



x86: Procedures and the Call Stack

The call stack discipline
x86 procedure call and return instructions
x86 calling conventions
x86 register-saving conventions

<https://cs.wellesley.edu/~cs240/>

x86: Procedures and the Call Stack

Outline

1. Motivation
 - a. (video 1) What we have seen so far
 - b. (video 1) Why we can't implement procedure calls with jumps alone
2. (video 1) High-level call stack example
3. Procedure control flow instructions: call and ret
4. Procedure call example (in depth!) on whiteboard
5. Caller vs/callee example
6. (*Covered in lab, video*) Recursion example

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Does the call stack really matter?

Alexa VanHattum 9:31 AM Yesterday!
looking at popcnt now!

```
avh veri/veri - (popcnt-expand)$ ./script/veri.sh --rule popcnt_8
    Blocking waiting for file lock on build directory
    Compiling cranelift-isle-veri v0.1.0
    Writing generated file:
    /var/folders/9r/4bqb01xs6f08kpv59bk68cpc000gn/T/tmp.AsPxX054/clif_lower.is
    le
#990    popcnt_8

thread 'main' has overflowed its stack
fatal runtime error: stack overflow
```

mbm 9:36 AM
There's an environment variable you can set to increase stack size. The default is not that big.

Alexa VanHattum 10:49 AM And process with a loop instead of recursion

Yes, the call stack comes up in large-scale software engineering/research!

Why procedures?

Why functions? Why methods?

```
int contains_char(char* haystack, char needle) {
    while (*haystack != '\0') {
        if (*haystack == needle) return 1;
        haystack++;
    }
    return 0;
}
```

Answer: procedural abstraction

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Implementing procedures

1. How does a caller pass arguments to a procedure?

Have we already seen
how this is done?



2. How does a caller receive a return value from a procedure?



3. How does a procedure know where to return
(what code to execute next when done)?



4. Where does a procedure store local variables?



5. How do procedures share limited registers and memory?

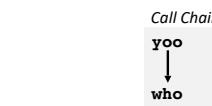
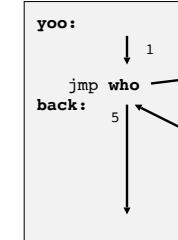


Procedure call/return: Jump?

```
yoo(...) {  
    • • •  
    who();  
    • • •  
}
```

```
who(...) {  
    • • •  
    • • •  
    • • •  
}
```

```
ru(...){  
    • • •  
}
```



But what if we want to call a function from multiple places in the code?

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Procedure call/return: Jump? Broken!

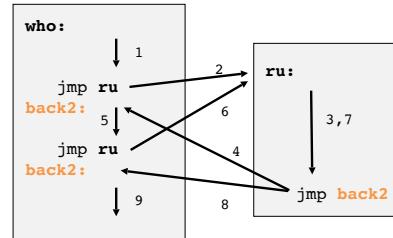
```
yoo(...) {  
    • • •  
    who();  
    • • •  
}
```

```
who(...) {  
    • • •  
    ru();  
    • • •  
    ru();  
    • • •  
}
```

```
ru(...){  
    • • •  
}
```

Call Chain

```
yoo  
↓  
who  
↓  
ru  
↓  
ru
```



But what if we want to call a function from multiple places in the code?

Broken: needs to track context.

Implementing procedures

requires **separate storage per call!**
(not just per procedure)

Have we already seen
how this is done?

1. How does a caller pass arguments to a procedure?



2. How does a caller receive a return value from a procedure?



3. How does a procedure know where to return
(what code to execute next when done)?



4. Where does a procedure store local variables?



1. How do procedures share limited registers and memory?



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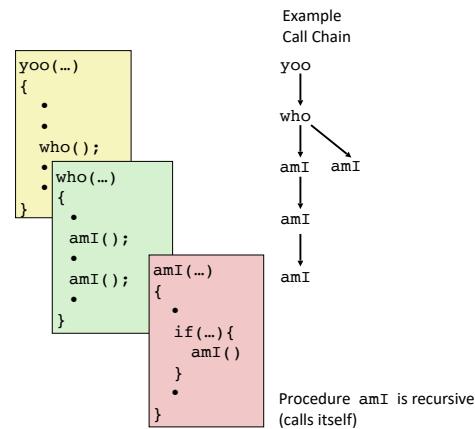
Memory Layout

reminder

Addr	Perm	Contents	Managed by	Initialized
2 ^{N-1} ↑	RW	Procedure context	Compiler	Run-time
	RW	Dynamic data structures	Programmer, malloc/free, new/ GC	Run-time
	RW	Global variables/ static data structures	Compiler/ Assembler/Linker	Startup
	R	String literals	Compiler/ Assembler/Linker	Startup
0	X	Instructions	Compiler/ Assembler/Linker	Startup

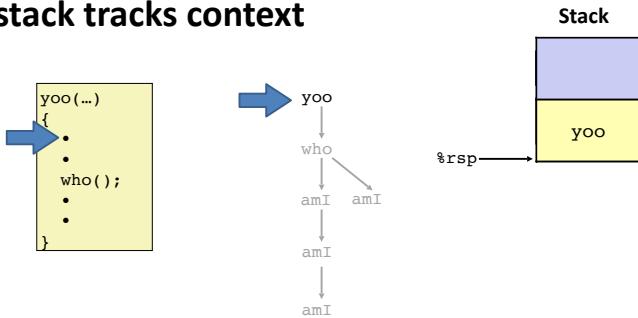
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Call stack tracks context

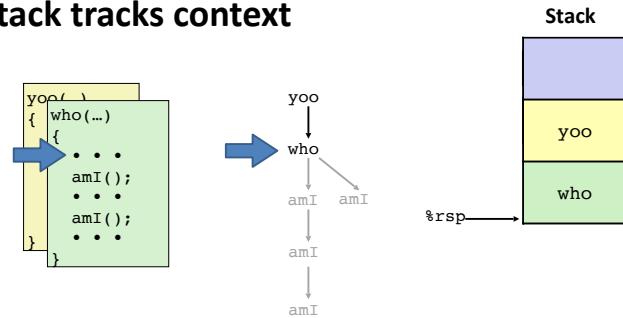


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Call stack tracks context



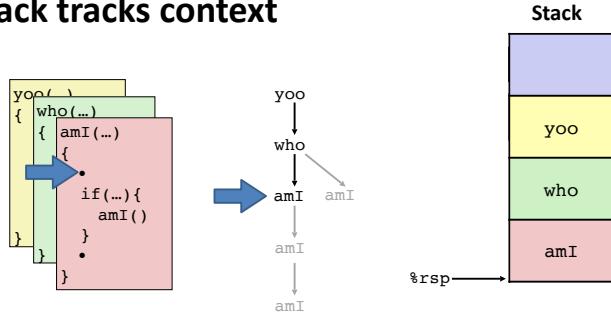
Call stack tracks context



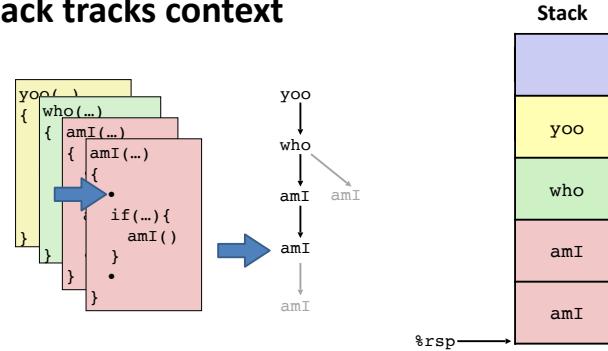
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Call stack tracks context



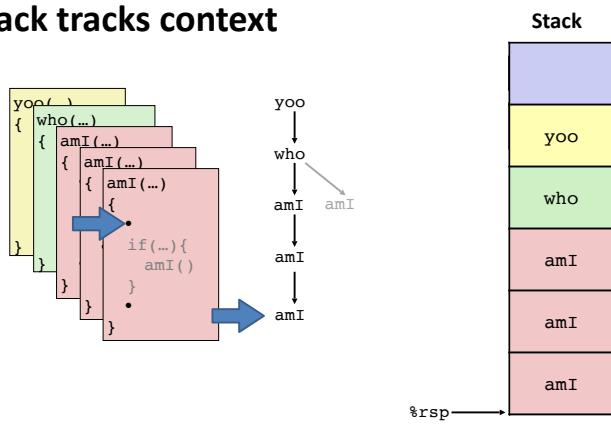
Call stack tracks context



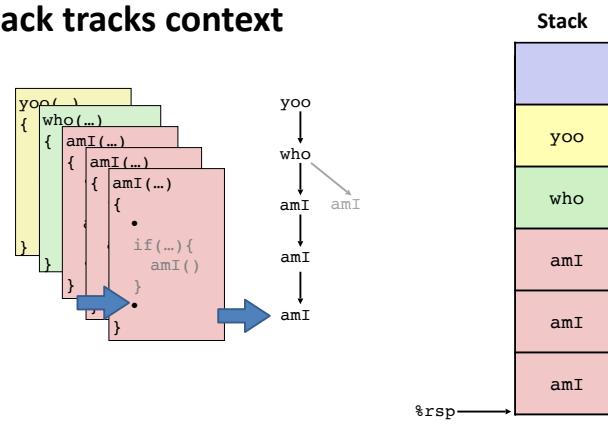
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Call stack tracks context



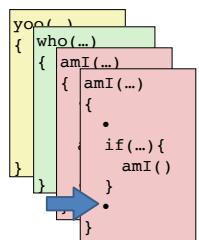
Call stack tracks context



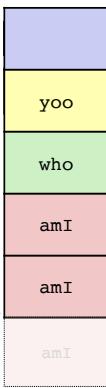
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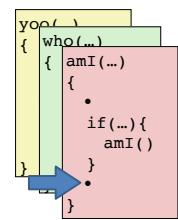
Call stack tracks context



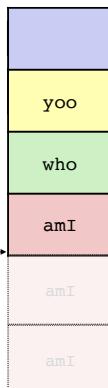
Stack



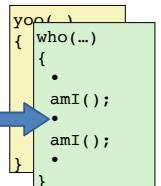
Call stack tracks context



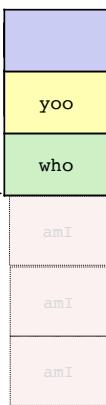
Stack



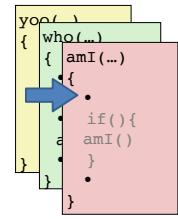
Call stack tracks context



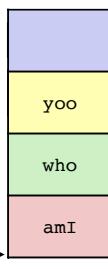
Stack



Call stack tracks context



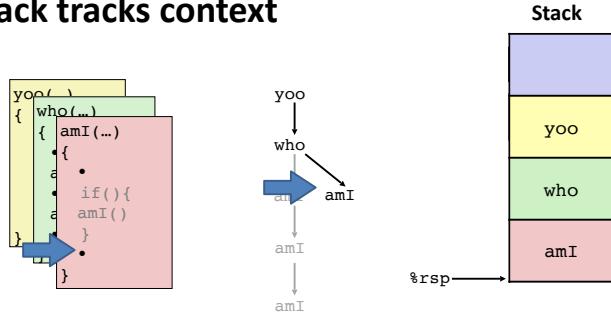
Stack



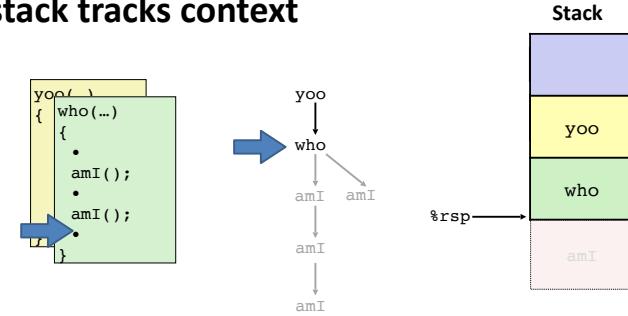
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Call stack tracks context



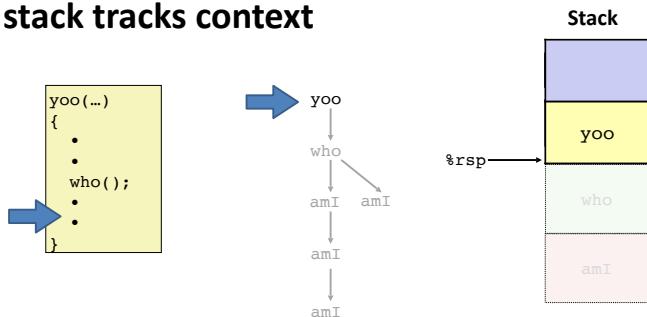
Call stack tracks context



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Call stack tracks context



In x86, where is the return address located for a "call" instruction?

In the instructions of the caller function

In the instructions of the callee (being called) function

On the call stack, once per function

On the call stack, once per call

None of the above

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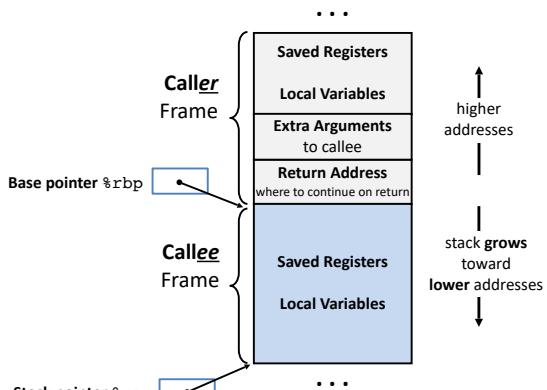
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The call stack supports procedures

Stack frame: section of stack used by one procedure `call` to store context while running.

Procedure code manages stack frames explicitly.

- **Setup:** allocate space at start of procedure.
- **Cleanup:** deallocate space before return.



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Procedure control flow instructions

Procedure call: `callq target`

1. Push return address on stack
2. Jump to `target`

Return address: Address of instruction after `call`.

```
400544: callq 400550 <mult2>
400549: movq %rax,(%rbx)
```

Procedure return: `retq`

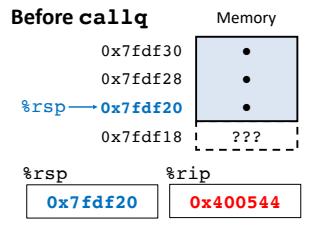
1. Pop return address from stack
2. Jump to return address

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Call example



```
0000000000400540 <multstore>:
.
400544: callq 400550 <mult2>
400549: mov %rax,(%rbx)
.
0000000000400550 <mult2>:
400550: mov %rdi,%rax
.
400557: retq
```



callq target

1. Push return address on stack
2. Jump to `target`

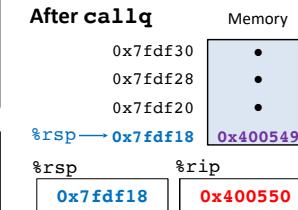
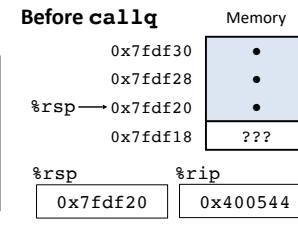
Call example



```
0000000000400540 <multstore>:
.
400544: callq 400550 <mult2>
400549: mov %rax,(%rbx)
.
0000000000400550 <mult2>:
400550: mov %rdi,%rax
.
400557: retq
```

callq target

1. Push return address on stack
2. Jump to `target`



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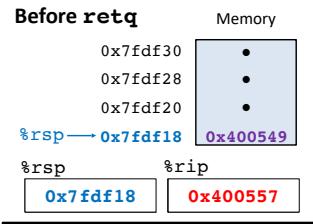
Return example



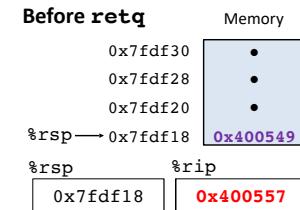
```
0000000000400540 <multstore>:
.
.
400544: callq 400550 <mult2>
400549: mov %rax,(%rbx)
.
.
```

```
0000000000400550 <mult2>:
400550: mov %rdi,%rax
.
.
400557: retq
```

retq
1. Pop return address from stack
2. Jump to return address



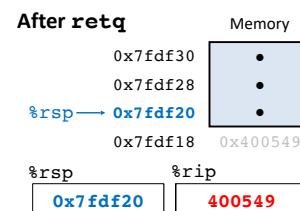
Return example



```
0000000000400540 <multstore>:
.
.
400544: callq 400550 <mult2>
400549: mov %rax,(%rbx)
.
.
```

```
0000000000400550 <mult2>:
400550: mov %rdi,%rax
.
.
400557: retq
```

retq
1. Pop return address from stack
2. Jump to return address



Procedure data flow conventions

Recall:

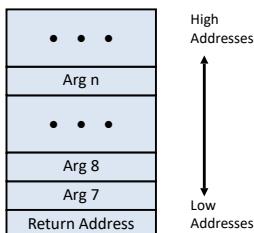
First 6 arguments: passed in registers

Arg 1	%rdi	Diane's
Arg 2	%rsi	Silk
Arg 3	%rdx	Dress
Arg 4	%rcx	Costs
Arg 5	%r8	\$8
Arg 6	%r9	9

Return value: passed in %rax

%rax

Remaining arguments:
passed on **stack** (in memory)



Procedure call / stack frame example

```
step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq
```

```
long step_up() {
    long v1 = 240;
    long v2 = increment(&v1, 61);
    return v1+v2;
}
```

Passes address of local variable (in stack).

```
increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq
```

```
long increment(long* p, long val) {
    long x = *p;
    long y = x + val;
    *p = y;
    return x;
}
```

Uses memory through pointer.

Procedure call example (step 0)

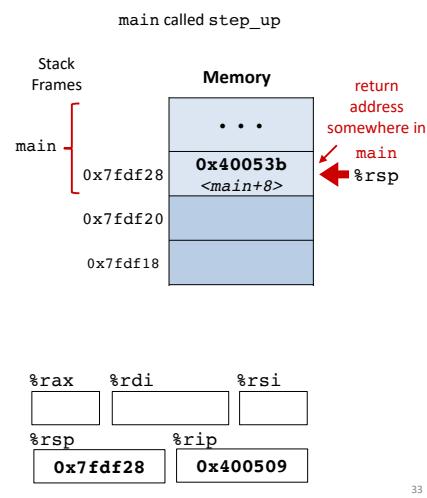
```

long step_up() {
    long v1 = 240;
    long v2 = increment(&v1, 61);
    return v1+v2;
}

step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq

increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq

```



Procedure call example (step 1)

```

long step_up() {
    long v1 = 240;
    long v2 = increment(&v1, 61);
    return v1+v2;
}

step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq

increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq

```

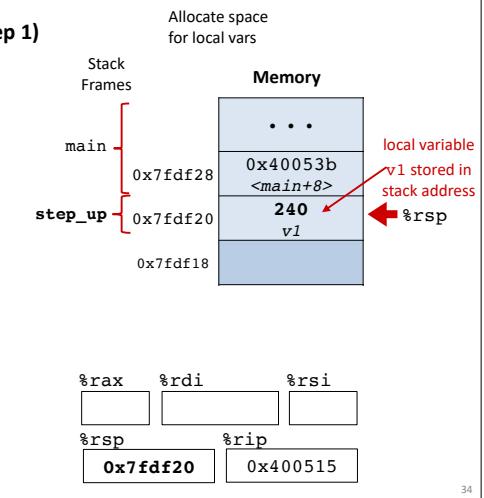
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long step_up() {
    long v1 = 240;
    long v2 = increment(&v1, 61);
    return v1+v2;
}

step_up:
400509: subq $8, %rsp
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400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq

increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq

```



Procedure call example (step 2)

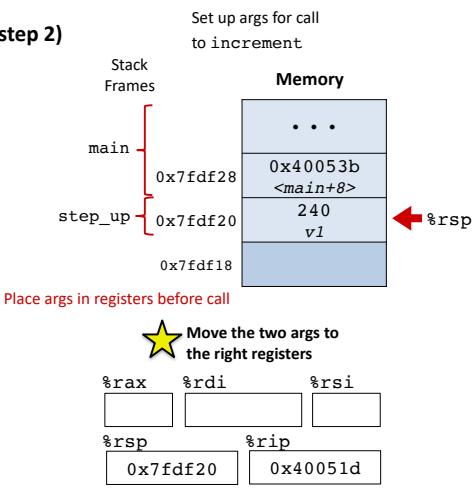
```

long step_up() {
    long v1 = 240;
    long v2 = increment(&v1, 61);
    return v1+v2;
}

step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq

increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq

```



Procedure call example (step 2)

```

long step_up() {
    long v1 = 240;
    long v2 = increment(&v1, 61);
    return v1+v2;
}

step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq

increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq

```

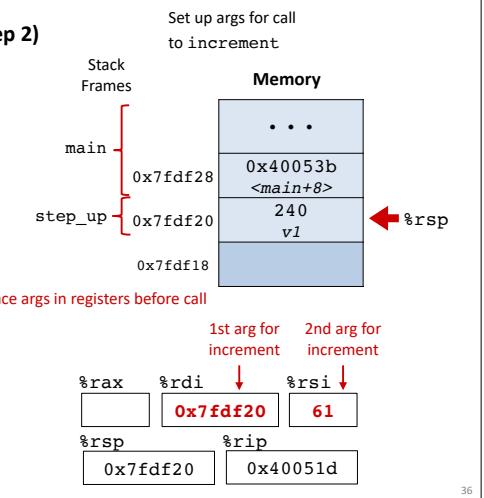
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    long v1 = 240;
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    return v1+v2;
}

step_up:
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400522: addq (%rsp), %rax
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4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq

```

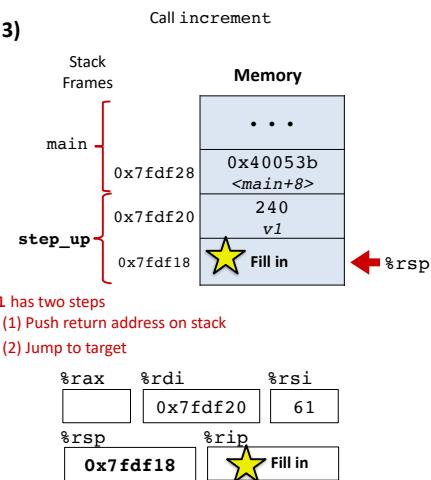


Procedure call example (step 3)

```
long step_up() {
    long v1 = 240;
    long v2 = increment(&v1, 61);
    return v1+v2;
}

step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq

increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq
```



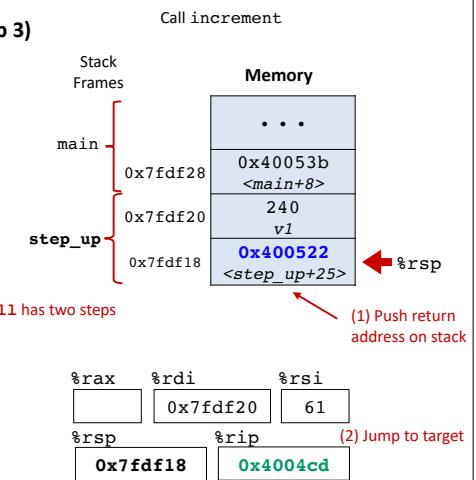
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Procedure call example (step 3)

```
long step_up() {
    long v1 = 240;
    long v2 = increment(&v1, 61);
    return v1+v2;
}

step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq

increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq
```



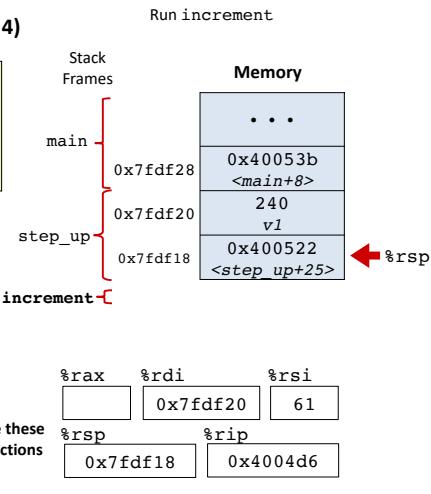
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Procedure call example (step 4)

```
long increment(long* p, long val) {
    long x = *p;
    long y = x + val;
    *p = y;
    return x;
}

step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq

increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq
```



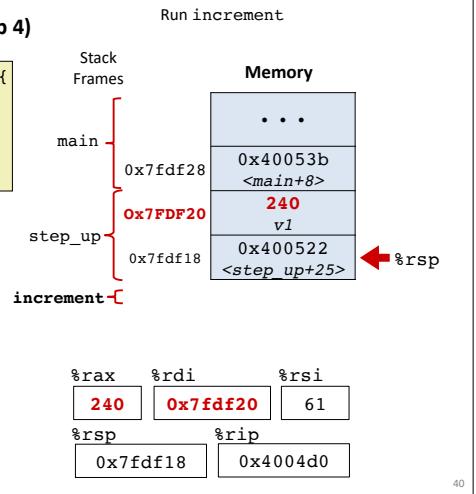
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Procedure call example (step 4)

```
long increment(long* p, long val) {
    long x = *p;
    long y = x + val;
    *p = y;
    return x;
}

step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq

increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq
```



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Procedure call example (step 4)

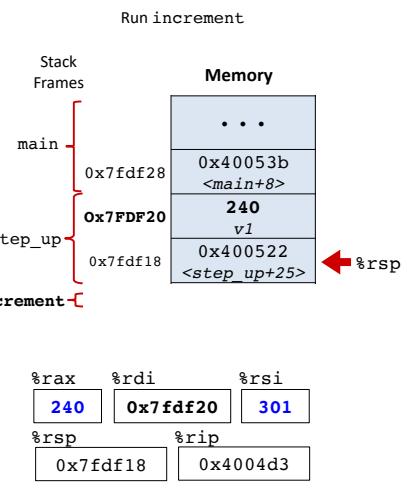
```

lo long increment(long* p, long val) {
    long x = *p;
    long y = x + val;
    *p = y;
    return x;
}

step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq

increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq

```



Procedure call example (step 4)

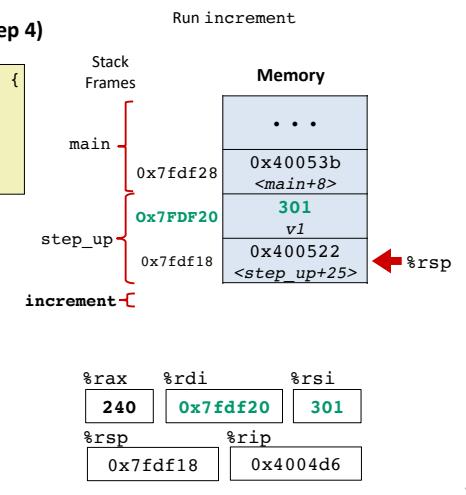
```

lo long increment(long* p, long val) {
    long x = *p;
    long y = x + val;
    *p = y;
    return x;
}

step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq

increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq

```



Procedure call example (step 5a)

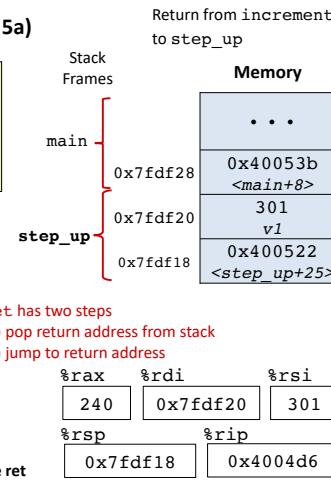
```

lo long increment(long* p, long val) {
    long x = *p;
    long y = x + val;
    *p = y;
    return x;
}

step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq

increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq

```



Procedure call example (step 5b)

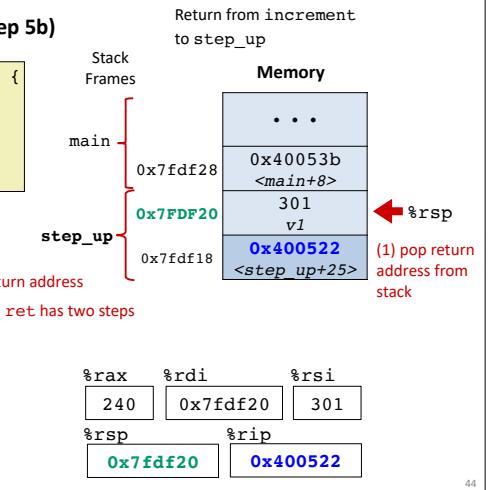
```

lo long increment(long* p, long val) {
    long x = *p;
    long y = x + val;
    *p = y;
    return x;
}

step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq

increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq

```

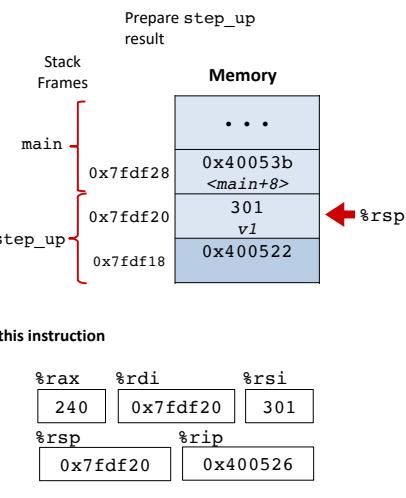


Procedure call example (step 6)

```
long step_up() {
    long v1 = 240;
    long v2 = increment(&v1, 61);
    return v1+v2;
}

step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq

increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq
```

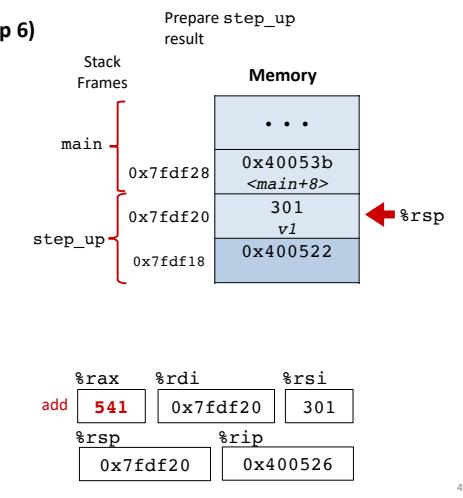


Procedure call example (step 6)

```
long step_up() {
    long v1 = 240;
    long v2 = increment(&v1, 61);
    return v1+v2;
}

step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq

increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq
```

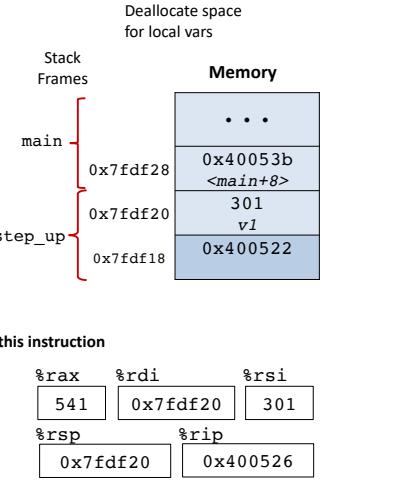


Procedure call example (step 7)

```
long step_up() {
    long v1 = 240;
    long v2 = increment(&v1, 61);
    return v1+v2;
}

step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq

increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq
```

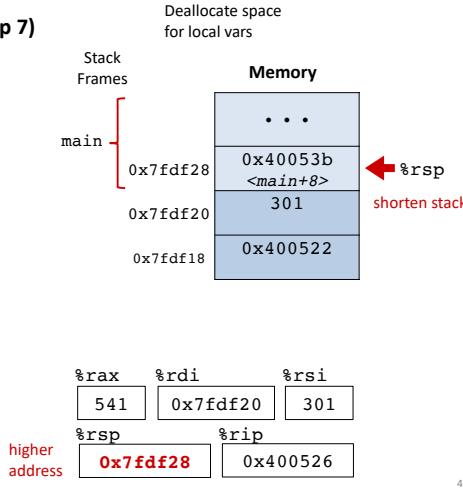


Procedure call example (step 7)

```
long step_up() {
    long v1 = 240;
    long v2 = increment(&v1, 61);
    return v1+v2;
}

step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq

increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq
```

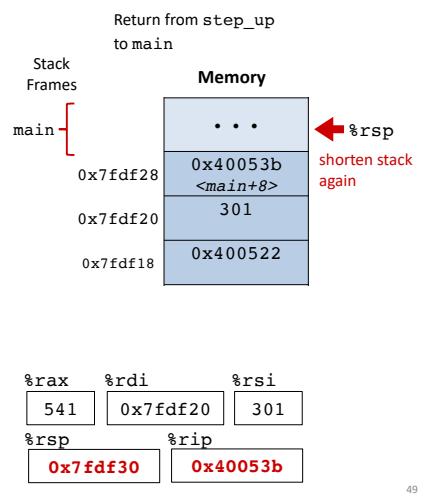


Procedure call example (step 8)

```
long step_up() {
    long v1 = 240;
    long v2 = increment(&v1, 61);
    return v1+v2;
}

step_up:
400509: subq $8, %rsp
40050d: movq $240, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq (%rsp), %rax
400526: addq $8, %rsp
40052a: retq

increment:
4004cd: movq (%rdi), %rax
4004d0: addq %rax, %rsi
4004d3: movq %rsi, (%rdi)
4004d6: retq
```



Implementing procedures

Have we now seen how this is done?

- How does a caller pass arguments to a procedure? ✓
- How does a caller receive a return value from a procedure? ✓
- How does a procedure know where to return (what code to execute next when done)? ✓
- Where does a procedure store local variables? ✓
- How do procedures share limited registers and memory? ??

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Register saving conventions

yoo calls who:

<i>Caller</i>	<i>Callee</i>
.....
who();
}	

Will register contents still be there after a procedure call?

```
yoo:
    ...
    movq $12345, %rbx
    call who
    addq %rbx, %rax
    ...
    ret
```

```
who:
    ...
    addq %rdi, %rbx
    ...
    ret
```

Conventions:
Caller Save
Callee Save

x86-64 register conventions

%rax	Return value – Caller saved	%r8	Argument #5 – Caller saved
%rbx	Callee saved	%r9	Argument #6 – Caller saved
%rcx	Argument #4 – Caller saved	%r10	Caller saved
%rdx	Argument #3 – Caller saved	%r11	Caller Saved
%rsi	Argument #2 – Caller saved	%r12	Callee saved
%rdi	Argument #1 – Caller saved	%r13	Callee saved
%rsp	Stack pointer	%r14	Callee saved
%rbp	Callee saved	%r15	Callee saved

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Callee-save example (step 0)

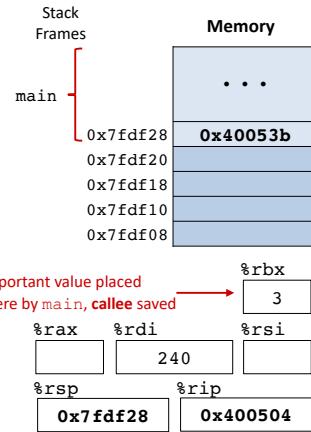
Similar function, but now takes an arg for the local variable

```
long step_by(long x) {
    long v1 = x;
    long v2 = increment(&v1, 61);
    return x + v2;
}
```

```
step_by:
400504: pushq %rbx
400506: movq %rdi, %rbx
400509: subq $16, %rsp
40050d: movq %rdi, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq %rbx, %rax
400525: addq $16, %rsp
400529: popq %rbx
40052b: retq
```

caller saved: %rax, %rdi, %rsi
callee saved: %rbx

main called step_by(240)



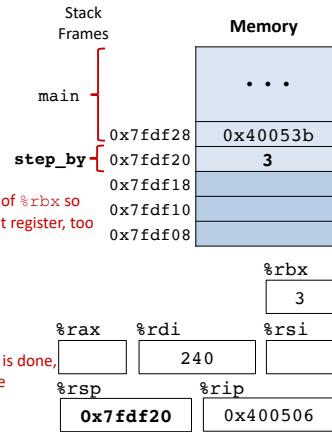
Callee-save example (step 1)

Save register %rbx

```
long step_by(long x) {
    long v1 = x;
    long v2 = increment(&v1, 61);
    return x + v2;
}
```

```
step_by:
400504: pushq %rbx ← Save the value of %rbx so we can use that register, too
400506: movq %rdi, %rbx
400509: subq $16, %rsp
40050d: movq %rdi, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq %rbx, %rax
400525: addq $16, %rsp
400529: popq %rbx ← Once this function is done, restore saved value
40052b: retq
```

caller saved: %rax, %rdi, %rsi
callee saved: %rbx



Callee-save example (step 2)

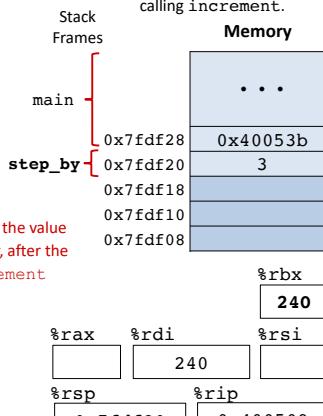
Copy argument x to %rbx for continued use after calling increment.

```
long step_by(long x) {
    long v1 = x;
    long v2 = increment(&v1, 61);
    return x + v2;
}
```

```
step_by:
400504: pushq %rbx
400506: movq %rdi, %rbx ← Need to save the value
400509: subq $16, %rsp
40050d: movq %rdi, (%rsp) x to use later, after the
400515: movq %rsp, %rdi call to increment
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq %rbx, %rax
400525: addq $16, %rsp
400529: popq %rbx
40052b: retq
```

caller saved: %rax, %rdi, %rsi
callee saved: %rbx

Stack Frames
main ...
step_by 0x7fdf20 3
0x7fdf18
0x7fdf10
0x7fdf08



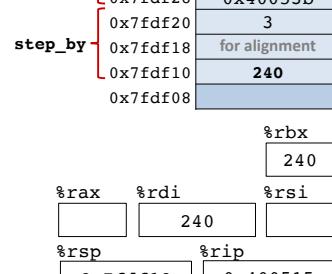
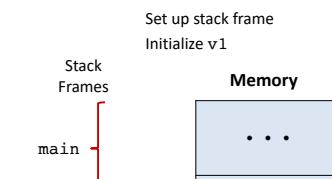
Callee-save example (step 3)

Set up stack frame
Initialize v1

```
long step_by(long x) {
    long v1 = x;
    long v2 = increment(&v1, 61);
    return x + v2;
}
```

```
step_by:
400504: pushq %rbx
400506: movq %rdi, %rbx
400509: subq $16, %rsp
40050d: movq %rdi, (%rsp) ← Convention: at call, %rsp must be a multiple of 16
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq %rbx, %rax
400525: addq $16, %rsp
400529: popq %rbx
40052b: retq
```

caller saved: %rax, %rdi, %rsi
callee saved: %rbx

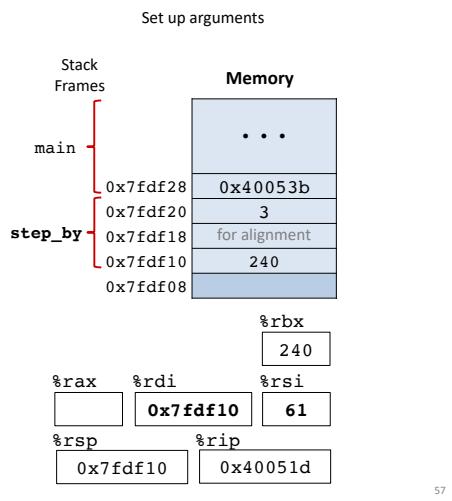


Callee-save example (step 4)

```
long step_by(long x) {
    long v1 = x;
    long v2 = increment(&v1, 61);
    return x + v2;
}

step_by:
400504: pushq %rbx
400506: movq %rdi, %rbx
400509: subq $16, %rsp
40050d: movq %rdi, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq %rbx, %rax
400525: addq $16, %rsp
400529: popq %rbx
40052b: retq

caller saved: %rax, %rdi, %rsi
callee saved: %rbx
```

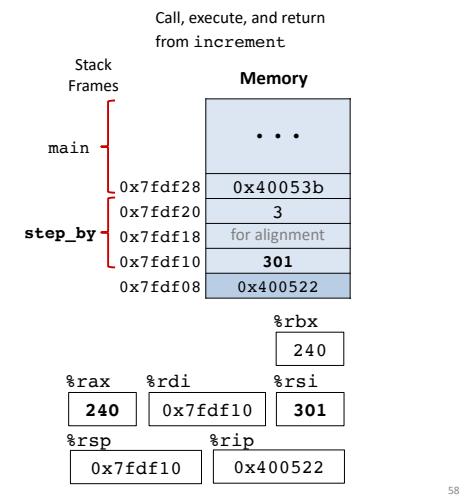


Callee-save example (step 5)

```
long step_by(long x) {
    long v1 = x;
    long v2 = increment(&v1, 61);
    return x + v2;
}

step_by:
400504: pushq %rbx
400506: movq %rdi, %rbx
400509: subq $16, %rsp
40050d: movq %rdi, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq %rbx, %rax
400525: addq $16, %rsp
400529: popq %rbx
40052b: retq

caller saved: %rax, %rdi, %rsi
callee saved: %rbx
```

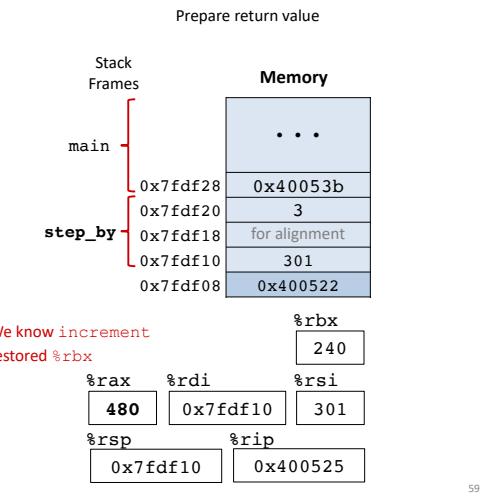


Callee-save example (step 6)

```
long step_by(long x) {
    long v1 = x;
    long v2 = increment(&v1, 61);
    return x + v2;
}

step_by:
400504: pushq %rbx
400506: movq %rdi, %rbx
400509: subq $16, %rsp
40050d: movq %rdi, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq %rbx, %rax
400525: addq $16, %rsp
400529: popq %rbx
40052b: retq

caller saved: %rax, %rdi, %rsi
callee saved: %rbx
```

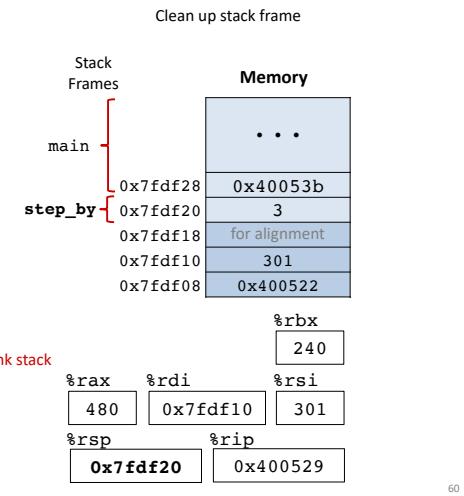


Callee-save example (step 7)

```
long step_by(long x) {
    long v1 = x;
    long v2 = increment(&v1, 61);
    return x + v2;
}

step_by:
400504: pushq %rbx
400506: movq %rdi, %rbx
400509: subq $16, %rsp
40050d: movq %rdi, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq %rbx, %rax
400525: addq $16, %rsp
400529: popq %rbx
40052b: retq

caller saved: %rax, %rdi, %rsi
callee saved: %rbx
```

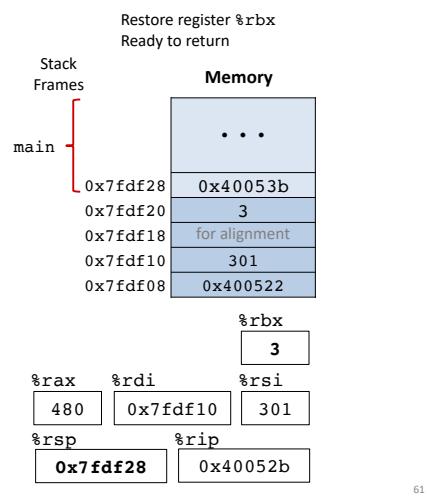


Callee-save example (step 8)

```
long step_by(long x) {
    long v1 = x;
    long v2 = increment(&v1, 61);
    return x + v2;
}

step_by:
400504: pushq %rbx
400506: movq %rdi, %rbx
400509: subq $16, %rsp
40050d: movq %rdi, (%rsp)
400515: movq %rsp, %rdi
400518: movl $61, %esi
40051d: callq 4004cd <increment>
400522: addq %rbx, %rax
400525: addq $16, %rsp
400529: popq %rbx           Restore %rbx ← for main
40052b: retq

caller saved: %rax, %rdi, %rsi
callee saved: %rbx
```



Recursion example: code

```
long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

```
pcount:
4005dd: movl $0, %eax
4005e2: testq %rdi, %rdi
4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
4005e8: movq %rdi, %rbx
4005eb: andl $1, %ebx
4005ee: shrq %rdi
4005f1: callq pcount
4005f6: addq %rbx, %rax
4005f9: popq %rbx
.L6:
4005fa: rep
4005fb: retq
```

base case/
condition

recursive
case

x&1 in %rbx
across call

save/restore

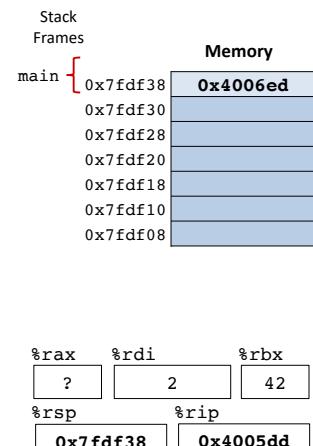
%rbx (callee-save)

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Recursion Example: pcount(2)

```
long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}

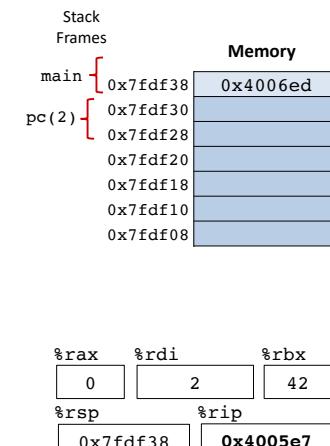
pcount:
4005dd: movl $0, %eax
4005e2: testq %rdi, %rdi
4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
4005e8: movq %rdi, %rbx
4005eb: andl $1, %ebx
4005ee: shrq %rdi
4005f1: callq pcount
4005f6: addq %rbx, %rax
4005f9: popq %rbx
.L6:
4005fa: rep
4005fb: retq
```



Recursion Example: pcount(2)

```
long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

```
pcount:
4005dd: movl $0, %eax
4005e2: testq %rdi, %rdi
4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
4005e8: movq %rdi, %rbx
4005eb: andl $1, %ebx
4005ee: shrq %rdi
4005f1: callq pcount
4005f6: addq %rbx, %rax
4005f9: popq %rbx
.L6:
4005fa: rep
4005fb: retq
```

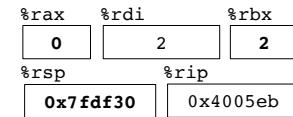
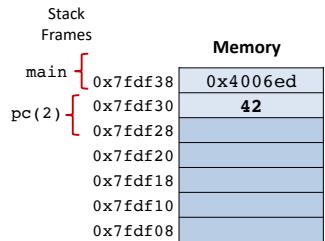


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Recursion Example: pcount(2)

```
long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

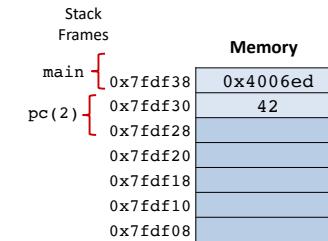
```
pcount:
4005dd: movl $0, %eax
4005e2: testq %rdi, %rdi
4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
4005e8: movq %rdi, %rbx
4005eb: andl $1, %ebx
4005ee: shrq %rdi
4005f1: callq pcount
4005f6: addq %rbx, %rax
4005f9: popq %rbx
.L6:
4005fa: rep
4005fb: retq
```



Recursion Example: pcount(2)

```
long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

```
pcount:
4005dd: movl $0, %eax
4005e2: testq %rdi, %rdi
4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
4005e8: movq %rdi, %rbx
4005eb: andl $1, %ebx
4005ee: shrq %rdi
4005f1: callq pcount
4005f6: addq %rbx, %rax
4005f9: popq %rbx
.L6:
4005fa: rep
4005fb: retq
```

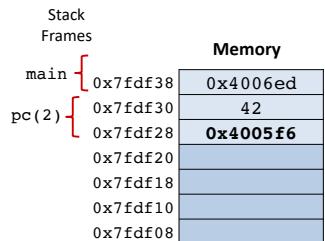


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Recursion Example: pcount(2) → pcount(1)

```
long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

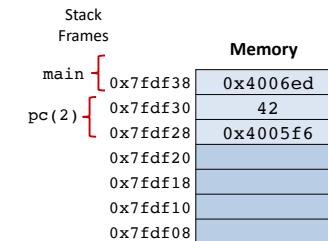
```
pcount:
4005dd: movl $0, %eax
4005e2: testq %rdi, %rdi
4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
4005e8: movq %rdi, %rbx
4005eb: andl $1, %ebx
4005ee: shrq %rdi
4005f1: callq pcount
4005f6: addq %rbx, %rax
4005f9: popq %rbx
.L6:
4005fa: rep
4005fb: retq
```



Recursion Example: pcount(2) → pcount(1)

```
long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

```
pcount:
4005dd: movl $0, %eax
4005e2: testq %rdi, %rdi
4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
4005e8: movq %rdi, %rbx
4005eb: andl $1, %ebx
4005ee: shrq %rdi
4005f1: callq pcount
4005f6: addq %rbx, %rax
4005f9: popq %rbx
.L6:
4005fa: rep
4005fb: retq
```

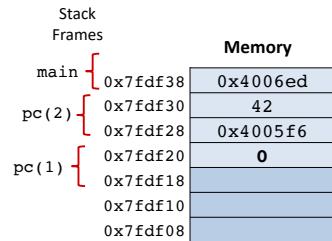


68

Recursion Example: pcount(2) → pcount(1)

```
1 long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

```
pcount:
4005dd: movl $0, %eax
4005e2: testq %rdi, %rdi
4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
4005e8: movq %rdi, %rbx
4005eb: andl $1, %ebx
4005ee: shrq %rdi
4005f1: callq pcount
4005f6: addq %rbx, %rax
4005f9: popq %rbx
.L6:
4005fa: rep
4005fb: retq
```

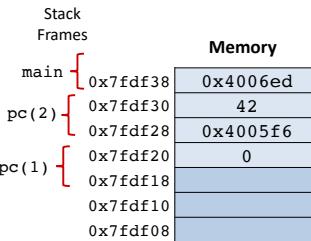


69

Recursion Example: pcount(2) → pcount(1)

```
1 long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

```
pcount:
4005dd: movl $0, %eax
4005e2: testq %rdi, %rdi
4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
4005e8: movq %rdi, %rbx
4005eb: andl $1, %ebx
4005ee: shrq %rdi
4005f1: callq pcount
4005f6: addq %rbx, %rax
4005f9: popq %rbx
.L6:
4005fa: rep
4005fb: retq
```

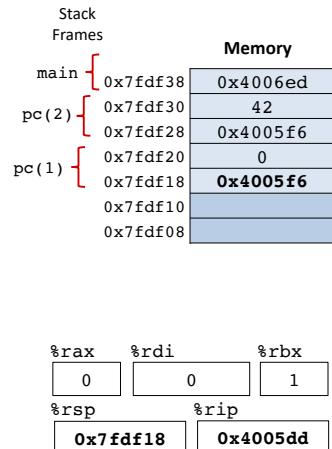


70

Recursion Example: pcount(2) → pcount(1) → pcount(0)

```
1 long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

```
pcount:
4005dd: movl $0, %eax
4005e2: testq %rdi, %rdi
4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
4005e8: movq %rdi, %rbx
4005eb: andl $1, %ebx
4005ee: shrq %rdi
4005f1: callq pcount
4005f6: addq %rbx, %rax
4005f9: popq %rbx
.L6:
4005fa: rep
4005fb: retq
```

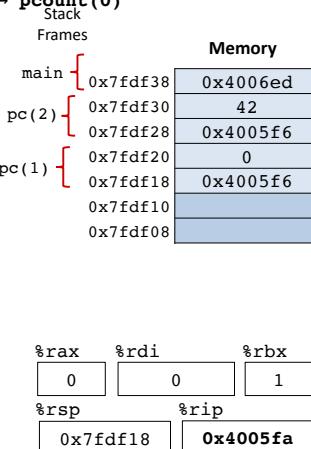


71

Recursion Example: pcount(2) → pcount(1) → pcount(0)

```
1 long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

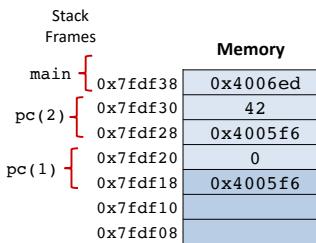
```
pcount:
4005dd: movl $0, %eax
4005e2: testq %rdi, %rdi
4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
4005e8: movq %rdi, %rbx
4005eb: andl $1, %ebx
4005ee: shrq %rdi
4005f1: callq pcount
4005f6: addq %rbx, %rax
4005f9: popq %rbx
.L6:
4005fa: rep
4005fb: retq
```



72

Recursion Example: pcount(2) → pcount(1) → pcount(0)

```
1 long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

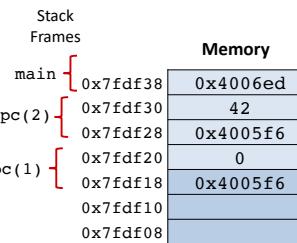


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Recursion Example: pcount(2) → pcount(1) → pcount(0)

```
1 long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

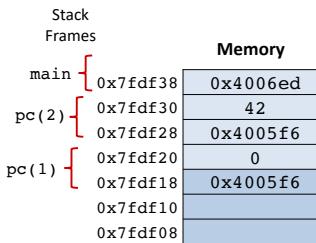
```
pcount:
4005dd: movl $0, %eax
4005e2: testq %rdi, %rdi
4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
4005e8: movq %rdi, %rbx
4005eb: andl $1, %ebx
4005ee: shrq %rdi
4005f1: callq pcount
4005f6: addq %rbx, %rax
4005f9: popq %rbx
.L6:
4005fa: rep
4005fb: retq
```



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Recursion Example: pcount(2) → pcount(1) → pcount(0)

```
1 long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

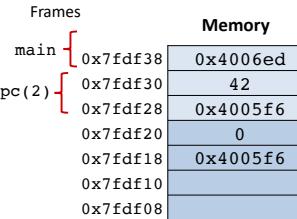


75

Recursion Example: pcount(2) → pcount(1) → pcount(0)

```
1 long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

```
pcount:
4005dd: movl $0, %eax
4005e2: testq %rdi, %rdi
4005e5: je 4005fa <.L6>
4005e7: pushq %rbx
4005e8: movq %rdi, %rbx
4005eb: andl $1, %ebx
4005ee: shrq %rdi
4005f1: callq pcount
4005f6: addq %rbx, %rax
4005f9: popq %rbx
.L6:
4005fa: rep
4005fb: retq
```

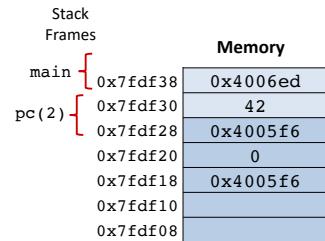


76

Recursion Example: pcount(2) → pcount(1) → pcount(0)

```
long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}

pcount:
    4005dd: movl $0, %eax
    4005e2: testq %rdi, %rdi
    4005e5: je 4005fa <.L6>
    4005e7: pushq %rbx
    4005e8: movq %rdi, %rbx
    4005eb: andl $1, %ebx
    4005ee: shrq %rdi
    4005f1: callq pcount
    4005f6: addq %rbx, %rax
    4005f9: popq %rbx
.L6:
    4005fa: rep
    4005fb: retq
```

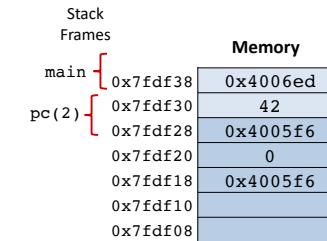


77

Recursion Example: pcount(2) → pcount(1) → pcount(0)

```
long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

```
pcount:
    4005dd: movl $0, %eax
    4005e2: testq %rdi, %rdi
    4005e5: je 4005fa <.L6>
    4005e7: pushq %rbx
    4005e8: movq %rdi, %rbx
    4005eb: andl $1, %ebx
    4005ee: shrq %rdi
    4005f1: callq pcount
    4005f6: addq %rbx, %rax
    4005f9: popq %rbx
.L6:
    4005fa: rep
    4005fb: retq
```

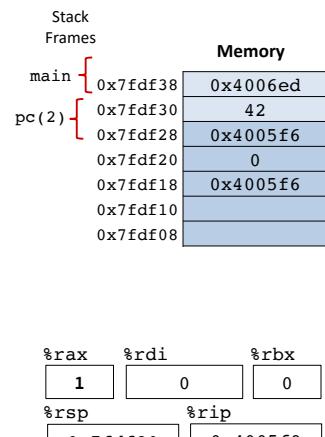


78

Recursion Example: pcount(2) → pcount(1) → pcount(0)

```
long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}

pcount:
    4005dd: movl $0, %eax
    4005e2: testq %rdi, %rdi
    4005e5: je 4005fa <.L6>
    4005e7: pushq %rbx
    4005e8: movq %rdi, %rbx
    4005eb: andl $1, %ebx
    4005ee: shrq %rdi
    4005f1: callq pcount
    4005f6: addq %rbx, %rax
    4005f9: popq %rbx
.L6:
    4005fa: rep
    4005fb: retq
```

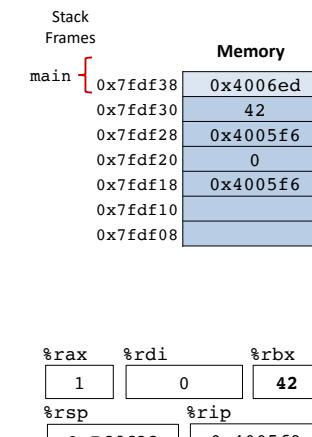


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Recursion Example: pcount(2) → pcount(1) → pcount(0)

```
long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}
```

```
pcount:
    4005dd: movl $0, %eax
    4005e2: testq %rdi, %rdi
    4005e5: je 4005fa <.L6>
    4005e7: pushq %rbx
    4005e8: movq %rdi, %rbx
    4005eb: andl $1, %ebx
    4005ee: shrq %rdi
    4005f1: callq pcount
    4005f6: addq %rbx, %rax
    4005f9: popq %rbx
.L6:
    4005fa: rep
    4005fb: retq
```

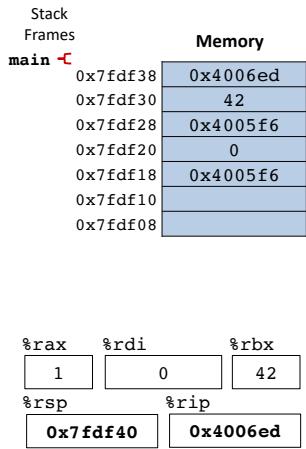


80

Recursion Example: pcount(2) → pcount(1) → pcount(0)

```
long pcount(unsigned long x) {
    if (x == 0) {
        return 0;
    } else {
        return (x & 1) + pcount(x >> 1);
    }
}

pcount:
4005dd:    movl $0, %eax
4005e2:    testq %rdi, %rdi
4005e5:    je 4005fa <.L6>
4005e7:    pushq %rbx
4005e8:    movq %rdi, %rbx
4005eb:    andl $1, %ebx
4005ee:    shrq %rdi
4005f1:    callq pcount
4005f6:    addq %rbx, %rax
4005f9:    popq %rbx
.L6:
4005fa:    rep
4005fb:    retq
```



Stack storage example (1)

```
long int call_proc()
{
    long x1 = 1;
    int x2 = 2;
    short x3 = 3;
    char x4 = 4;
    proc(x1, &x1, x2, &x2,
          x3, &x3, x4, &x4);
    return (x1+x2)*(x3-x4);
}
```

Return address to caller of call_proc ←%rsp


```
call_proc:
subq $32,%rsp
movq $1,16(%rsp) # x1
movl $2,24(%rsp) # x2
movw $3,28(%rsp) # x3
movb $4,31(%rsp) # x4
• • •
```

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Stack storage example (2) Allocate local vars

optional

```
long int call_proc()
{
    long x1 = 1;
    int x2 = 2;
    short x3 = 3;
    char x4 = 4;
    proc(x1, &x1, x2, &x2,
          x3, &x3, x4, &x4);
    return (x1+x2)*(x3-x4);
}
```

```
call_proc:
subq $32,%rsp
movq $1,16(%rsp) # x1
movl $2,24(%rsp) # x2
movw $3,28(%rsp) # x3
movb $4,31(%rsp) # x4
• • •
```

Return address to caller of call_proc

x4	x3	x2
x1		

24
16
8
←%rsp

Stack storage example (3) setup args to proc

optional

```
long int call_proc()
{
    long x1 = 1;
    int x2 = 2;
    short x3 = 3;
    char x4 = 4;
    proc(x1, &x1, x2, &x2,
          x3, &x3, x4, &x4);
    return (x1+x2)*(x3-x4);
}
```

Return address to caller of call_proc

x4	x3	x2
x1		
Arg 8		
Arg 7		

24
16
8
←%rsp

```
call_proc:
• • •
leaq 24(%rsp),%rcx # &x2
leaq 16(%rsp),%rsi # &x1
leaq 31(%rsp),%rax # &x4
movq %rax,8(%rsp) # ...
movl $4,%rsp # 4
leaq 28(%rsp),%r9 # &x3
movl $3,%r8d # 3
movl $2,%edx # 2
movq $1,%rdi # 1
call proc
• • •
```

Arguments passed in (in order):
%rdi, %rsi, %rdx, %rcx, %r8, %r9

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Stack storage example

(4) after call to proc

```
long int call_proc()
{
    long x1 = 1;
    int x2 = 2;
    short x3 = 3;
    char x4 = 4;
    proc(x1, &x1, x2, &x2,
          x3, &x3, x4, &x4);
    return (x1+x2)*(x3-x4);
}
```

Return address to caller of call_proc		
x4	x3	x2
		24
	x1	16
	Arg 8	8
	Arg 7	←%rsp

optional

```
call_proc:
    ...
    movswl 28(%rsp),%eax # x3
    movsbl 31(%rsp),%edx # x4
    subl %edx,%eax # x3-x4
    cltq # sign-extend %eax->%rax
    movslq 24(%rsp),%rdx # x2
    addq 16(%rsp),%rdx # x1+x2
    imulq %rdx,%rax # *
    addq $32,%rsp
    ret
```

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Stack storage example

(5) deallocate local vars

```
long int call_proc()
{
    long x1 = 1;
    int x2 = 2;
    short x3 = 3;
    char x4 = 4;
    proc(x1, &x1, x2, &x2,
          x3, &x3, x4, &x4);
    return (x1+x2)*(x3-x4);
}
```

optional

Return address to caller of call_proc		
		←%rsp

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Procedure Summary

call, ret, push, pop

Stack discipline fits procedure call / return.*

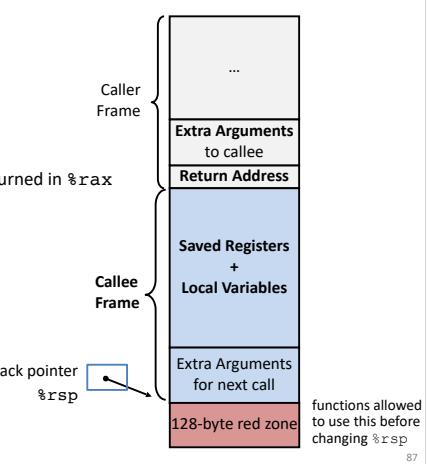
If P calls Q: Q (and calls by Q) returns before P

Conventions support arbitrary function calls.

Register-save conventions.

Stack frame saves extra args or local variables. Result returned in %rax

%rax	Return value - Caller saved
%rbx	Callee saved
%rcx	Argument #4 - Caller saved
%rdx	Argument #3 - Caller saved
%rsi	Argument #2 - Caller saved
%rdi	Argument #1 - Caller saved
%rsp	Stack pointer
%rbp	Callee saved



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