CS240 Laboratory 6 Disassembly and Reverse Engineering

Memory Layout

| Kernel above 0x7fffffff | Addr | | Perm | Contents | Managed by | Initialized |
|--|--------------------|----------|------|---|---------------------------------------|-------------|
| Stack | 2 ^ℕ -1↑ | Stack | RW | Procedure context | Compiler | Run-time |
| below Ox7fffffff grows down | | ↑ | | | | |
| Heap above Data segment | | Heap | RW | Dynamic data structures | Programmer, malloc/free, new/GC | Run-time |
| Data segment statics and literals | | Statics | RW | Global variables/ static data structures | Compiler/ Assembler/Linker | Startup |
| | | Literals | R | String literals | Compiler/ Assembler/Linker | Startup |
| Text segment starts at 0x400000 | | Text | х | Instructions | Compiler/ Assembler/Linker | Startup |
| addresses below 0x400000 reserved for operating system | 0 | | | | | |

Instructions

The size of the data being referenced is often specified with an additional character:

b (byte)w (2 bytes)l (4 bytes), orq (8 bytes).

Operand Types

Immediate \$0x400, \$-533

Register: %rax,%rbx,%rcx,%rdx,%rsi,%rdi,%rbp,%rsp, %r8,%r9,%r10,%r11,%r12,%r13,%r14,%r15

some have special purpose: **%rsp** is stack pointer, **%rax** always used to return value from functions

Memory -0x18(%rsp)

Most General Form:

D(Rb,Ri,S) Mem[Reg[Rb] + S*Reg[Ri] + D]

D: Constant "displacement" value represented in 1, 2, or 4 bytes

Rb: Base register: Any register

Ri: Index register: Any except %esp (or %rsp if 64-bit); %ebp unlikely

S: Scale: 1, 2, 4, or 8 (why these numbers?)

Special Cases: can use any combination of D, Rb, Ri and S

| (Rb,Ri) | Mem[Reg[Rb]+Reg[Ri]] | (S=1, D=0) |
|-----------|------------------------|------------|
| D(Rb,Ri) | Mem[Reg[Rb]+Reg[Ri]+D] | (S=1) |
| (Rb,Ri,S) | Mem[Reg[Rb]+S*Reg[Ri]] | (D=0) |

Types of Instruction

Moving Data

movl Src,Dest // copy 4 bytes from source to destination

Arithmetic/Logical operations – 2 operands

| add | Src,Dest |
|------|-----------|
| sub | Src,Dest |
| imul | Src,Dest |
| shr | Src,Dest |
| sar | Src, Dest |
| shl | Src,Dest |
| sal | Src, Dest |
| shr | Src,Dest |
| xor | Src,Dest |
| and | Src,Dest |
| or | Src,Dest |
| mul | Src,Dest |
| imul | Src,Dest |
| div | Src,Dest |
| idiv | Src,Dest |

Arithmetic/Logical operations - 1 operand

| inc | Dest |
|-----|------|
| del | Dest |
| neg | Dest |
| not | Dest |

Setting Condition Codes Explicitly – used for control flow

| cmp Src2,Src1 | sets flags based on value of Src2 – Src1, |
|----------------|--|
| | discards result |
| test Src2,Src1 | sets flags based on a & b, discards result |

Control Flow

Conditional jump instructions in X86 implement the following high-level constructs:

- if (condition) then {...} else {...}
- while (condition) {...}
- do {...} while (condition)
- for (initialization; condition; iterative) {...}

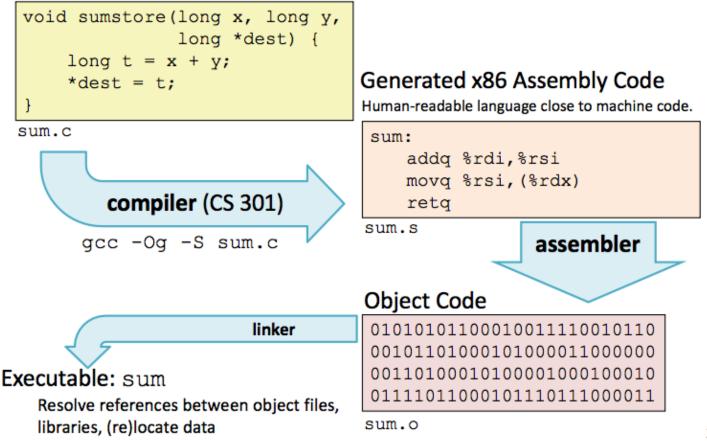
Unconditional jumps are used for high-level constructs such as:

- break
- continue

| X | Condition | Description | |
|-----|----------------|---------------------------|--|
| jmp | 1 | Unconditional | |
| je | ZF | Equal / Zero | |
| jne | ~ZF | Not Equal / Not Zero | |
| js | SF | Negative | |
| jns | ~SF | Nonnegative | |
| jg | ~ (SF^OF) &~ZF | Greater (Signed) | |
| jge | ~ (SF^OF) | Greater or Equal (Signed) | |
| j1 | (SF^OF) | Less (Signed) | |
| jle | (SF^OF) ZF | Less or Equal (Signed) | |
| ja | ~CF6~ZF | Above (unsigned) | |
| jЪ | CF | Below (unsigned) | |

Turning C into Machine Code





- X86 instructions can be in different order from C code
- Some C expressions require multiple X86 instructions
- Some X86 instructions can cover multiple C expressions
- Compiler optimization can do some surprising things!
- Local or temporary variables can be stored in registers or on the stack

Function Calling Conventions

- Arguments for functions are stored in registers, in the following order: arg1 arg6: %rdi,%rsi,%rdx,%rcx,%r8,%r9
- If there are more than 6 parameters for a function, the rest of the arguments are stored on the stack before the function is called
- Return value from function is always in %rax

The compiler will use only part of a register if the value stored there will fit in less than 64 bits (8 bytes). This is an optimization that makes instructions a bit shorter.

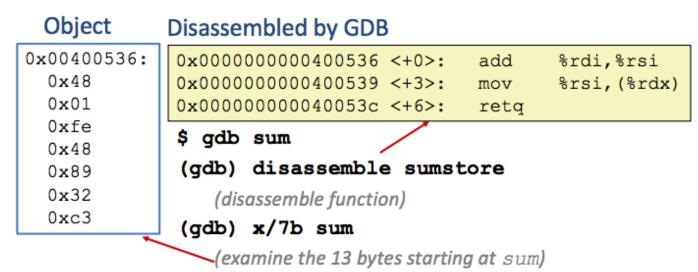
So, in the code, you may see register names of the following form, all of which refer to %rax:

%**r**ax = 8 byte value %**e**ax = 4 byte value %ax = 2 byte value %al = 1 byte value

Tools

Tools can be used to examine bytes of object code (executable program) and reconstruct (reverse engineer) the assembly source.

gdb – disassembles an executable file into the associated assembly language representation, and provides tools for memory and register examination, single step execution, breakpoints, etc.



objdump

can also be used to disassemble and display information

\$ objdump –t p

Prints out the program's symbol table. The symbol table includes the names of all functions and global variables, the names of all the functions the called, and their addresses.

\$ objdump -d p

Object Code

Disassembled version

0x401040 <sum>: 00401040 <_sum>: 0x55 0: 55 push %ebp mov %esp,%ebp 0x89 1: 89 e5 8b 45 0c mov 0xc(%ebp),%eax 0xe5 3: 03 45 08 add 0x8(%ebp),%eax 0x8b6: 0x45 9: 89 ec mov %ebp,%esp 5d pop %ebp 0x0cb: 0x03 **c**3 c: ret 0x45 0x08 0x89 0xec 0x5d0xc3

strings

\$ strings –t x p

Displays the printable strings in your program.