



Integer Representation

Representation of integers: unsigned and signed
Modular arithmetic and overflow
Sign extension
Shifting and arithmetic
Multiplication
Casting

Fixed-width integer encodings

Unsigned $\subset \mathbb{N}$ non-negative integers only

Signed $\subset \mathbb{Z}$ both negative and non-negative integers

n bits offer only 2^n distinct values.

Terminology:

“Most-significant” bit(s)
or “high-order” bit(s)

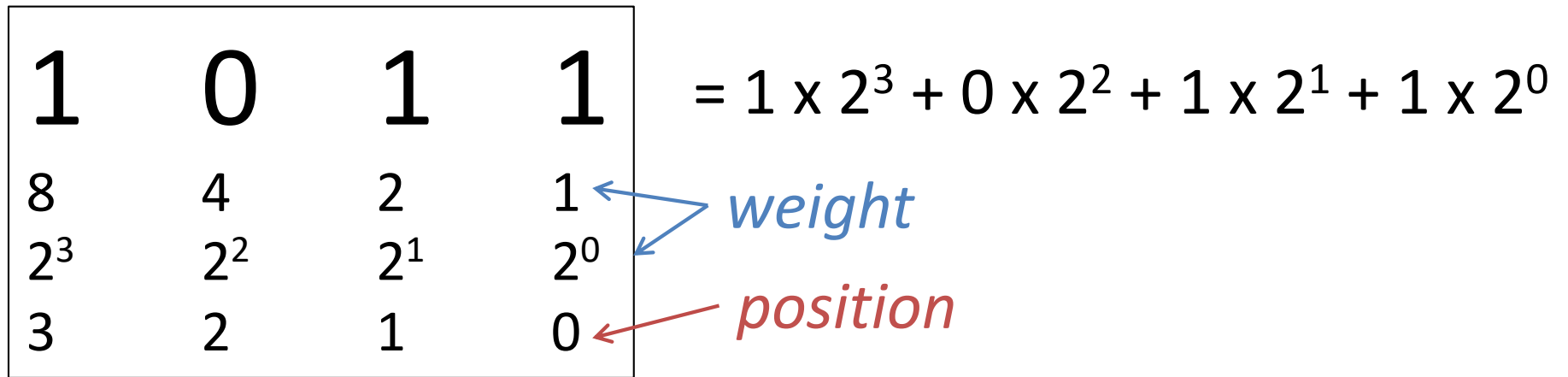
“Least-significant” bit(s)
or “low-order” bit(s)

MSB

0110010110101001

LSB

(4-bit) unsigned integer representation



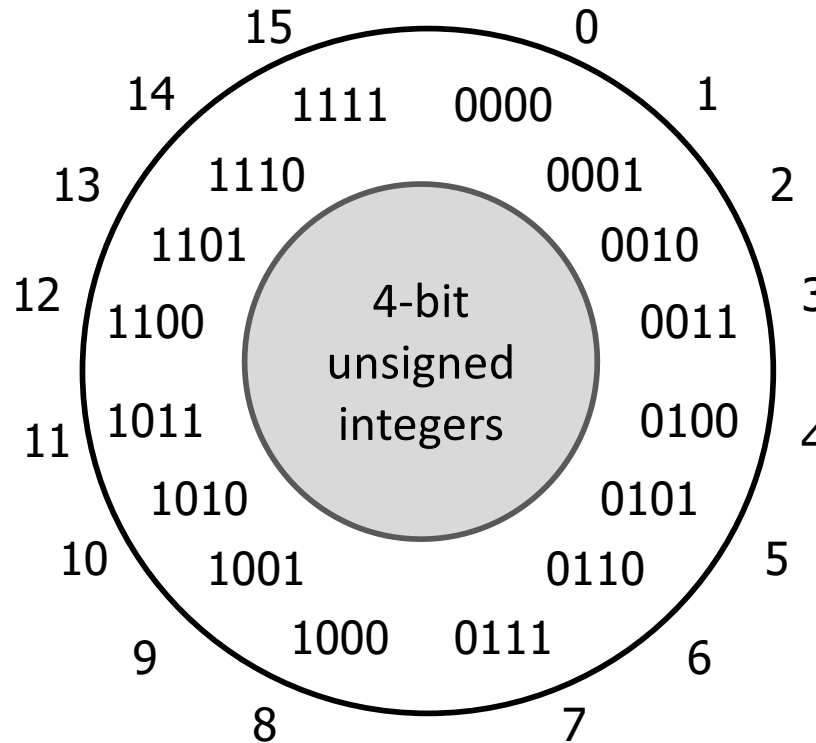
n -bit unsigned integers:

minimum =

maximum =

modular arithmetic, overflow

$$\begin{array}{r} 11 \quad 1011 \\ + 2 \quad + 0010 \\ \hline \end{array}$$



$$\begin{array}{r} 13 \quad 1101 \\ + 5 \quad + 0101 \\ \hline \end{array}$$

$x+y$ in n -bit unsigned arithmetic is

in math

unsigned overflow =
=

Unsigned addition *overflows* if and only if

sign-magnitude



Most-significant bit (MSB) is *sign bit*

0 means non-negative 1 means negative

Remaining bits are an unsigned magnitude

8-bit sign-magnitude: _____

00000000 represents _____

01111111 represents _____

10000101 represents _____

10000000 represents _____

Anything weird here?

Arithmetic?

Example:

$$4 - 3 \neq 4 + (-3)$$



$$\begin{array}{r} 00000100 \\ + 10000011 \\ \hline \end{array}$$

Zero?



(4-bit) two's complement signed integer representation



1	0	1	1
-2^3	2^2	2^1	2^0

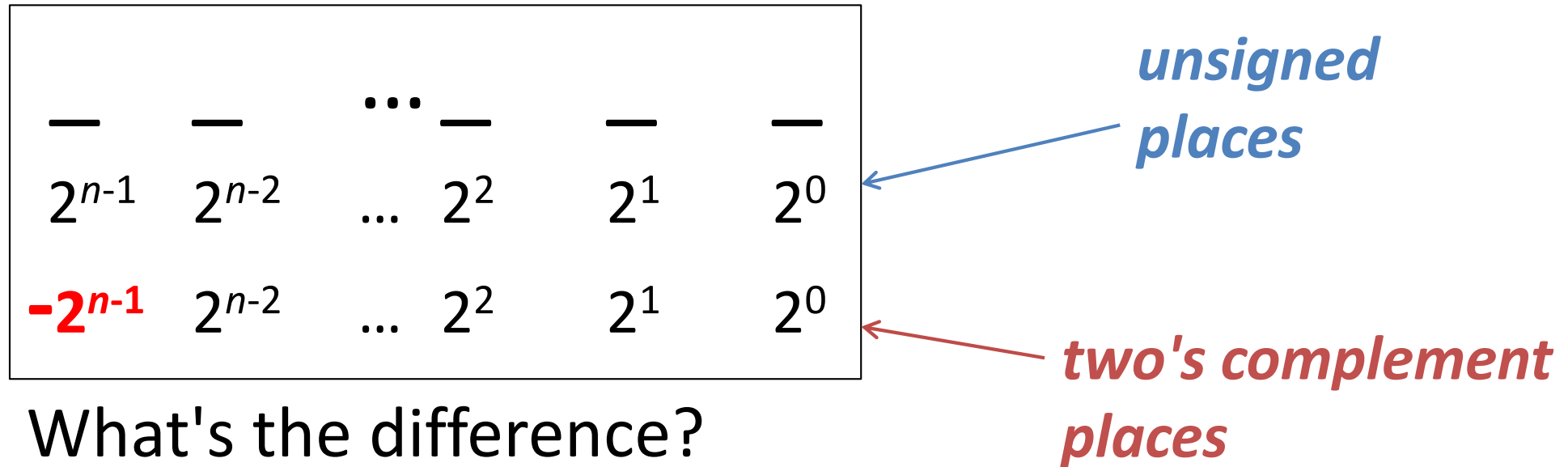
$$= 1 \times -2^3 + 0 \times 2^2 + 1 \times 2^1 + 1 \times 2^0$$

4-bit two's complement integers:

minimum =

maximum =

two's complement vs. unsigned



n-bit unsigned numbers:

minimum =

maximum =

8-bit representations



0 0 0 0 1 0 0 1

1 0 0 0 0 0 0 1

1 1 1 1 1 1 1 1

0 0 1 0 0 1 1 1

n-bit two's complement numbers:

minimum =

maximum =

4-bit unsigned vs. 4-bit two's complement

1 0 1 1

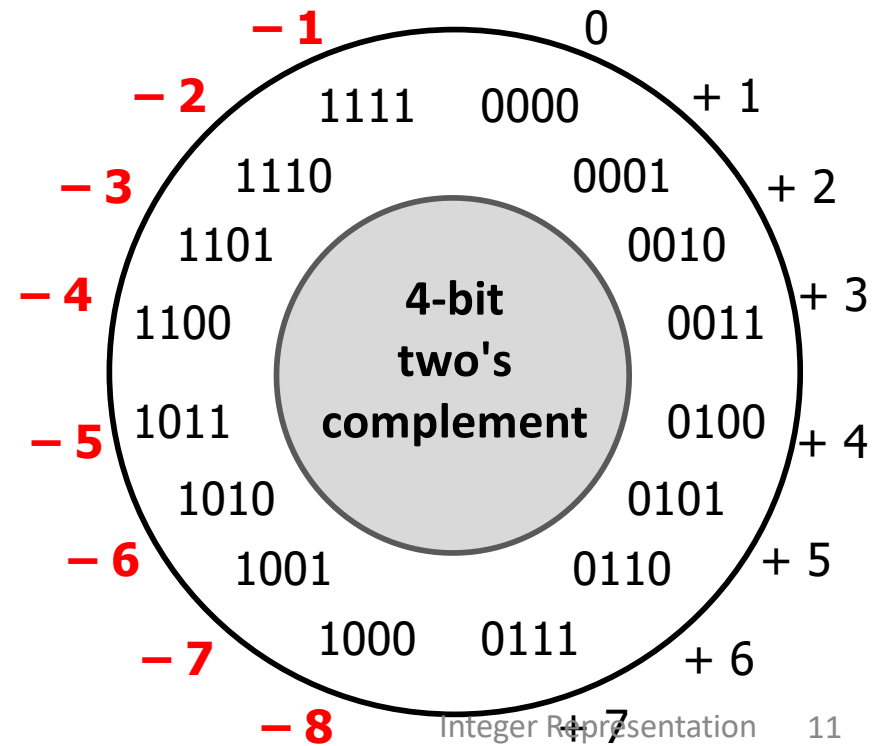
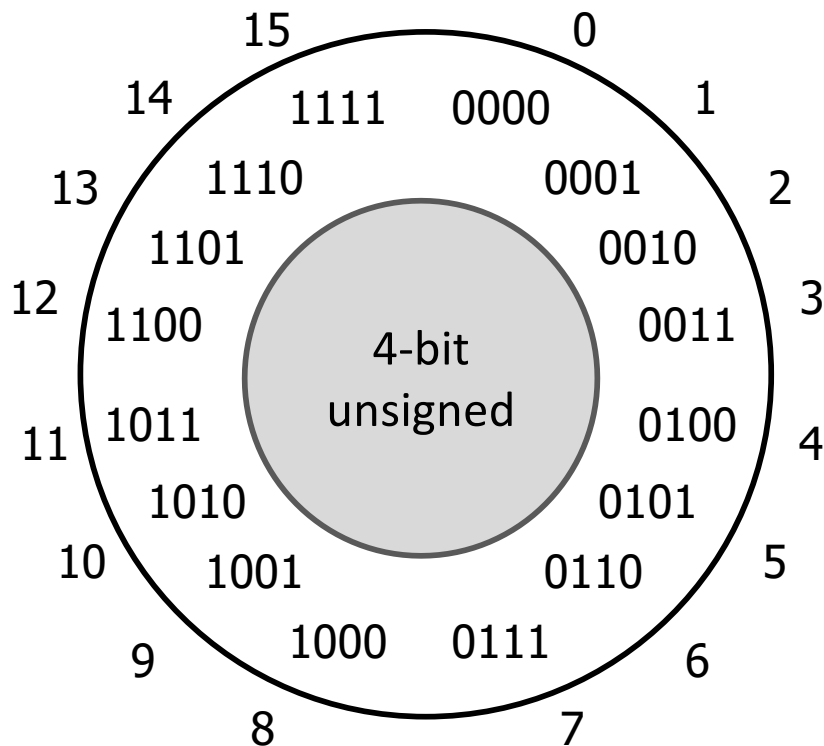
$$1 \times 2^3 + 0 \times 2^2 + 1 \times 2^1 + 1 \times 2^0$$

$$1 \times -2^3 + 0 \times 2^2 + 1 \times 2^1 + 1 \times 2^0$$

11

difference = ___ = 2

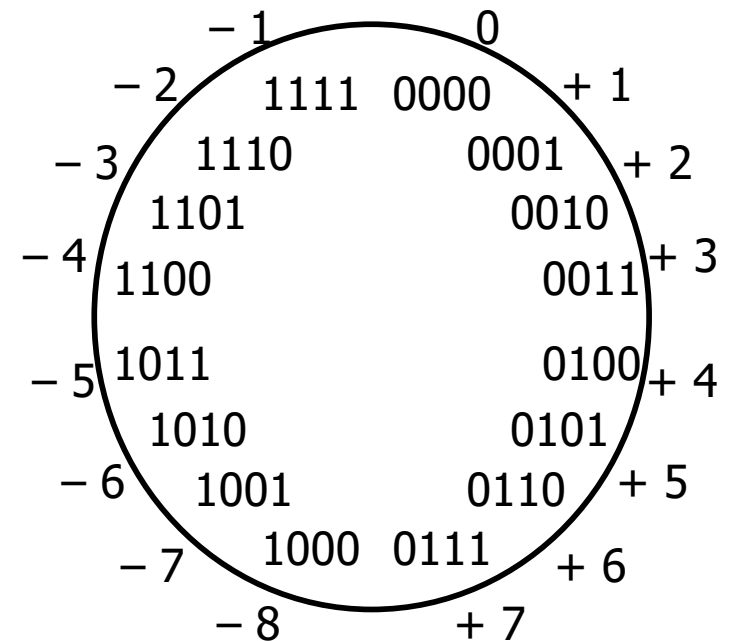
-5



two's complement addition

2	0010	-2	1110
<u>+ 3</u>	<u>+ 0011</u>	<u>+ -3</u>	<u>+ 1101</u>

-2	1110	2	0010
<u>+ 3</u>	<u>+ 0011</u>	<u>+ -3</u>	<u>+ 1101</u>



Modular Arithmetic

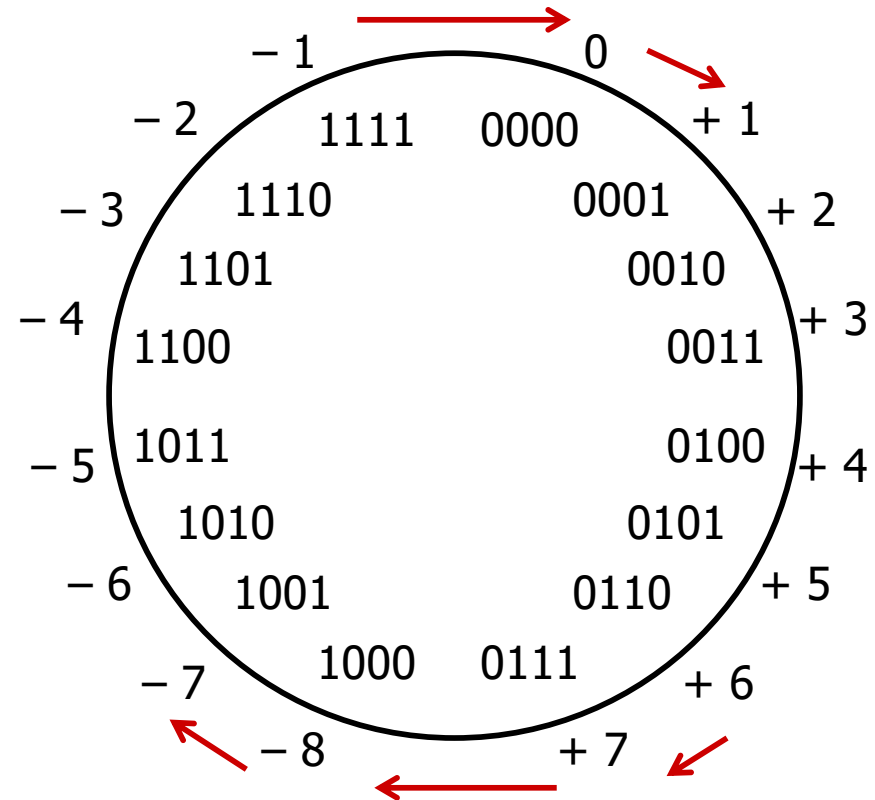
two's complement *overflow*

Addition *overflows*

if and only if
if and only if

$$\begin{array}{r} -1 \qquad 1111 \\ + 2 \qquad + 0010 \\ \hline \end{array}$$

$$\begin{array}{r} 6 \qquad 0110 \\ + 3 \qquad + 0011 \\ \hline \end{array}$$



Modular Arithmetic

Some CPUs/languages raise exceptions on overflow.
C and Java cruise along silently... Feature? Oops?

Reliability

Ariane 5 Rocket, 1996

Exploded due to **cast** of 64-bit floating-point number to 16-bit signed number.
Overflow.



Boeing 787, 2015



"... a **Model 787 airplane** ... can lose all alternating current (AC) electrical power ... caused by a **software counter** internal to the GCUs that will **overflow** after **248 days** of continuous power. We are issuing this AD to prevent loss of all AC electrical power, which could result in **loss of control of the airplane.**"
--FAA, April 2015

A few reasons two's complement is awesome

Addition, subtraction, hardware

Sign

Negative one

Complement rules



Another derivation

How should we represent 8-bit negatives?

- For all positive integers x , we want the representations of x and $-x$ to sum to zero.
- We want to use the standard addition algorithm.

$$\begin{array}{r} 00000001 \\ + \\ \hline 00000000 \end{array} \quad \begin{array}{r} 00000010 \\ + \\ \hline 00000000 \end{array} \quad \begin{array}{r} 00000011 \\ + \\ \hline 00000000 \end{array}$$

- Find a rule to represent $-x$ where that works...

Convert/cast signed number to larger type.

0 0 0 0 0 0 1 0 8-bit 2

----- 0 0 0 0 0 0 1 0 16-bit 2

1 1 1 1 1 1 0 0 8-bit -4

----- 1 1 1 1 1 1 0 0 16-bit -4

Rule/name?

unsigned shifting and arithmetic

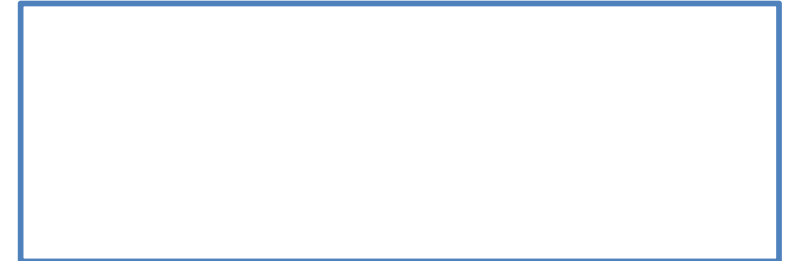
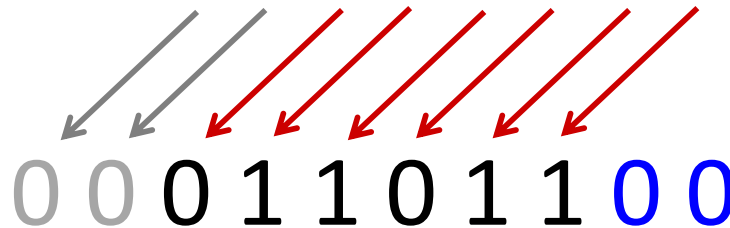
unsigned

$x = 27;$

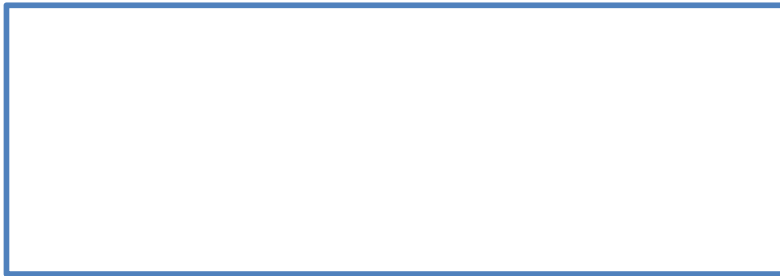
$y = x \ll 2;$

$y == 108$

0 0 0 1 1 0 1 1

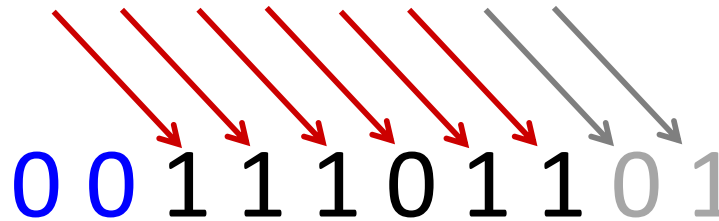


logical shift left



logical shift right

1 1 1 0 1 1 0 1



unsigned

$x = 237;$

$y = x \gg 2;$

$y == 59$

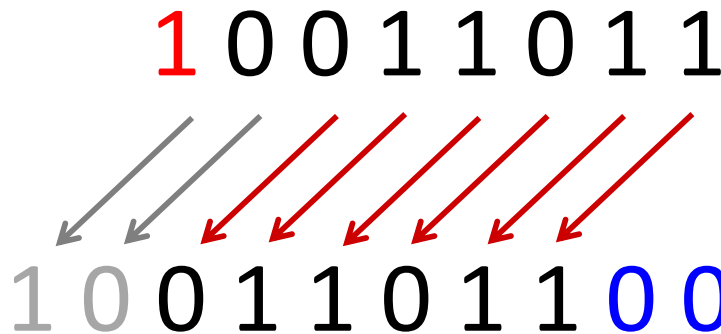
two's complement **shifting** and **arithmetic**

signed

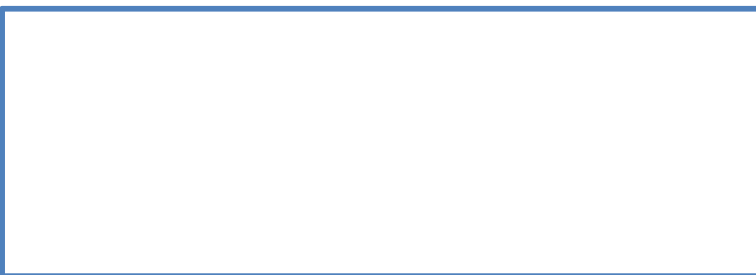
x = -101;

y = x << 2;

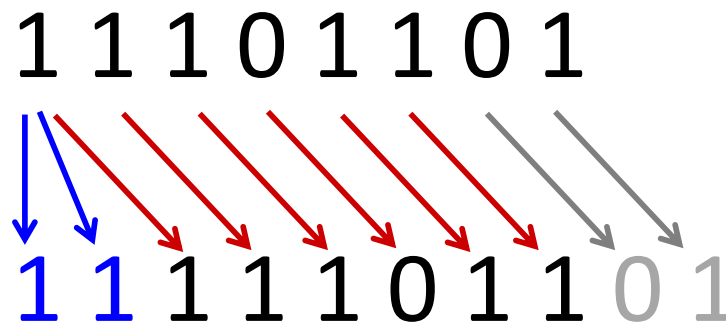
y == 108



logical shift left



arithmetic shift right



signed

x = -19;

y = x >> 2;

y == -5



shift-and-add

Available operations

$x \ll k$

implements $x * 2^k$

$x + y$

Implement $y = x * 24$ using only \ll , $+$, and integer literals

What does this function compute?

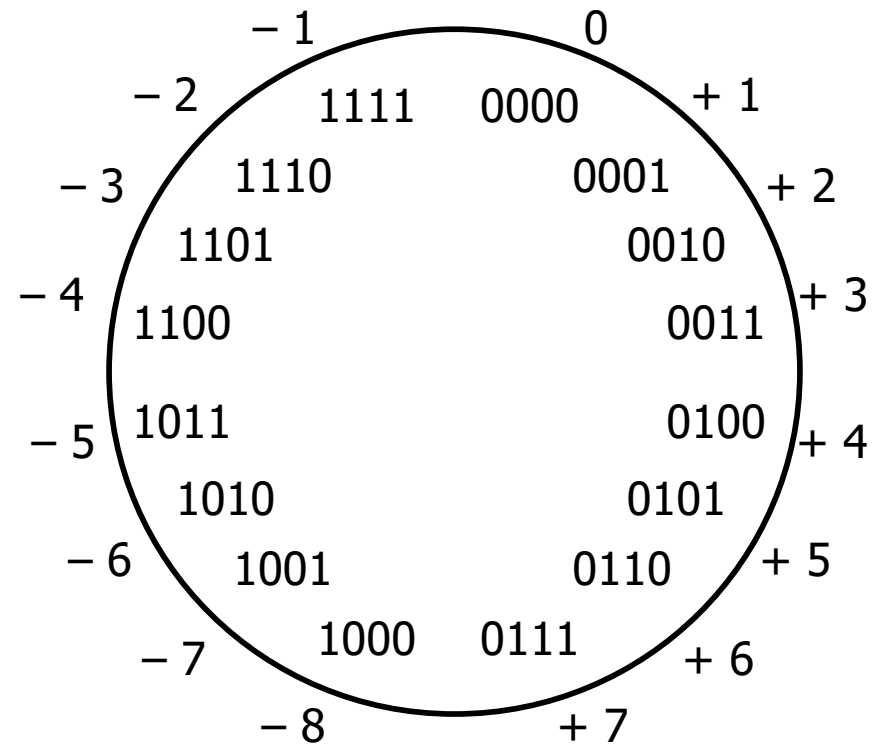


```
unsigned puzzle(unsigned x, unsigned y) {  
    unsigned result = 0;  
    for (unsigned i = 0; i < 32; i++) {  
        if (y & (1 << i)) {  
            result = result + (x << i);  
        }  
    }  
    return result;  
}
```

multiplication

$$\begin{array}{r} 2 \\ \times 3 \\ \hline 6 \end{array} \quad \begin{array}{r} 0010 \\ \times 0011 \\ \hline 00000110 \end{array}$$

$$\begin{array}{r} -2 \\ \times 2 \\ \hline -4 \end{array} \quad \begin{array}{r} 1110 \\ \times 0010 \\ \hline 11111100 \end{array}$$



Modular Arithmetic

multiplication

$$\begin{array}{r} 5 \\ \times 4 \\ \hline \end{array}$$

$$\begin{array}{r} \cancel{20} \\ 4 \end{array}$$

$$-3$$

$$\begin{array}{r} \times 7 \\ \hline \end{array}$$

$$\begin{array}{r} \cancel{-21} \\ -5 \end{array}$$

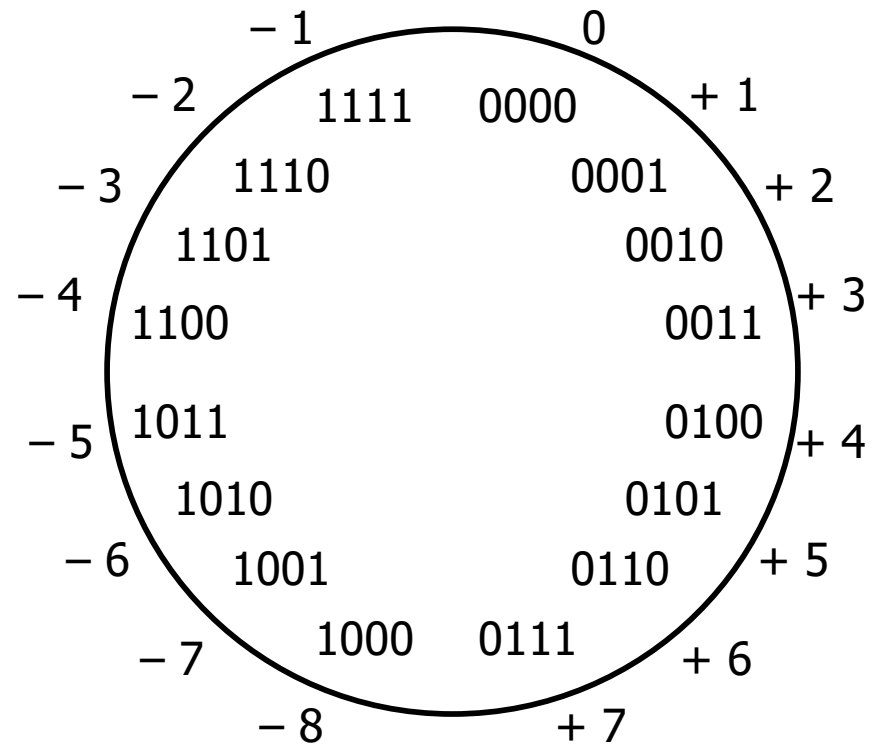
$$-5$$

$$\begin{array}{r} 0101 \\ \times 0100 \\ \hline \end{array}$$

$$00010100$$

$$\begin{array}{r} 1101 \\ \times 0111 \\ \hline \end{array}$$

$$11101011$$



Modular Arithmetic

multiplication

$$\begin{array}{r} 5 \\ \times 5 \\ \hline \end{array}$$

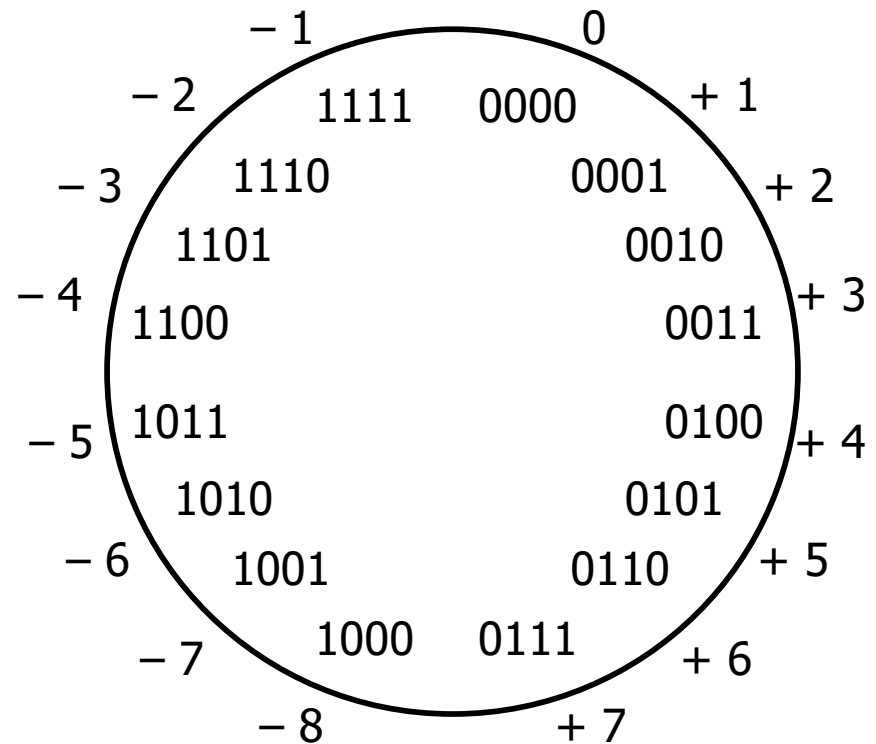
$$\begin{array}{r} \cancel{25} \\ -7 \end{array}$$

$$\begin{array}{r} -2 \\ \times 6 \\ \hline \end{array}$$

$$\begin{array}{r} \cancel{-12} \\ 4 \end{array}$$

$$\begin{array}{r} 0101 \\ \times 0101 \\ \hline 00011001 \end{array}$$

$$\begin{array}{r} 1110 \\ \times 0110 \\ \hline 11110100 \end{array}$$



Modular Arithmetic



Casting Integers in C

Number literals: `37` is signed, `37U` is unsigned

Integer Casting: *bits unchanged, just reinterpreted.*

Explicit casting:

```
int tx = (int) 73U;    // still 73
unsigned uy = (unsigned) -4; // big positive #
```

Implicit casting: Actually does

```
tx = ux;    // tx = (int)ux;
uy = ty;    // uy = (unsigned)ty;
void foo(int z) { ... }
foo(ux);    // foo((int)ux);
if (tx < ux) ... // if ((unsigned)tx < ux) ...
```

More Implicit Casting in C



If you mix unsigned and signed in a single expression, then *signed values are implicitly cast to unsigned*.

How are the argument bits interpreted?

Argument ₁	Op	Argument ₂	Type	Result
0	==	0U	unsigned	1
-1	<	0	signed	1
-1	<	0U	unsigned	0
2147483647	<	-2147483647-1		
2147483647U	<	-2147483647-1		
-1	<	-2		
(unsigned)-1	<	-2		
2147483647	<	2147483648U		
2147483647	<	(int)2147483648U		

Note: $T_{min} = -2,147,483,648$ $T_{max} = 2,147,483,647$

T_{min} must be written as $-2147483647-1$ (see pg. 77 of CSAPP for details)