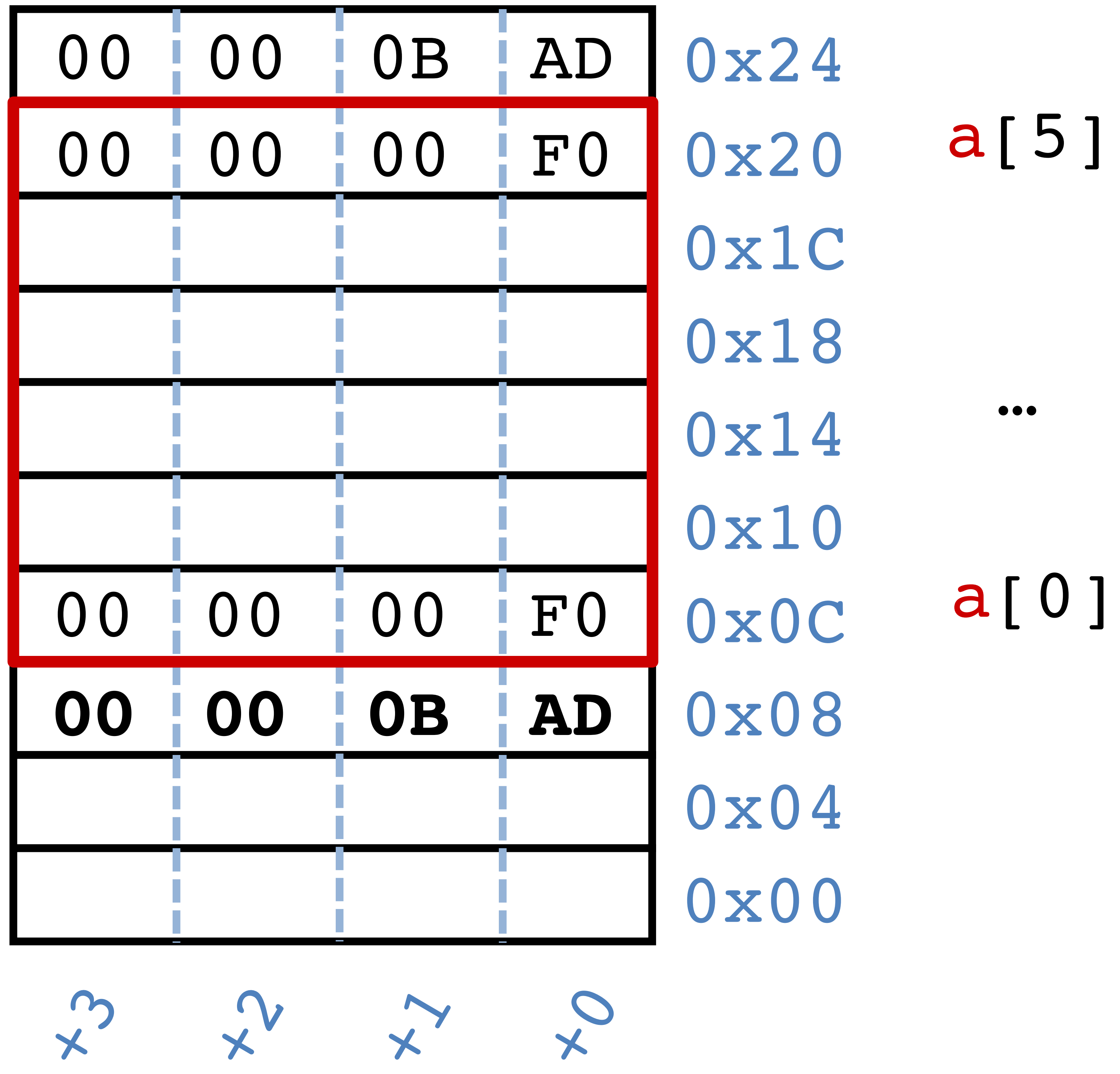

+3	+2	+1	+0		
00	00	0B	AD	0x24	
00	00	00	F0	0x20	a[5]
				0x1C	
				0x18	
				0x14	…
				0x10	
00	00	00	F0	0x0C	a[0]
00	**00**	**0B**	**AD**	0x08	
				0x04	
				0x00	

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```

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a[6] = 0xBAD;
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```
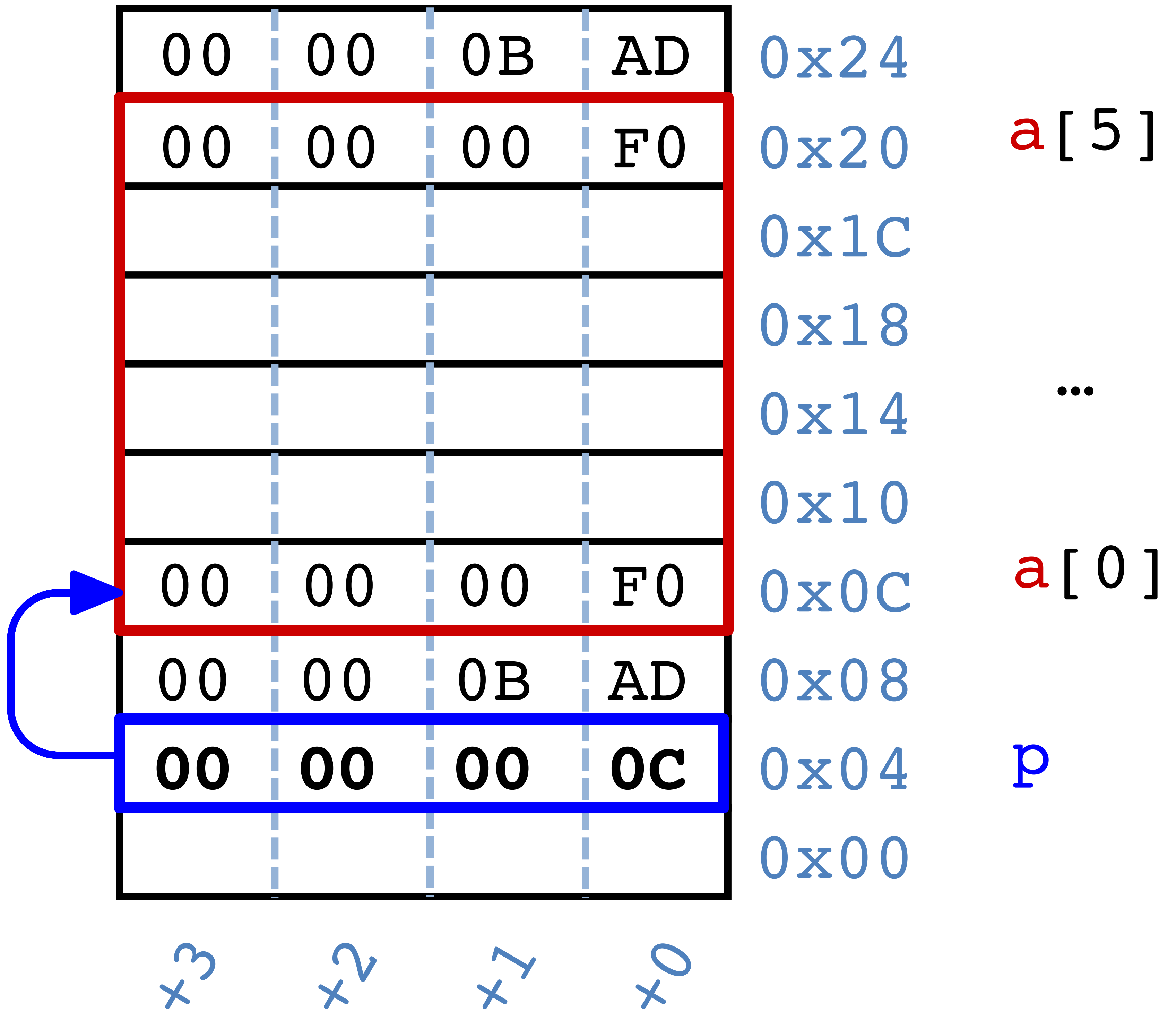

Pointers:

```
int* p;
```

equivalent {

```
p = a;
p = &a[0];
```



+3	+2	+1	+0		
00	00	0B	AD	0x24	
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				0x1C	
				0x18	
				0x14	…
				0x10	
00	00	00	F0	0x0C	a[0]
00	00	0B	AD	0x08	
00	**00**	**00**	**0C**	0x04	p
				0x00	

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```

No bounds check:

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a[6] = 0xBAD;
a[-1] = 0xBAD;
```

Pointers:

```
int* p;
p = a;
p = &a[0];
*p = 0xA;
```

equivalent: `p = a;` and `p = &a[0];`

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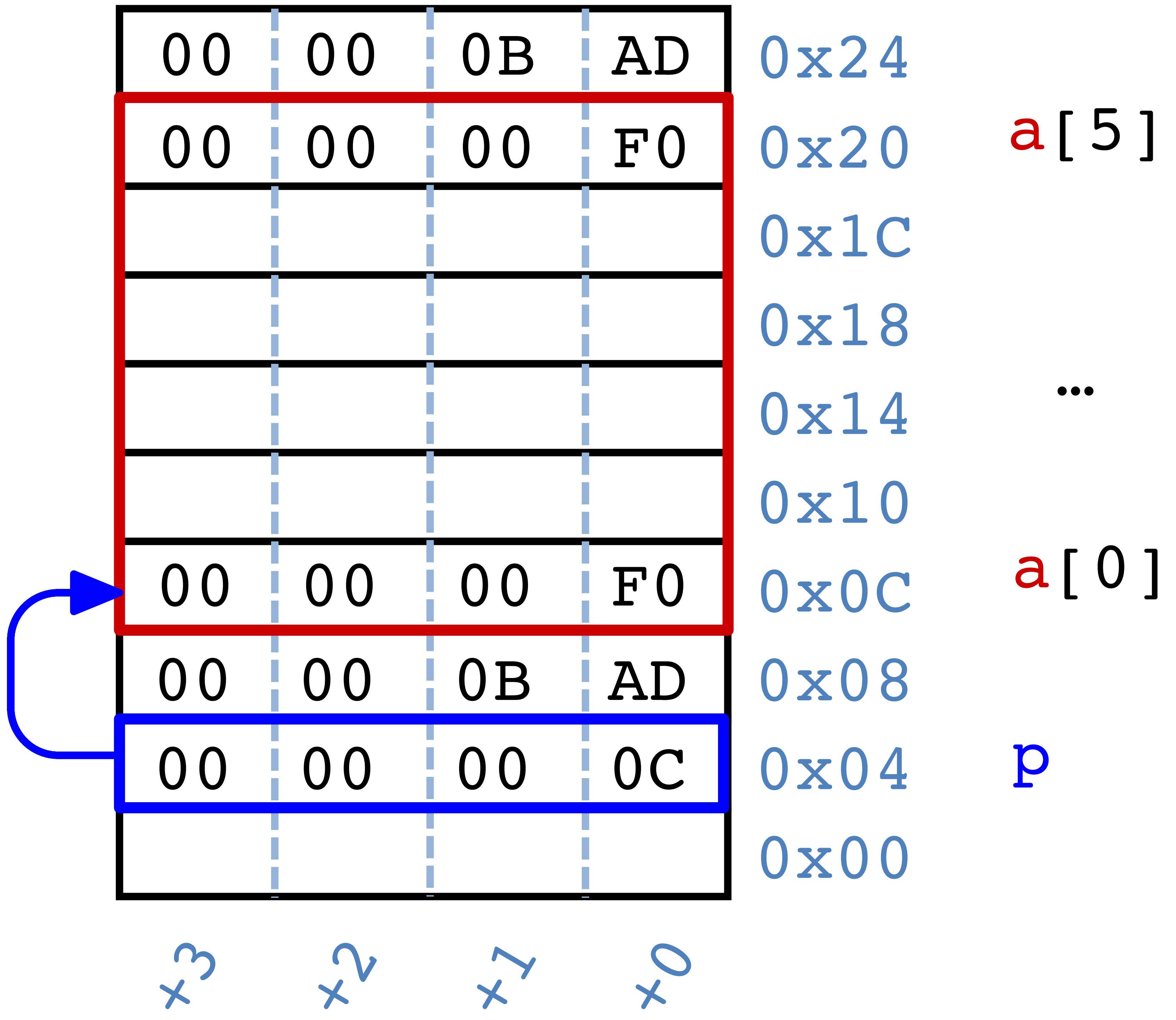

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```

Pointers:

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int* p;
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p = &a[0];
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```
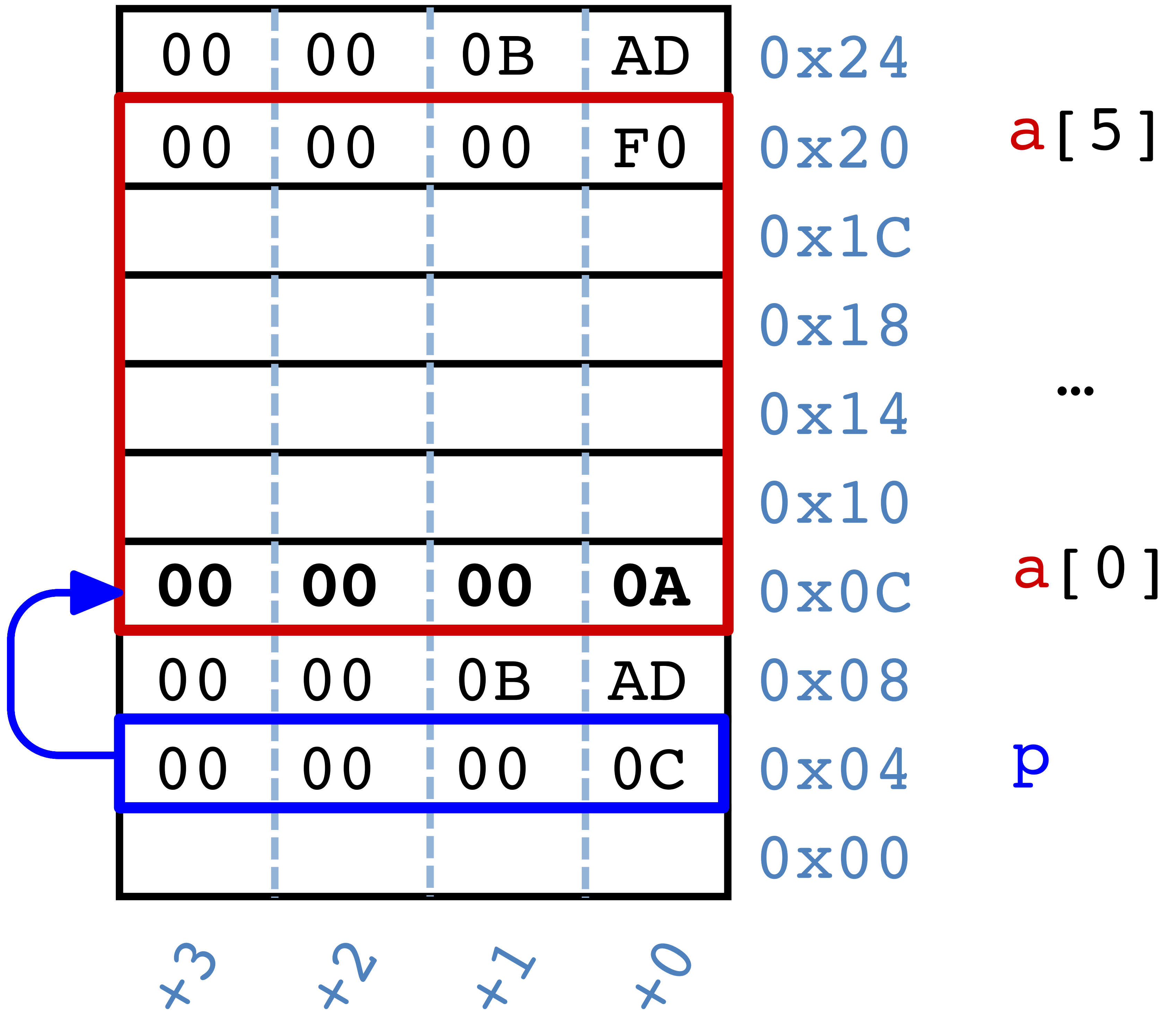

equivalent { p = a; p = &a[0];



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equivalent {

```
p = a;
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```

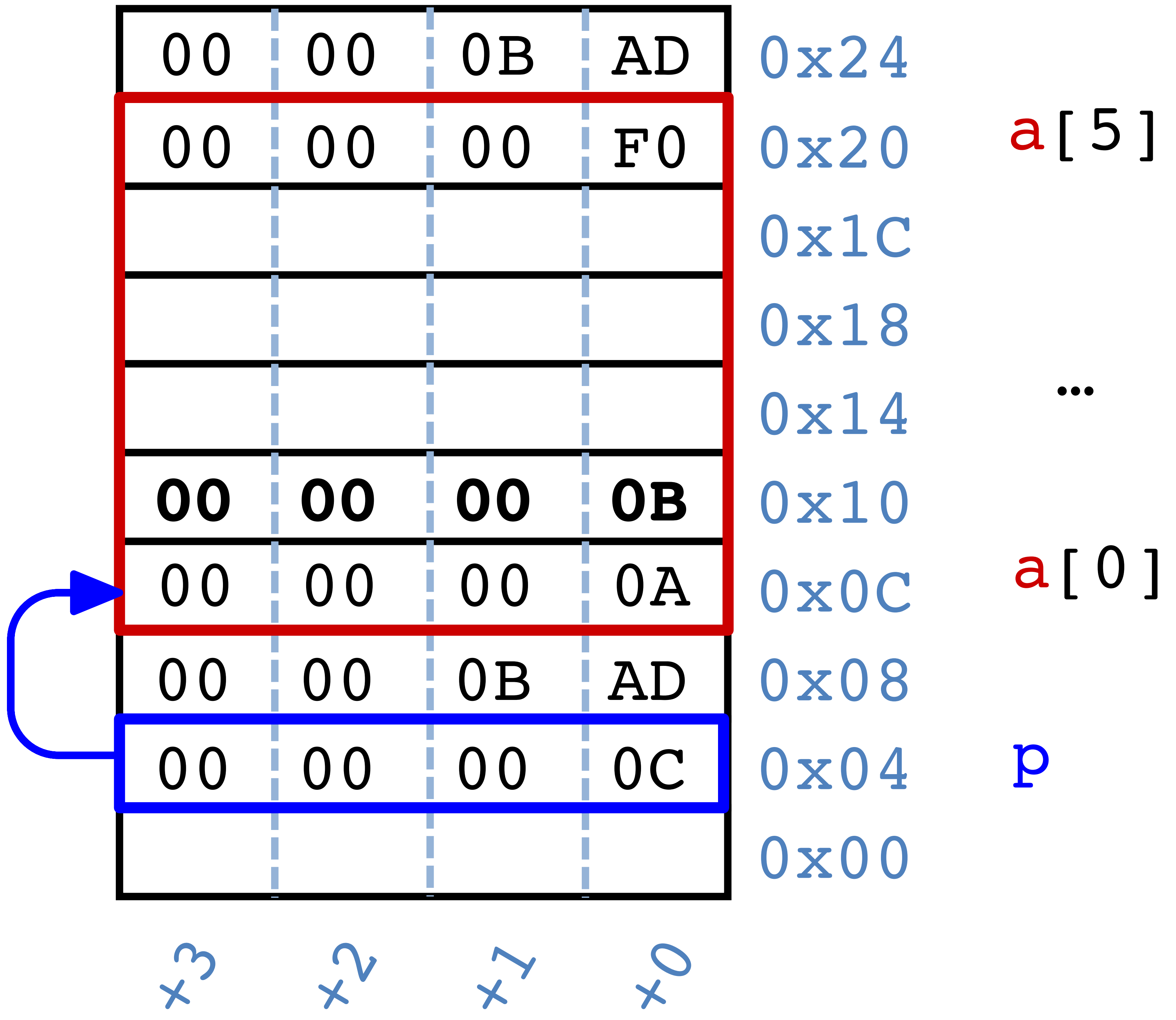
```
*p = 0xA;
```

equivalent {

```
p[1] = 0xB;
*(p + 1) = 0xB;
```

array indexing = address arithmetic

Both are scaled by the size of the type.



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				0x00	
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```

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```

equivalent

```
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```

```
*p = 0xA;
```

equivalent

```
p[1] = 0xB;
*(p + 1) = 0xB;
```

```
p = p + 2;
```

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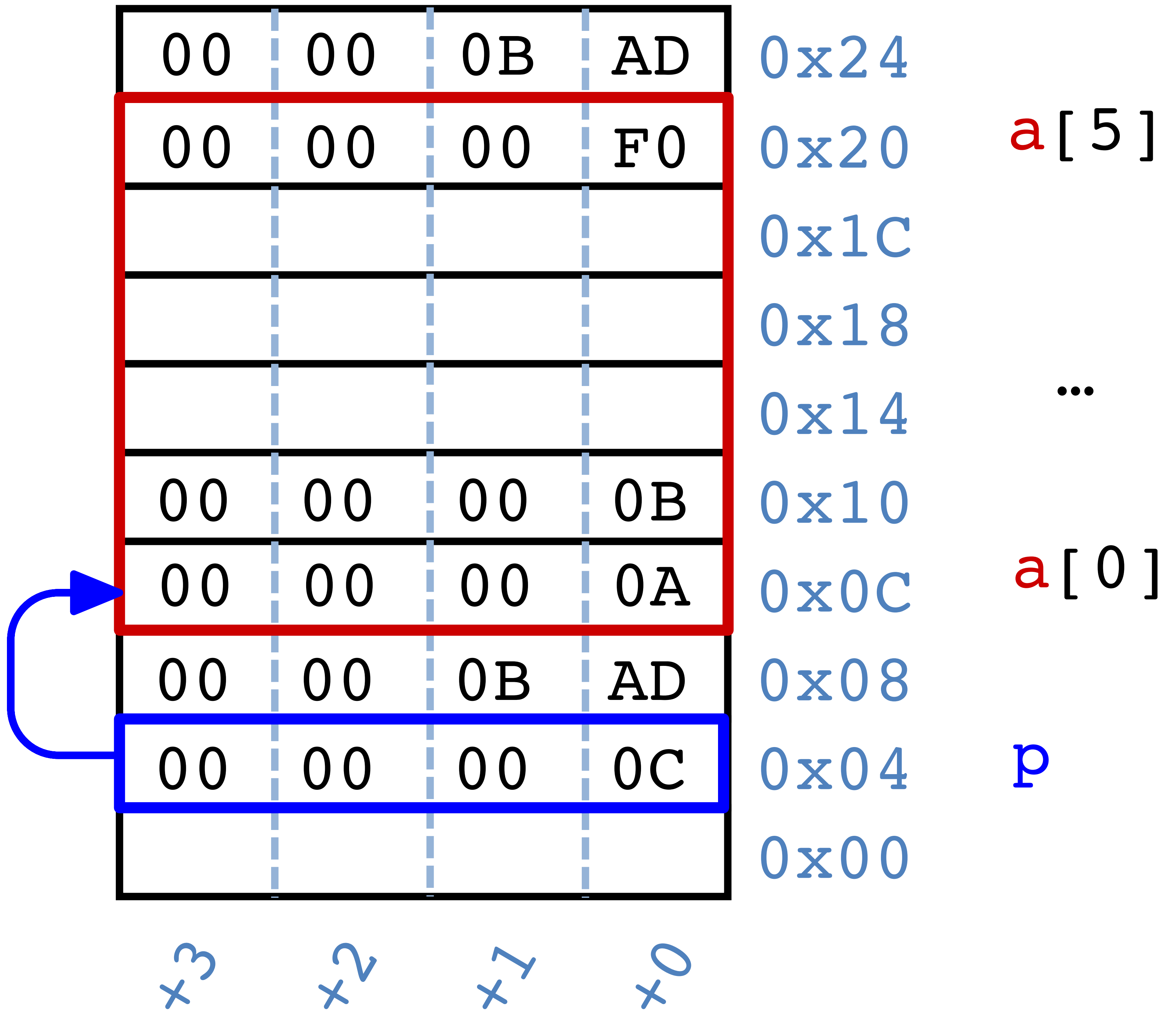

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				0x18	
				0x14	…
00	00	00	0B	0x10	
00	00	00	0A	0x0C	a[0]
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				0x00	
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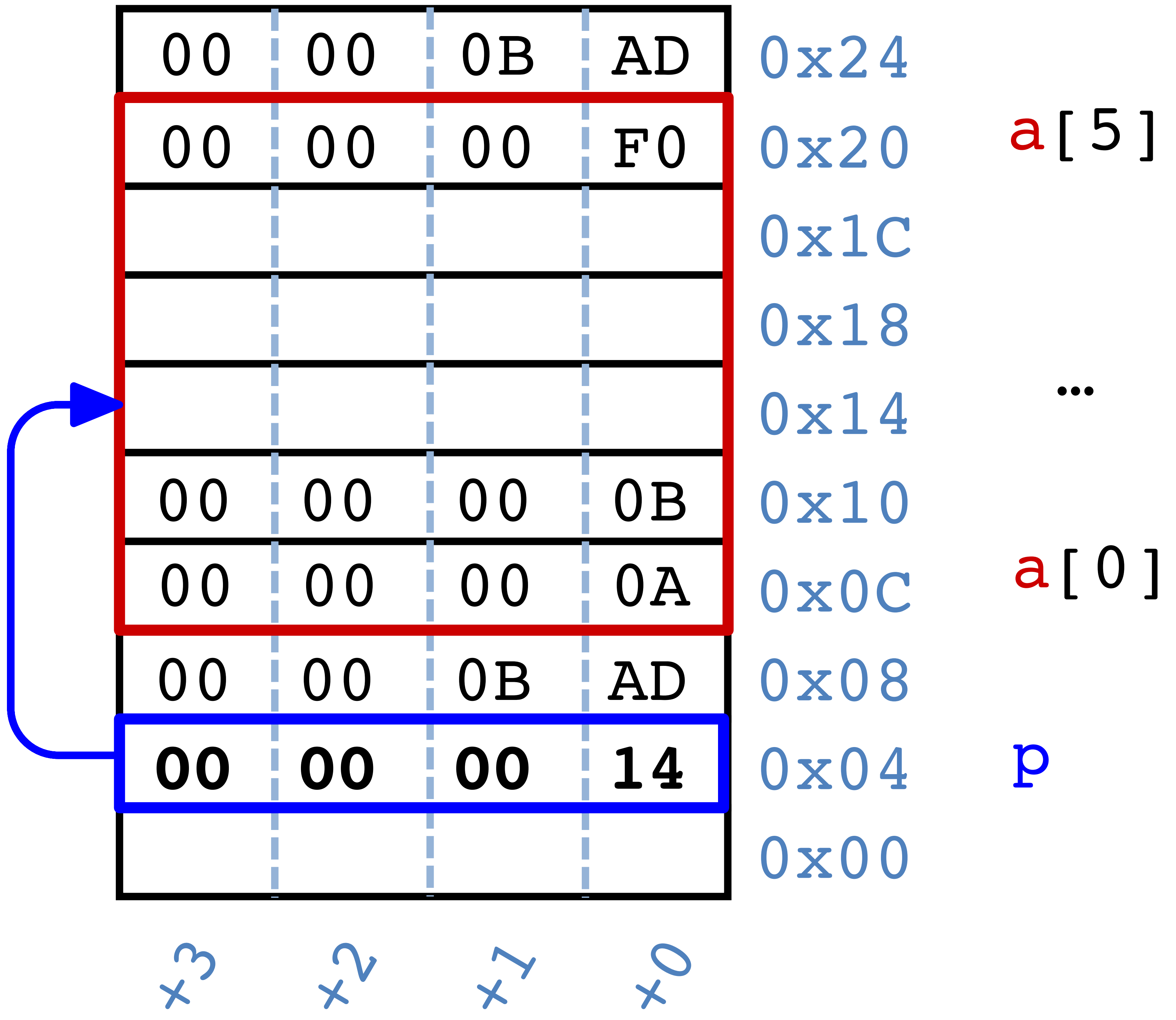

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00	00	0B	AD	0x24	
00	00	00	F0	0x20	a[5]
				0x1C	
				0x18	
				0x14	…
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				0x00	

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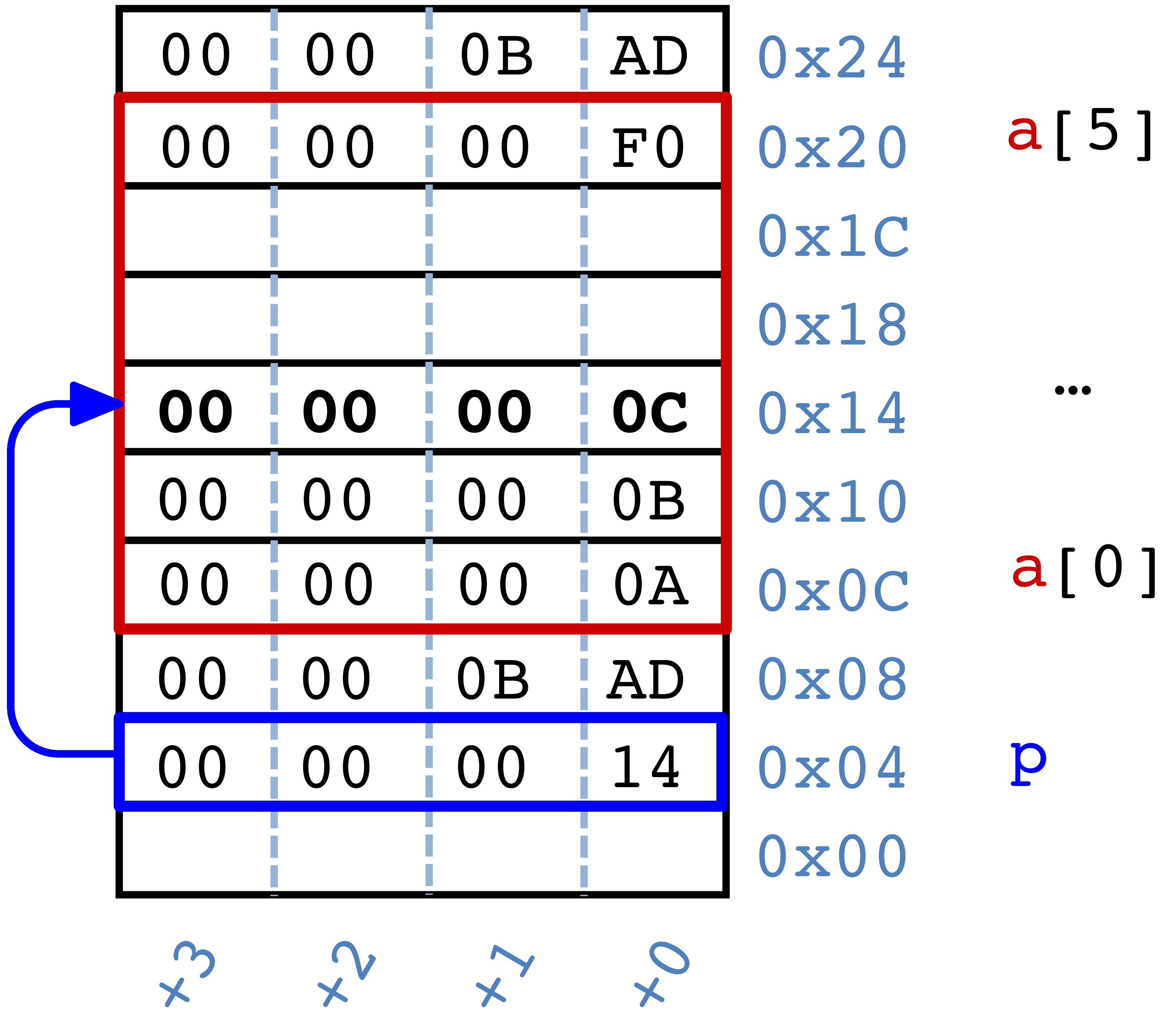
Both are scaled by the size of the type.

```
*p = a[1] + 1;
```



+3	+2	+1	+0		
00	00	0B	AD	0x24	
00	00	00	F0	0x20	a[5]
				0x1C	
				0x18	
00	**00**	**00**	**0C**	0x14	…
00	00	00	0B	0x10	
00	00	00	0A	0x0C	a[0]
00	00	0B	AD	0x08	
00	00	00	14	0x04	p
				0x00	

Assume p has type int *. Are `p[2] = 5` and `*(p + 2) = 5` equivalent? What about `p[2] = 5` and `*p + 2 = 5`?

No; No.

No; Yes.

Yes; No.

Yes; Yes.

C: Array allocation

Basic Principle

T *A*[*N*];

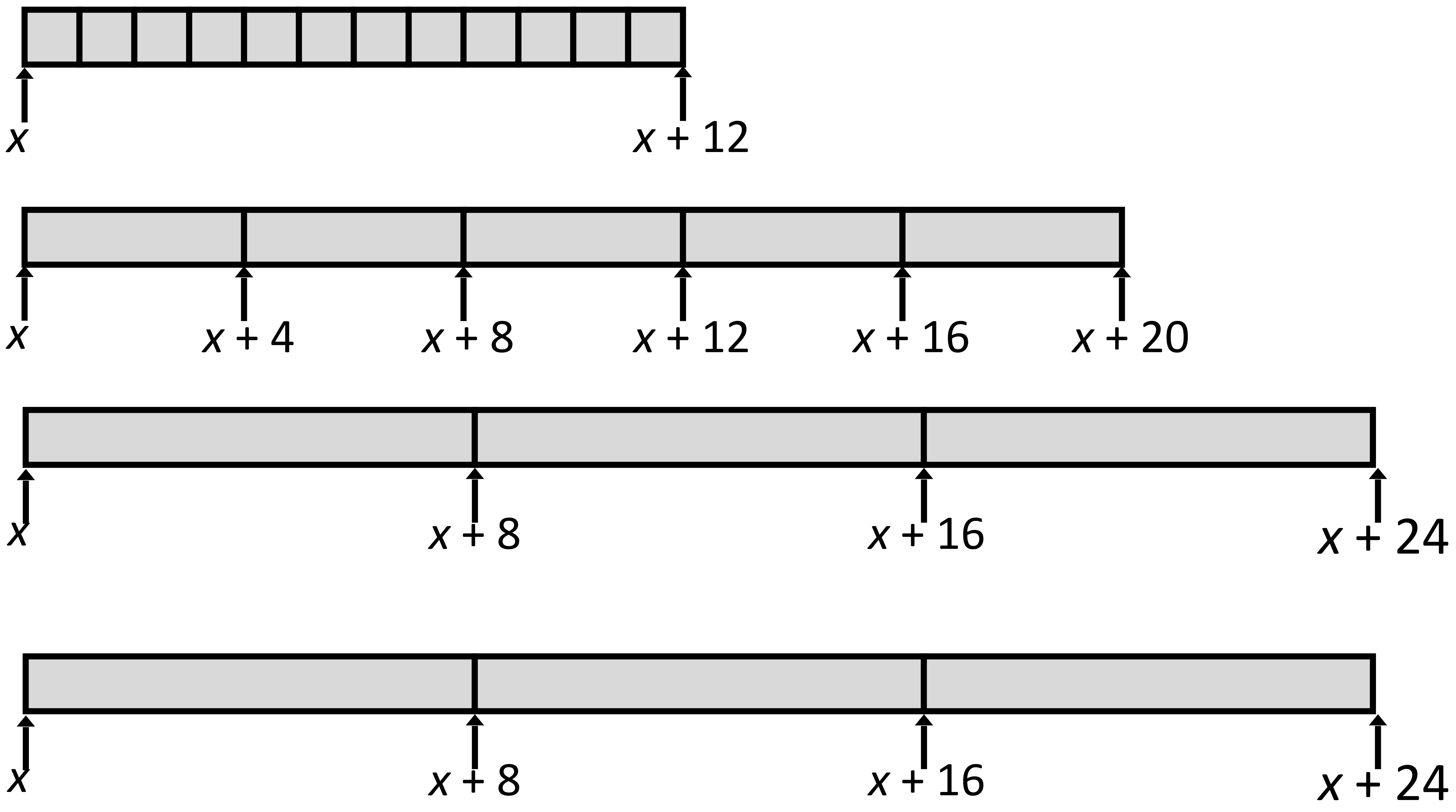
Array of length *N* with elements of type *T* and name *A*

Contiguous block of *N**sizeof(*T*) bytes of memory

Use *sizeof* to determine proper size in C.

```
char string[12];
int val[5];
double a[3];
char* p[3];
```

(or `char *p[3];`)



x86-64

size depends on the machine word size

C: Array access

ex

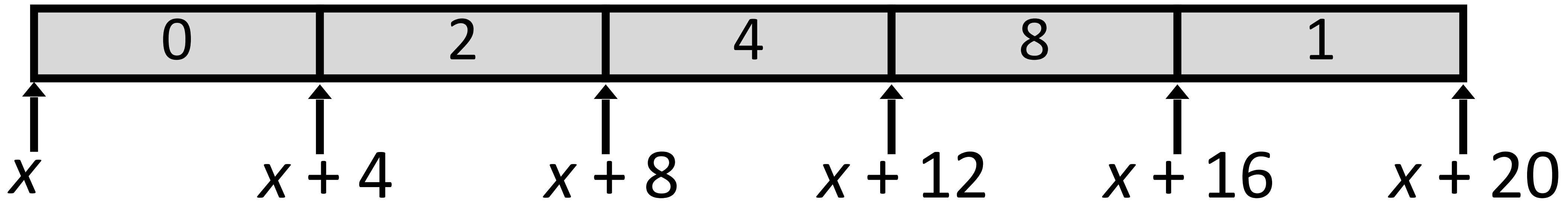
Basic Principle

T *A*[*N*];

Array of length *N* with elements of type *T* and name *A*

Identifier *A* has type *T**

`int val[5];`



Expression	Type	Value
`val[4]`	`int`	`1`
`val`	`int *`	
`val+1`	`int *`	
`&val[2]`	`int *`	
`val[5]`	`int`	
`*(val+1)`	`int`	
`val + i`	`int *`	

Representing strings

A C-style string is represented by an array of bytes (`char`).

- Elements are one-byte ASCII codes for each character.
- ASCII = American Standard Code for Information Interchange

32	space	48	0	64	@	80	P	96	`	112	p
33	!	49	1	65	A	81	Q	97	a	113	q
34	”	50	2	66	B	82	R	98	b	114	r
35	#	51	3	67	C	83	S	99	c	115	s
36	$	52	4	68	D	84	T	100	d	116	t
37	%	53	5	69	E	85	U	101	e	117	u
38	&	54	6	70	F	86	V	102	f	118	v
39	’	55	7	71	G	87	W	103	g	119	w
40	(	56	8	72	H	88	X	104	h	120	x
41	)	57	9	73	I	89	Y	105	I	121	y
42	*	58	:	74	J	90	Z	106	j	122	z
43	+	59	;	75	K	91	[	107	k	123	{
44	,	60	<	76	L	92	\	108	l	124	|
45	-	61	=	77	M	93	]	109	m	125	}
46	.	62	>	78	N	94	^	110	n	126	~
47	/	63	?	79	O	95	_	111	o	127	del

C: Null-terminated strings

ex

C strings: arrays of ASCII characters ending with *null character*.

Why?

0x57	0x65	0x6C	0x6C	0x65	0x73	0x6C	0x65	0x79	0x20	0x43	0x53	0x00
'W'	'e'	'l'	'l'	'e'	's'	'l'	'e'	'y'	' '	'C'	'S'	'\0'

Does Endianness matter for strings?

```
int string_length(char str[]) {







}
```

C: * *and* [ ]

ex

C programmers often use * where you might expect []:

e.g., char*:

- pointer to a char
- pointer to the first char in a string of unknown length

```
int strcmp(char* a, char* b);
```

C: 0 vs. '\0' vs. NULL

0

Name:	zero
Type:	`int`
Size:	4 bytes
Value:	`0x00000000`
Usage:	The integer zero.

'\0'

Name:	null character
Type:	`char`
Size:	1 byte
Value:	`0x00`
Usage:	Terminator for C strings.

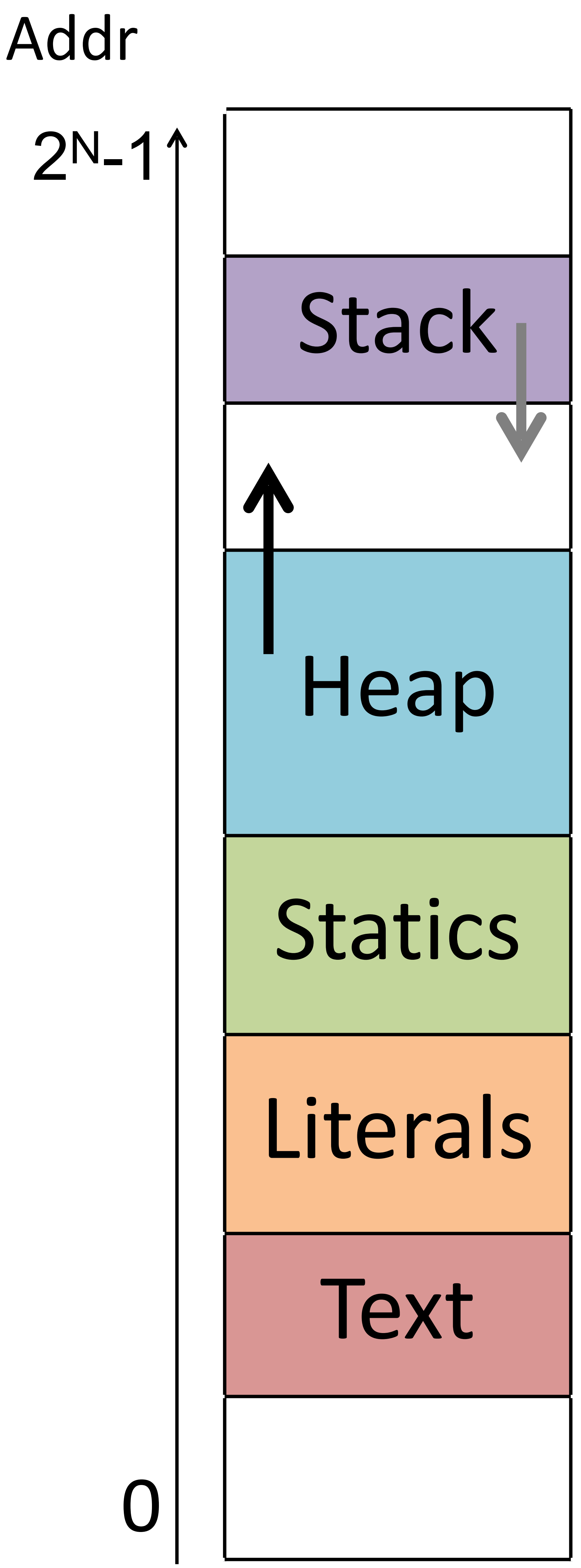
NULL

Name:	null pointer / null reference / null address
Type:	`void*`
Size:	1 word (= 8 bytes on a 64-bit architecture)
Value:	`0x0000000000000000`
Usage:	The absence of a pointer where one is expected.
	Address 0 is inaccessible, so `*NULL` is invalid; it crashes.

Is it important/necessary to encode the null character or the null pointer as 0x0?

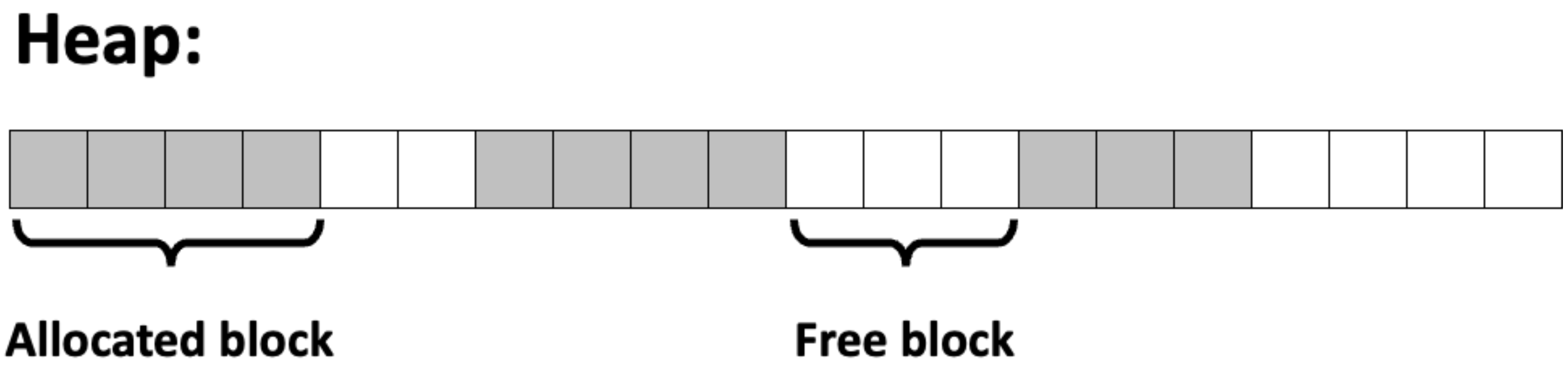
What happens if a programmer mixes up these "zeroey" values?

Memory address-space layout



	Perm	Contents	Managed by	Initialized
Stack	RW	Procedure context	Compiler	Run time
Heap	RW	Dynamic data structures	Programmer, malloc/free, new/GC	Run time
Statics	RW	Global variables/ static data structures	Compiler/ Assembler/Linker	Startup
Literals	R	String literals	Compiler/ Assembler/Linker	Startup
Text	X	Instructions	Compiler/ Assembler/Linker	Startup

C: Dynamic memory allocation in the heap



Managed by memory allocator:

pointer to newly allocated block of at least that size

number of contiguous bytes required

```
void* malloc(size_t size);

void free(void* ptr);
```

pointer to allocated block to free

C: standard memory allocator

```
#include <stdlib.h>     // include C standard library
```

```
void* malloc(size_t size)
```

Allocates a memory block of at least `size` bytes and returns its address.

If memory error (e.g., allocator has no space left), returns `NULL`.

Rules:

Check for error result.

Cast result to relevant pointer type.

Use sizeof(...) to determine size.

```
void free(void* ptr)
```

Deallocates the block referenced by `ptr`,
making its space available for new allocations.

`ptr` **must** be a `malloc` result that has not yet been `freed`.

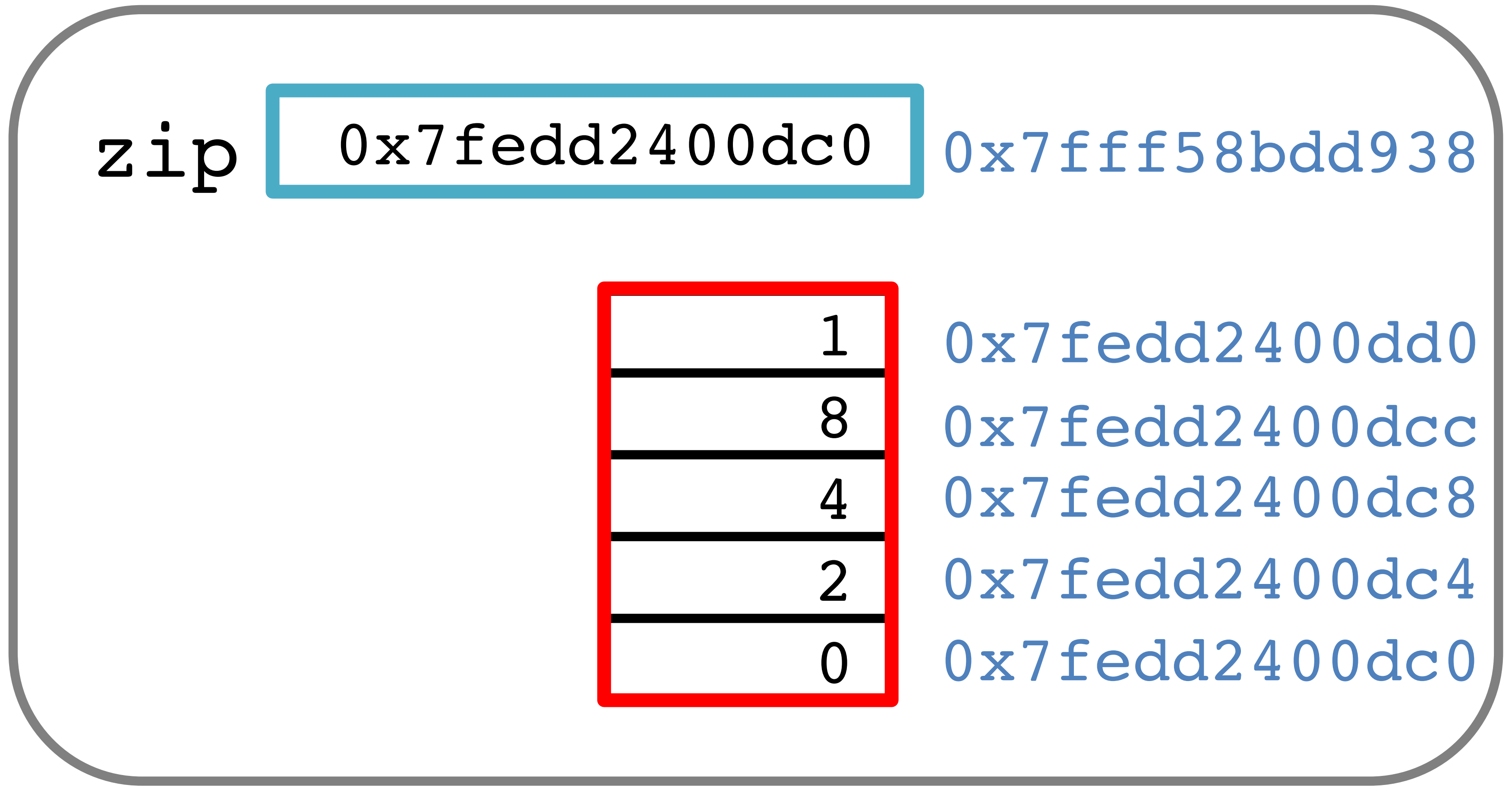
Rules:

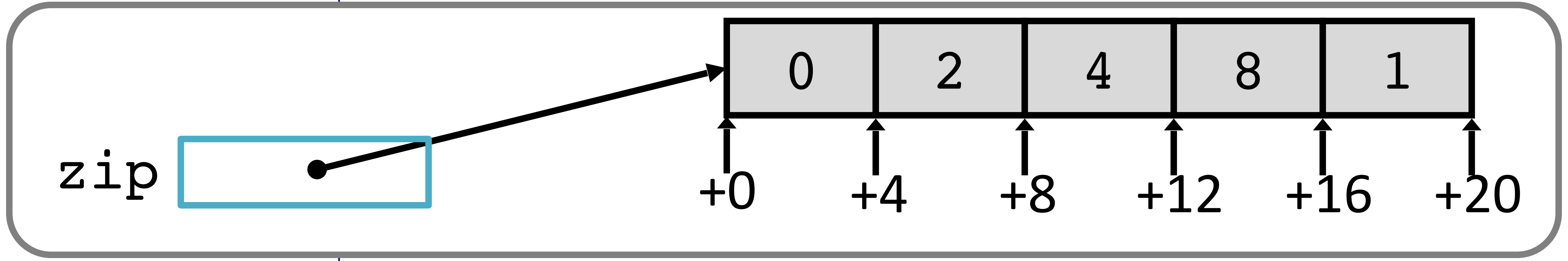
`ptr` must be a `malloc` result that has not yet been freed.

Do not use `*ptr` after freeing.

C: Dynamic array allocation

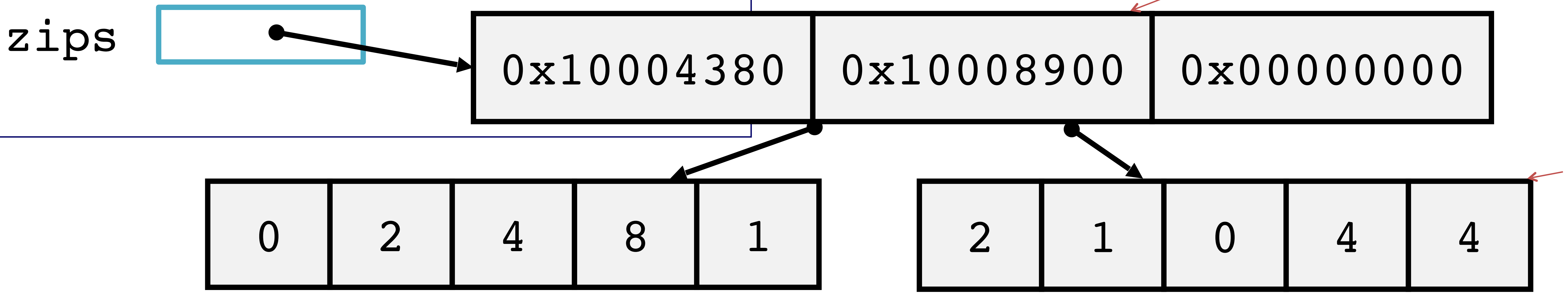
```
#define ZIP_LENGTH 5
int* zip = (int*)malloc(sizeof(int)*ZIP_LENGTH);
if (zip == NULL) {    // if error occurred
  perror("malloc");   // print error message
  exit(0);            // end the program
}

zip[0] = 0;
zip[1] = 2;
zip[2] = 4;
zip[3] = 8;
zip[4] = 1;

printf("zip is");
for (int i = 0; i < ZIP_LENGTH; i++) {
   printf(" %d", zip[i]);
}
printf("\n");

free(zip);
```





C: Array of pointers to arrays of `ints`

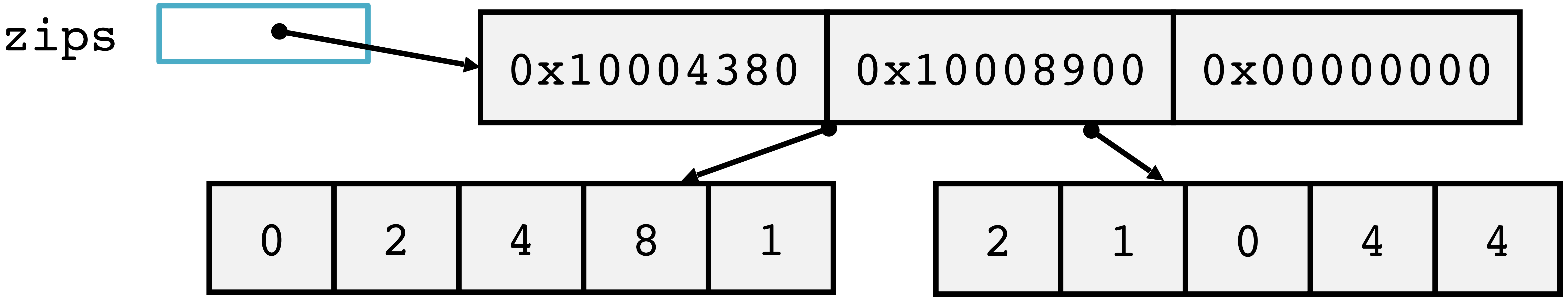
```
int** zips = (int**)malloc(sizeof(int*) * 3);

zips[0] = (int*)malloc(sizeof(int)*5);
int* zip0 = zips[0];
zip0[0] = 0;
zips[0][1] = 2;
zips[0][2] = 4;
zips[0][3] = 8;
zips[0][4] = 1;

zips[1] = (int*)malloc(sizeof(int)*5);
zips[1][0] = 2;
zips[1][1] = 1;
zips[1][2] = 0;
zips[1][3] = 4;
zips[1][4] = 4;

zips[2] = NULL;
```



Why terminate with NULL?

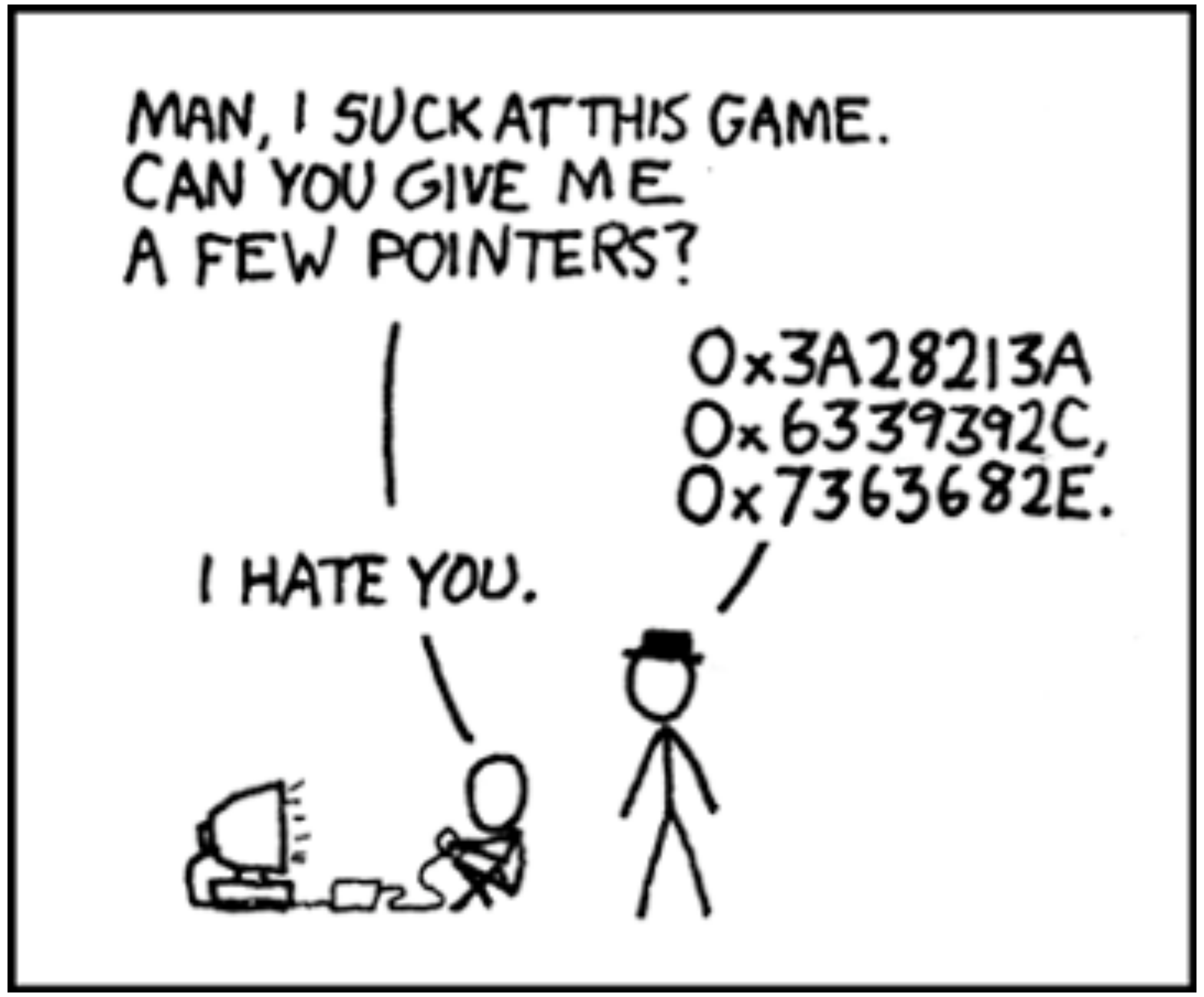
Why no NULL?

Zip code



```
// return a count of all zips that end with digit endNum
int zipCount(int* zips[], int endNum) {











}
```

http://xkcd.com/138/



C: `scanf` reads formatted input

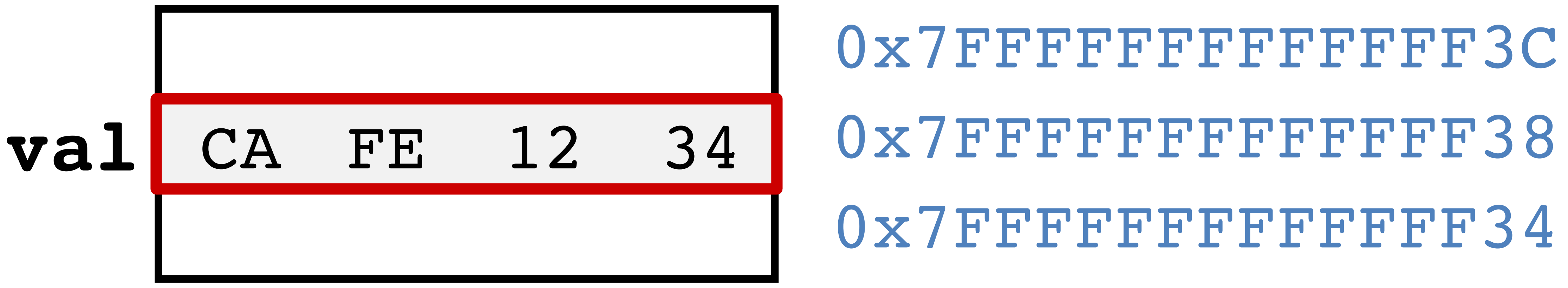
```
int val;

...

scanf("%d", &val);
```

Declared, but not initialized.
Holds anything.

Read one `int`
in **decimal**$_{10}$ format
from input.

Store it in memory
at this address.

Store in memory at the address given by the **address of `val`**: store input @ `0x7F…F38`.



```
              0x7FFFFFFFFFFFFF3C
val  CA FE 12 34   0x7FFFFFFFFFFFFF38
              0x7FFFFFFFFFFFFF34
```

C: Classic bug using `scanf`

!!!

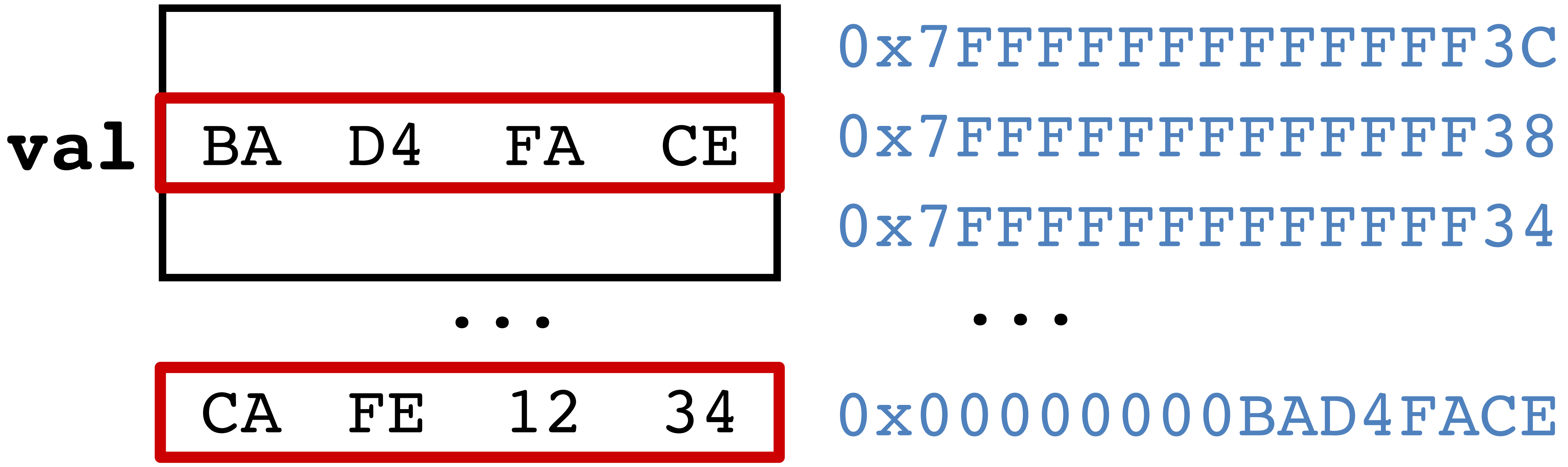
```
int val;

...

scanf("%d", val);
```

Declared, but not initialized. Holds anything.

Read one `int` in decimal$_{10}$ format **from input.**

Store it in memory at this address.

Store in memory at the address given by the **contents of `val`** (implicitly cast as a pointer): store input @ `0xBAD4FACE`.



```
       |             |   0x7FFFFFFFFFFFFF3C
val    | BA D4 FA CE |   0x7FFFFFFFFFFFFF38
       |             |   0x7FFFFFFFFFFFFF34
            ...               ...
       | CA FE 12 34 |   0x00000000BAD4FACE
```

Best case: 🤦! **crash immediately** with segmentation fault/bus error.

Bad case: 🤬 **silently corrupt data** stored @ `0xBAD4FACE`, **fail to store input in `val`**, *and keep going.*

Worst case: 💻🔥🧨🚀 program does literally anything.

C: Memory error messages

11: **segmentation fault** ("**segfault**", SIGSEGV)
accessing address outside legal area of memory

10: **bus error** (SIGBUS)
accessing misaligned or other problematic address

More to come on debugging!

OKAY, HUMAN.
HUH?
BEFORE YOU HIT 'COMPILE,' LISTEN UP.
YOU KNOW WHEN YOU'RE FALLING ASLEEP, AND YOU IMAGINE YOURSELF WALKING OR SOMETHING,
AND SUDDENLY YOU MISSTEP, STUMBLE, AND JOLT AWAKE?
YEAH!
WELL, THAT'S WHAT A SEGFAULT FEELS LIKE.
DOUBLE-CHECK YOUR DAMN POINTERS, OKAY?

http://xkcd.com/371/

C: Why?

Why learn C?

- Think like actual computer (abstraction close to machine level) without dealing with machine code.
- Understand just how much Your Favorite Language provides.
- Understand just how much Your Favorite Language might cost.
- Classic.
- Still (more) widely used (than it should be).
- Pitfalls still fuel devastating reliability and security failures today.

Why not use C?

- Probably not the right language for your next personal project.
- It "gets out of the programmer's way" … even when the programmer is unwittingly running toward a cliff.
- Advances in programming language design since the 70's have produced languages that fix C's problems while keeping strengths.