

CSAPP book is **very useful** and well-aligned with class for the remainder of the course.

C to Machine Code and x86 Basics

ISA context and x86 history

Translation tools: C --> assembly <--> machine code

x86 Basics:

- Registers
- Data movement instructions
- Memory addressing modes
- Arithmetic instructions

x86-64 registers

%rax	Return Value
%rbx	
%rcx	Argument 4
%rdx	Argument 3
%rsi	Argument 2
%rdi	Argument 1
%rsp	Special Purpose: Stack Pointer
%rbp	
%r8	Argument 5
%r9	Argument 6
%r10	
%r11	
%r12	
%r13	
%r14	
%r15	

64-bits / 8 bytes

sub-registers

1985: 32-bit extended register %eax
1978: 16-bit register %ax



historical artifacts



Some have special uses for particular instructions

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Pointers and Memory Addressing

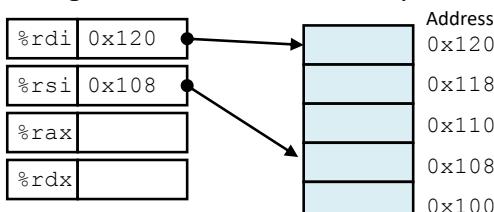
```
void swap(long* xp, long* yp) {
    long t0 = *xp;
    long t1 = *yp;
    *xp = t1;
    *yp = t0;
}
```

```
swap:
    movq (%rdi),%rax
    movq (%rsi),%rdx
    movq %rdx,(%rdi)
    movq %rax,(%rsi)
    retq
```

Registers

Register	Variable
%rdi	\leftrightarrow xp
%rsi	\leftrightarrow yp
%rax	\leftrightarrow t0
%rdx	\leftrightarrow t1

Memory



Address Computation Examples

ex

General Addressing Modes

D(Rb,Ri,S) $\text{Mem}[\text{Reg}[Rb]+S*\text{Reg}[Ri] + D]$

Register contents

%rdx	0xf000
%rcx	0x100

Special Cases:

Implicitly:

(Rb,Ri) $\text{Mem}[\text{Reg}[Rb]+\text{Reg}[Ri]]$

(S=1,D=0)

D(Rb,Ri) $\text{Mem}[\text{Reg}[Rb]+\text{Reg}[Ri]+D]$

(S=1)

(Rb,Ri,S) $\text{Mem}[\text{Reg}[Rb]+S*\text{Reg}[Ri]]$

(D=0)

Address Expression	Address Computation	Address
0x8(%rdx)		
(%rdx,%rcx)		
(%rdx,%rcx,4)		
0x80(,%rdx,2)		

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Compute address given by
this addressing mode expression
and store it here.

leaq Src, Dest

Load effective address



DOES NOT ACCESS MEMORY

Uses: "address of" "Lovely Efficient Arithmetic"

$p = \&x[i];$ $x + k*I$, where $k = 1, 2, 4, \text{ or } 8$

leaq vs. movq

Registers
%rax
%rbx
%rcx 0x4
%rdx 0x100
%rdi
%rsi

Memory
0x400
0xf
0x8
0x110
0x10
0x1

Address
0x120
0x118
0x8
0x110
0x108
0x100

Assembly Code

```
leaq (%rdx,%rcx,4), %rax
movq (%rdx,%rcx,4), %rbx
leaq (%rdx), %rdi
movq (%rdx), %rsi
```

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Procedure Preview (more soon)

call, ret, push, pop

Procedure arguments passed in 6 registers:

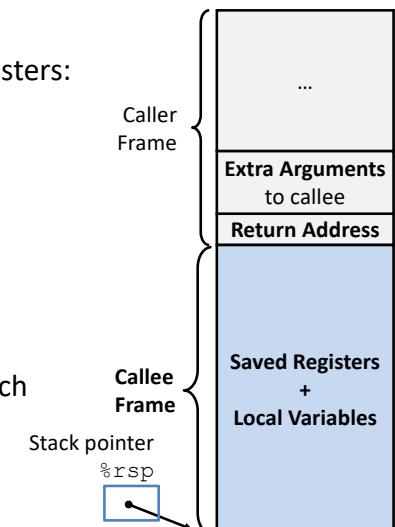
%rax	Return Value
%rbx	Argument 5
%rcx	Argument 4
%rdx	Argument 3
%rsi	Argument 2
%rdi	Argument 1
%rsp	Stack pointer
%rbp	%r15

Return value in %rax.

Allocate/push new *stack frame* for each procedure call.

Some local variables,
saved register values, extra arguments

Deallocate/pop frame before return.



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Arithmetic Operations

Two-operand instructions:

Format	Computation	
addq Src,Dest	$Dest = Dest + Src$	
subq Src,Dest	$Dest = Dest - Src$	← argument order
imulq Src,Dest	$Dest = Dest * Src$	
shlq Src,Dest	$Dest = Dest \ll Src$	a.k.a salq
sarq Src,Dest	$Dest = Dest \gg Src$	Arithmetic
shrq Src,Dest	$Dest = Dest \gg Src$	Logical
xorq Src,Dest	$Dest = Dest \wedge Src$	
andq Src,Dest	$Dest = Dest \& Src$	
orq Src,Dest	$Dest = Dest Src$	

One-operand (unary) instructions

incq Dest	$Dest = Dest + 1$	increment
decq Dest	$Dest = Dest - 1$	decrement
negq Dest	$Dest = -Dest$	negate
notq Dest	$Dest = \sim Dest$	bitwise complement

See CSAPP 3.5.5 for: mulq, cqto, idivq, divq

leaq for arithmetic

```
long arith(long x, long y,
           long z){
    long t1 = x+y;
    long t2 = z+t1;
    long t3 = x+4;
    long t4 = y * 48;
    long t5 = t3 + t4;
    long rval = t2 * t5;
    return rval;
}
```

```
arith:
    leaq    (%rdi,%rsi), %rax
    addq   %rdx, %rax
    leaq   (%rsi,%rsi,2), %rdx
    salq   $4, %rdx
    leaq   4(%rdi,%rdx), %rcx
    imulq  %rcx, %rax
    ret
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rdx	Argument z
%rax	
%rcx	

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Another example

```
long logical(long x, long y){  
    long t1 = x^y;  
    long t2 = t1 >> 17;  
    long mask = (1<<13) - 7;  
    long rval = t2 & mask;  
    return rval;  
}
```

Register	Use(s)
%rdi	Argument x
%rsi	Argument y
%rax	

```
logical:  
    movq %rdi,%rax  
    xorq %rsi,%rax  
    sarq $17,%rax  
    andq $8185,%rax  
    retq
```