

Representing Data Structures

Multidimensional arrays

Structs

Array Layout and Indexing

int val[5];



Write x86 code to load `val[i]` into `%eax`.

1. Assume:

- Base address of `val` **is in `%rdi`**
- `i` **is in `%rsi`**

2. Assume:

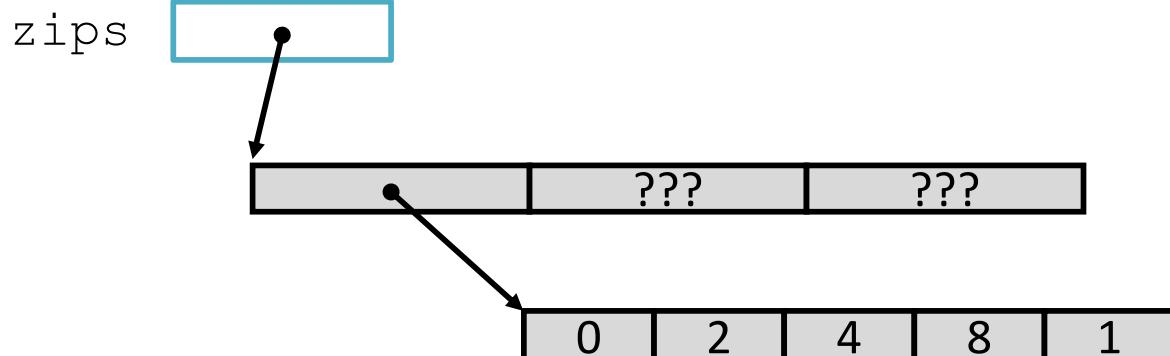
- Base address of `val` **is `28(%rsp)`**
- `i` **is in `%rcx`**

C: Arrays of pointers to arrays of ...

```
int** zips = (int**)malloc(sizeof(int*)*3);  
...  
zips[0] = (int*)malloc(sizeof(int)*5);  
...  
int* zip0 = zips[0];  
zip0[0] = 0;  
zips[0][1] = 2;  
zips[0][2] = 4;  
zips[0][3] = 8;  
zips[0][4] = 1;
```

Write x86 code to implement:

```
void copyleft(int** zips, long i, long j){  
    zips[i][j] = zips[i][j - 1];  
}
```



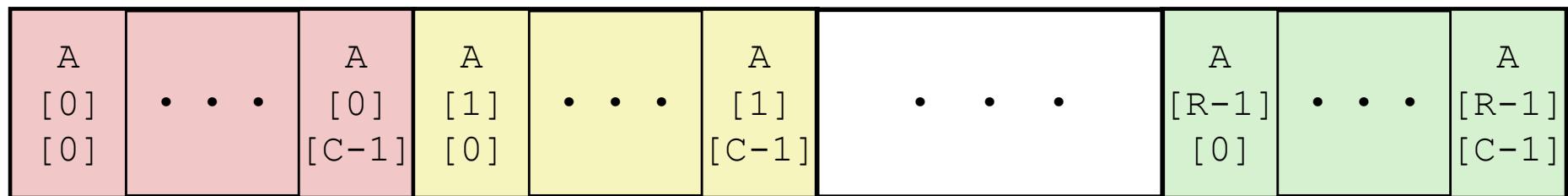
```
int[][] zips = new int[3][];  
zips[0] = new int[5] {0, 2, 4, 8, 1};
```

Java

Row-Major Nested Arrays

```
int a[R][C];
```

$$\begin{bmatrix} A[0][0] & \cdots & A[0][C-1] \\ \vdots & & \vdots \\ A[R-1][0] & \cdots & A[R-1][C-1] \end{bmatrix}$$



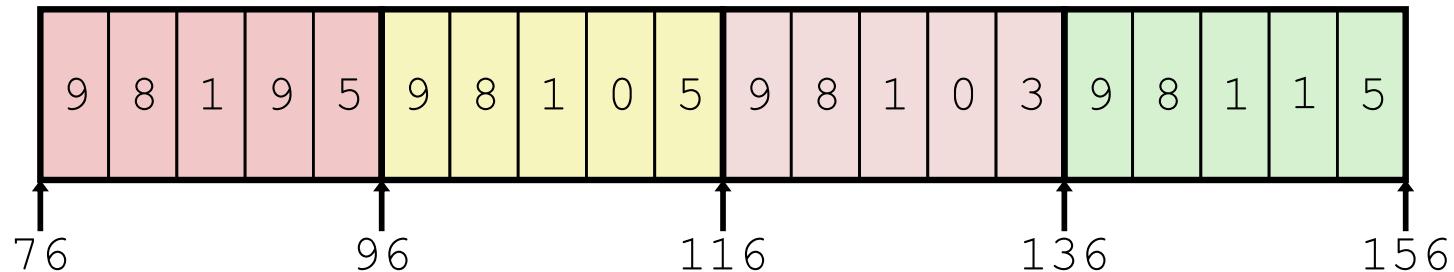
$\&a[i][j]$ is $a +$

```
int* b = (int*)a; // Can treat as larger 1D array
```

```
&a[i][j] == &b[_____]
```

Strange Referencing Examples

```
int sea[4][5];
```



Reference	Address	Value	Guaranteed?
sea[3][3]	$76 + 20 * 3 + 4 * 3 = 148$	1	Yes
sea[2][5]			
sea[2][-1]			
sea[4][-1]			
sea[0][19]			
sea[0][-1]			

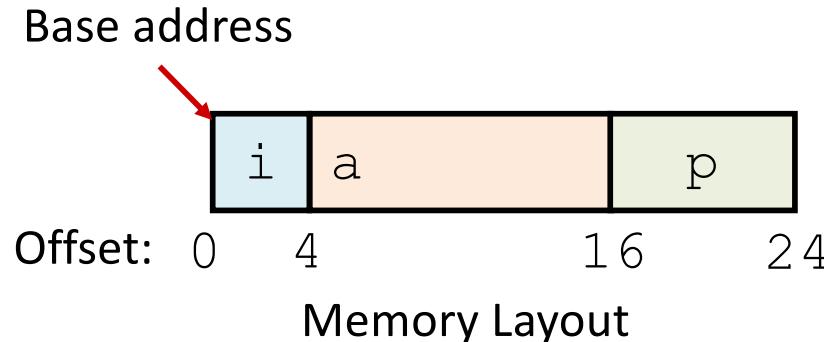
C does not do any bounds checking.

Row-major array layout is guaranteed.

C structs

```
struct rec {  
    int i;  
    int a[3];  
    int* p;  
};
```

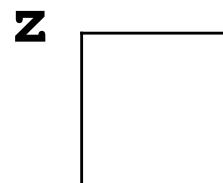
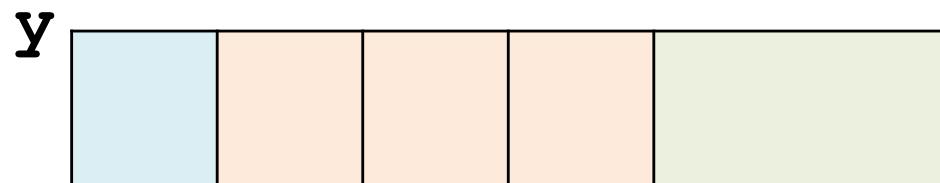
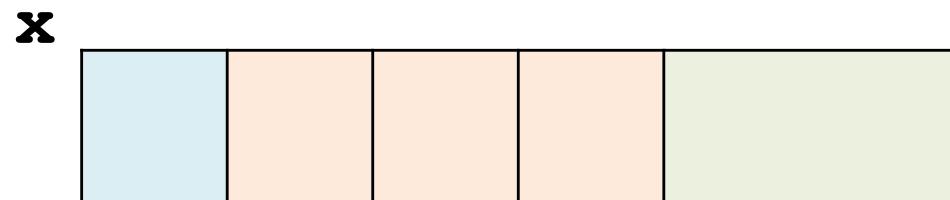
```
struct rec x;  
struct rec y;  
x.i = 1;  
x.a[1] = 2;  
x.p = &(x.i);  
// copy full struct:  
y = x;  
  
struct rec* z;  
z = &y;  
(*z).i++;  
// same as:  
z->i++
```



Like Java class/object without methods.

Compiler determines:

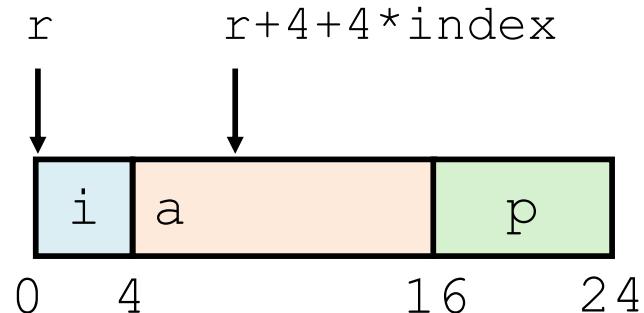
- Total size
- Offset of each field



Write x86.

Accessing Struct Field

```
struct rec {  
    int i;  
    int a[3];  
    int* p;  
};
```



```
int get_i_plus_elem(struct rec* r, int index) {  
    return r->i + r->a[index];  
}
```

```
movl 0(%rdi),%eax          # Mem[r+0]  
addl 4(%rdi,%rsi,4),%eax  # Mem[r+4*index+4]  
retq
```

typedef

```
// give type T another name: U
typedef T U;
```

```
// struct types can be verbose
struct ListNode { ... };

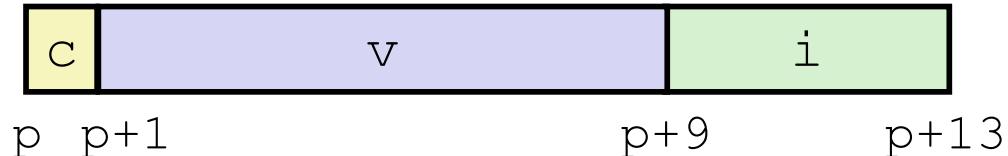
...
struct ListNode* n = ...;
```

```
// typedef can help
typedef struct list_node {
    ...
} ListNode;

...
ListNode* n = ...;
```

Struct alignment (1)

Unaligned Data



Aligned Data

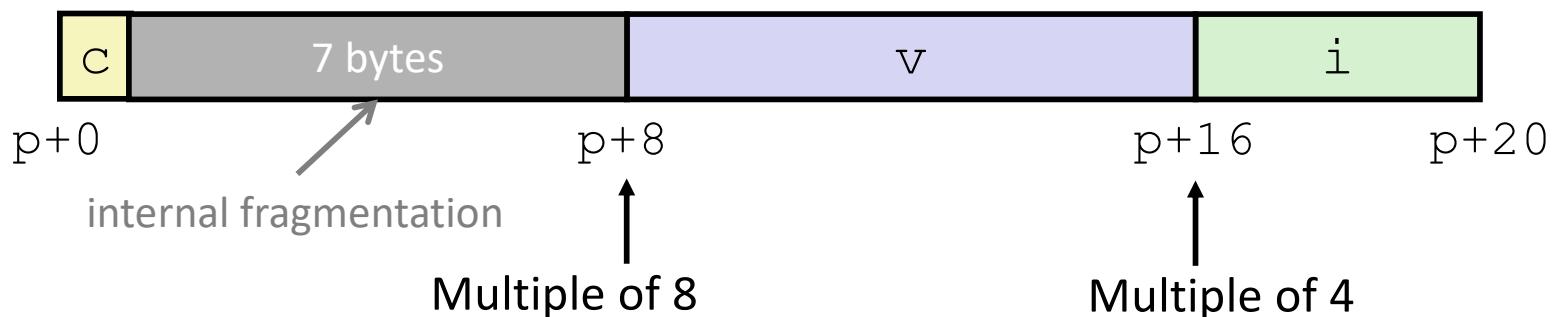
Primitive data type requires K bytes

Address must be multiple of K

C: align every struct field accordingly.

```
struct S1 {  
    char c;  
    double v;  
    int i;  
}* p;
```

Defines new struct type
and declares variable p
of type struct S1*



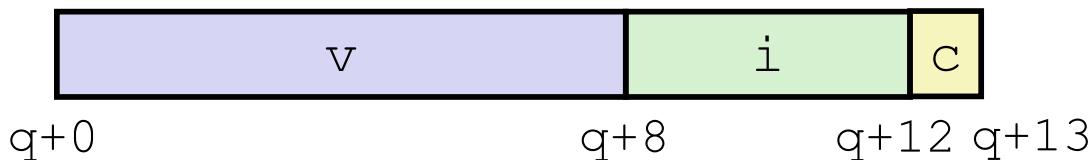
Struct packing

Put large data types first:

```
struct S1 {  
    char c;  
    double v;  
    int i;  
} * p;
```

programmer

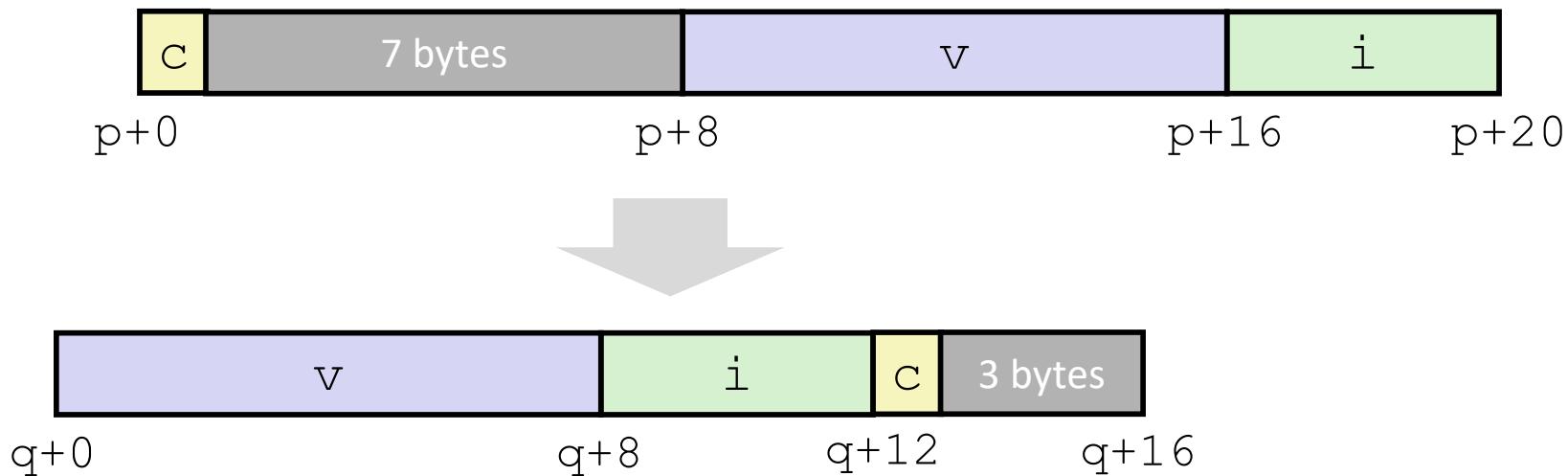
```
struct S2 {  
    double v;  
    int i;  
    char c;  
} * q;
```



But actually...

Struct alignment (full)

- Base and total size must align largest internal primitive type.
- Fields must align their type's largest alignment requirement.



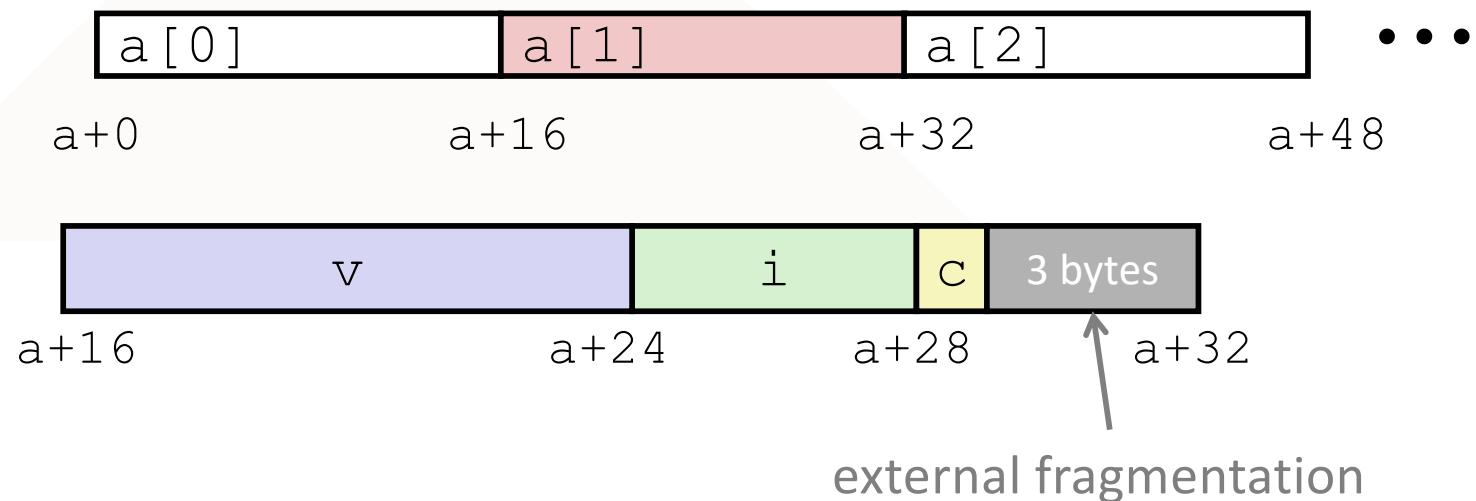
Array in struct

```
struct rec {  
    int i;  
    int a[3];  
    int* p;  
};
```

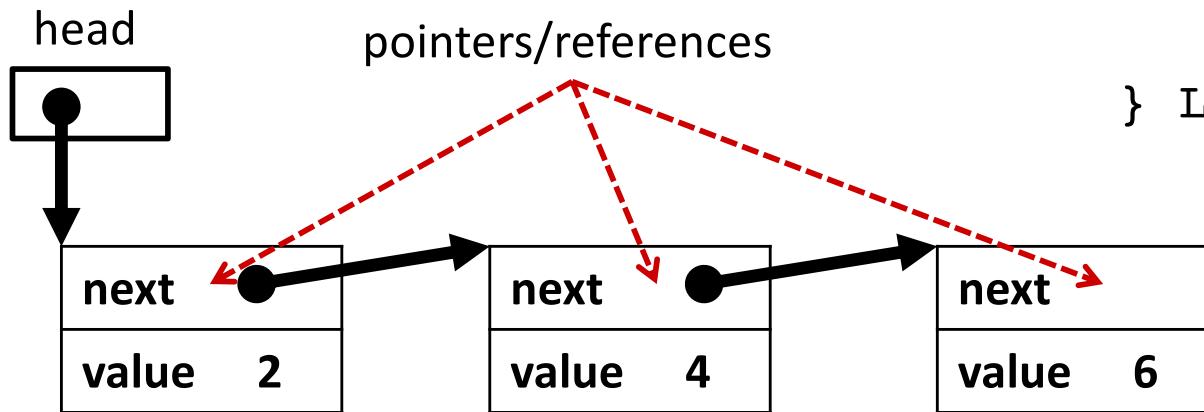


Struct in array

```
struct S2 {  
    double v;  
    int i;  
    char c;  
} a[10];
```



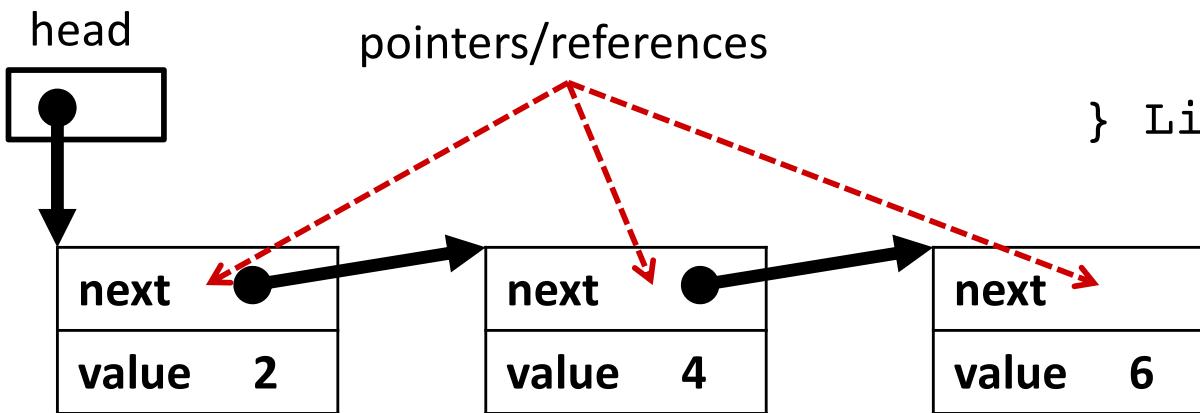
Linked Lists



```
typedef  
struct ListNode {  
    struct ListNode* next;  
    int value;  
} ListNode;
```

1. Choose memory layout for ListNode

Linked Lists



```
typedef
struct ListNode {
    struct ListNode* next;
    int value;
} ListNode;
```

2. Implement append in x86:

```
void append(ListNode* head, int x) { // assume head != NULL
    ListNode* cursor = head;
    while (cursor->next != NULL) { // find tail
        cursor = cursor->next;
    }
    ListNode* n = (ListNode*)malloc(sizeof(ListNode));
    // error checking omitted for x86 simplicity
    cursor->next = n;
    n->next = NULL;
    n->value = x;
}
```

Try a recursive version too.