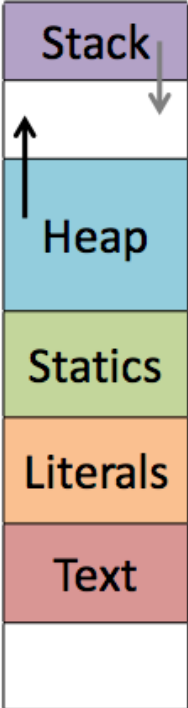


CS240 Laboratory 9

Disassembly and Reverse Engineering

Memory Layout

	Addr		Perm	Contents	Managed by	Initialized
Kernel above 0x7fffffff	Addr 2^N-1  0	Stack	RW	Procedure context	Compiler	Run-time
Stack below 0x7fffffff grows down						
Heap above Data segment		Heap	RW	Dynamic data structures	Programmer, malloc/free, new/GC	Run-time
Data segment statics and literals		Statics	RW	Global variables/ static data structures	Compiler/ Assembler/Linker	Startup
		Literals	R	String literals	Compiler/ Assembler/Linker	Startup
Text segment starts at 0x400000		Text	X	Instructions	Compiler/ Assembler/Linker	Startup

Instructions

Moving Data

mov Src, Dest

Note: the size of the data being referenced is often specified with an additional character:

b (byte)

w (2 bytes)

l (4 bytes), or

q (8 bytes).

Arithmetic/Logical operations – 2 operands

add Src, Dest

sub Src, Dest

imul Src, Dest

shr Src, Dest

sar Src, Dest

shl Src, Dest

sal Src, Dest

shr Src, Dest

xor Src, Dest

and Src, Dest

or Src, Dest

Memory -0x18(%rsp)

Most General Form:

$$D(Rb, Ri, S) \qquad \text{Mem}[\text{Reg}[Rb] + S * \text{Reg}[Ri] + D]$$

D: Constant "displacement" value represented in 1, 2, or 4 bytes

Rb: Base register: Any register

Ri: Index register: Any except %esp (or %rsp if 64-bit); %ebp unlikely

S: Scale: 1, 2, 4, or 8 (*why these numbers?*)

Special Cases: can use any combination of D, Rb, Ri and S

$$(Rb, Ri) \qquad \text{Mem}[\text{Reg}[Rb] + \text{Reg}[Ri]] \qquad (S=1, D=0)$$

$$D(Rb, Ri) \qquad \text{Mem}[\text{Reg}[Rb] + \text{Reg}[Ri] + D] \qquad (S=1)$$

$$(Rb, Ri, S) \qquad \text{Mem}[\text{Reg}[Rb] + S * \text{Reg}[Ri]] \qquad (D=0)$$

Control Flow

Conditional jump instructions in X86 implement the following high-level constructs:

- if (condition) then {...} else {...}
- while (condition) {...}
- do {...} while (condition)
- for (initialization; condition; iterative) {...}

Unconditional jumps are used for high-level constructs such as:

- break
- continue

jX	Condition	Description
jmp	1	Unconditional
je	ZF	Equal / Zero
jne	~ZF	Not Equal / Not Zero
ja	SF	Negative
jns	~SF	Nonnegative
jg	~ (SF^OF) & ~ZF	Greater (Signed)
jge	~ (SF^OF)	Greater or Equal (Signed)
jl	(SF^OF)	Less (Signed)
jle	(SF^OF) ZF	Less or Equal (Signed)
ja	~CF & ~ZF	Above (unsigned)
jb	CF	Below (unsigned)

Jump instructions encode the offset from next instruction to destination PC, instead of the absolute address of the destination (makes it easier to relocate the code)

Turning C into Machine Code

C Code

```
void sumstore(long x, long y,  
              long *dest) {  
    long t = x + y;  
    *dest = t;  
}
```

sum.c

Generated x86 Assembly Code

Human-readable language close to machine code.

```
sum:  
    addq %rdi,%rsi  
    movq %rsi,(%rdx)  
    retq
```

sum.s

compiler (CS 301)

gcc -Og -S sum.c

assembler

Object Code

```
01010101100010011110010110  
00101101000101000011000000  
00110100010100001000100010  
01111011000101110111000011
```

sum.o

linker

Executable: sum

Resolve references between object files,
libraries, (re)locate data

- X86 instructions can be in different order from C code
- Some C expressions require multiple X86 instructions
- Some X86 instructions can cover multiple C expressions
- Compiler optimization can do some surprising things!
- Local or temporary variables can be stored in registers or on the stack

Function Calling Conventions

- Arguments for functions are stored in registers, in the following order: arg1 – arg6: `%rdi,%rsi,%rdx,%rcx,%r8,%r9`
- If there are more than 6 parameters for a function, the rest of the arguments are stored on the stack before the function is called
- Return value from function is always in `%rax`

The compiler will use only part of a register if the value stored there will fit in less than 64 bits (8 bytes). This is an optimization that makes instructions a bit shorter.

So, in the code, you may see register names of the following form, all of which refer to `%rax`:

`%rax` = 8 byte value

`%eax` = 4 byte value

`%ax` = 2 byte value

`%al` = 1 byte value

Tools

Tools can be used to examine bytes of object code (executable program) and reconstruct (reverse engineer) the assembly source.

gdb – disassembles an executable file into the associated assembly language representation, and provides tools for memory and register examination, single step execution, breakpoints, etc.

Object	Disassembled by GDB
0x00400536:	0x0000000000400536 <+0>: add %rdi,%rsi
0x48	0x0000000000400539 <+3>: mov %rsi,(%rdx)
0x01	0x000000000040053c <+6>: retq
0xfe	
0x48	\$ gdb sum
0x89	(gdb) disassemble sumstore
0x32	(disassemble function)
0xc3	(gdb) x/7b sum
	(examine the 13 bytes starting at sum)

objdump

can also be used to disassemble and display information

```
$ objdump -t p
```

Prints out the program's symbol table. The symbol table includes the names of all functions and global variables, the names of all the functions the called, and their addresses.

\$ objdump -d p

Object Code

0x401040 <sum>:

0x55

0x89

0xe5

0x8b

0x45

0x0c

0x03

0x45

0x08

0x89

0xec

0x5d

0xc3

Disassembled version

00401040 <_sum>:

0: 55 push %ebp

1: 89 e5 mov %esp,%ebp

3: 8b 45 0c mov 0xc(%ebp),%eax

6: 03 45 08 add 0x8(%ebp),%eax

9: 89 ec mov %ebp,%esp

b: 5d pop %ebp

c: c3 ret

strings

\$ strings -t x p

Displays the printable strings in your program.