



CS 240

Foundations of Computer Systems

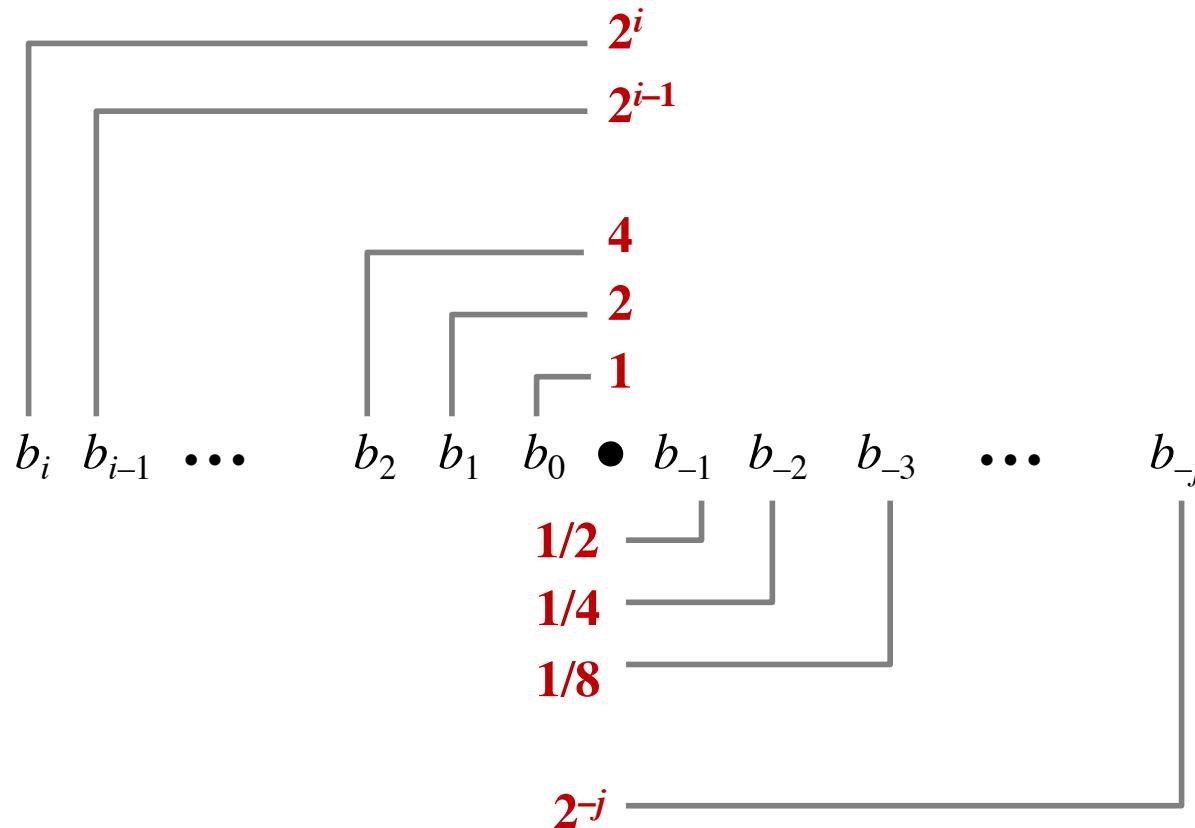


Floating Point Representation

Fractional binary numbers
IEEE floating-point standard
Floating-point operations and rounding
Lessons for programmers

Many more details we will skip (it's a 58-page standard...)
See CSAPP 2.4 for more detail.

Fractional Binary Numbers



$$\sum_{k=-j}^i b_k \cdot 2^k$$

Fractional Binary Numbers

Value	Representation
-------	----------------

5 and 3/4	
-----------	--

2 and 7/8	
-----------	--

47/64	
-------	--

Observations

Shift left =

Shift right =

Numbers of the form $0.\underset{\text{1}}{1}\underset{\text{1}}{1}\underset{\text{1}}{1}\underset{\text{1}}{1}\underset{\text{1}}{1}\dots_2$ are...?

Limitations:

Exact representation possible when?

$$1/3 = 0.\underset{\text{1}}{3}\underset{\text{1}}{3}\underset{\text{1}}{3}\dots_{10} = 0.$$

Fixed-Point Representation

Implied binary point.

b₇ b₆ b₅ b₄ b₃ [.] b₂ b₁ b₀

b₇ b₆ b₅ b₄ b₃ b₂ b₁ b₀ [.]

range: difference between largest and smallest representable numbers

precision: smallest difference between any two representable numbers

fixed point = fixed range, fixed precision

IEEE Floating Point Standard 754

IEEE = Institute of Electrical and Electronics Engineers

Numerical form:

$$V_{10} = (-1)^s * M * 2^E$$

Sign bit s determines whether number is negative or positive

Significand (mantissa) M usually a fractional value in range [1.0,2.0)

Exponent E weights value by a (-/+) power of two

Analogous to scientific notation

Representation:

MSB s = sign bit s

exp field encodes E (but is *not equal* to E)

frac field encodes M (but is *not equal* to M)



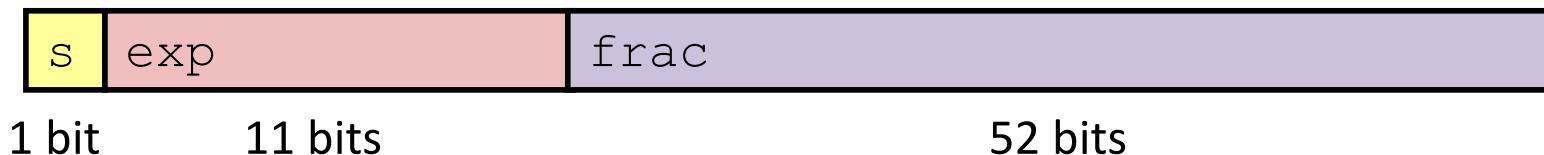
Numerically well-behaved, but hard to make fast in hardware

Precisions

Single precision (`float`): 32 bits



Double precision (`double`): 64 bits



Finite representation of infinite range...

Three kinds of values

$$V = (-1)^S * M * 2^E$$



1. Normalized: $M = 1.xxxxx\dots$

As in scientific notation: $0.011 \times 2^5 = 1.1 \times 2^3$

Representation advantage?

2. Denormalized, near zero: $M = 0.xxxxx\dots$, smallest E

Evenly space near zero.

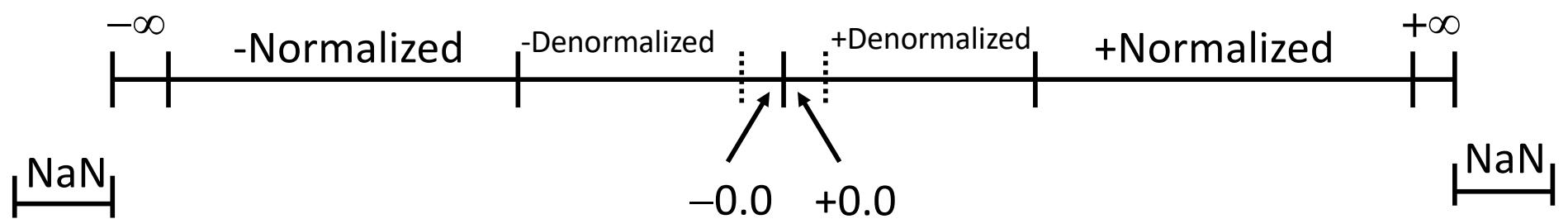
3. Special values:

0.0: $s = 0$ $\text{exp} = 00\dots0$ $\text{frac} = 00\dots0$

+inf, -inf: $\text{exp} = 11\dots1$ $\text{frac} = 00\dots0$
division by 0.0

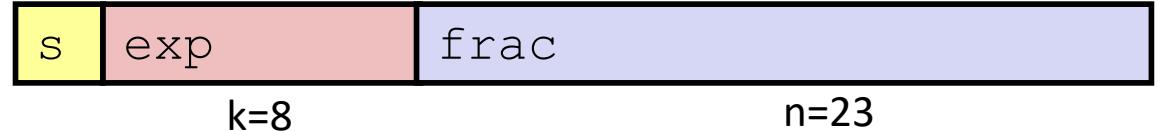
NaN (“Not a Number”): $\text{exp} = 11\dots1$ $\text{frac} \neq 00\dots0$
sqrt(-1), $\infty - \infty$, $\infty * 0$, etc.

Value distribution



Normalized values, with float example

$$V = (-1)^S * M * 2^E$$



Value: float f = 12345.0;

$$\begin{aligned}12345_{10} &= 11000000111001_2 \\&= 1.1000000111001_2 \times 2^{13} \text{ (normalized form)}\end{aligned}$$

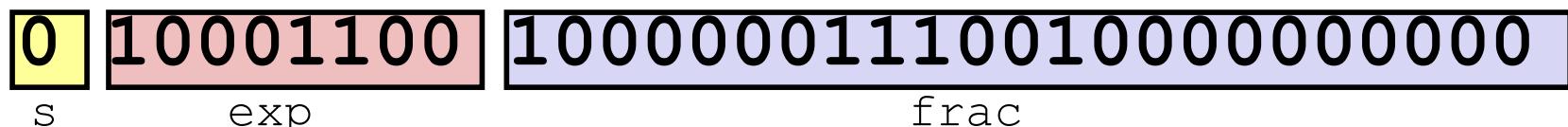
Significand:

$$\begin{aligned}M &= 1.\underline{1000000111001}_2 \\frac &= \underline{1000000111001}0000000000_2\end{aligned}$$

Exponent: $E = exp - Bias \rightarrow exp = E + Bias$

$$\begin{aligned}E &= 13 \\Bias &= 127 = 2^7 - 1 = 2^{k-1} - 1 && \text{Splits exponents roughly -/+} \\exp &= 140 = 10001100_2\end{aligned}$$

Result:



Denormalized Values: near zero

"Near zero": $\text{exp} = 000\dots0$

Exponent:

$$E = 1 + \text{exp} - \text{Bias} = 1 - \text{Bias} \quad \text{not: } \text{exp} - \text{Bias}$$

Significand: leading zero

$$M = 0.\text{xxxx...x}_2$$

$$\text{frac} = \text{xxx...x}$$

Cases:

$$\text{exp} = 000\dots0, \text{frac} = 000\dots0 \qquad 0.0, -0.0$$

$$\text{exp} = 000\dots0, \text{frac} \neq 000\dots0$$

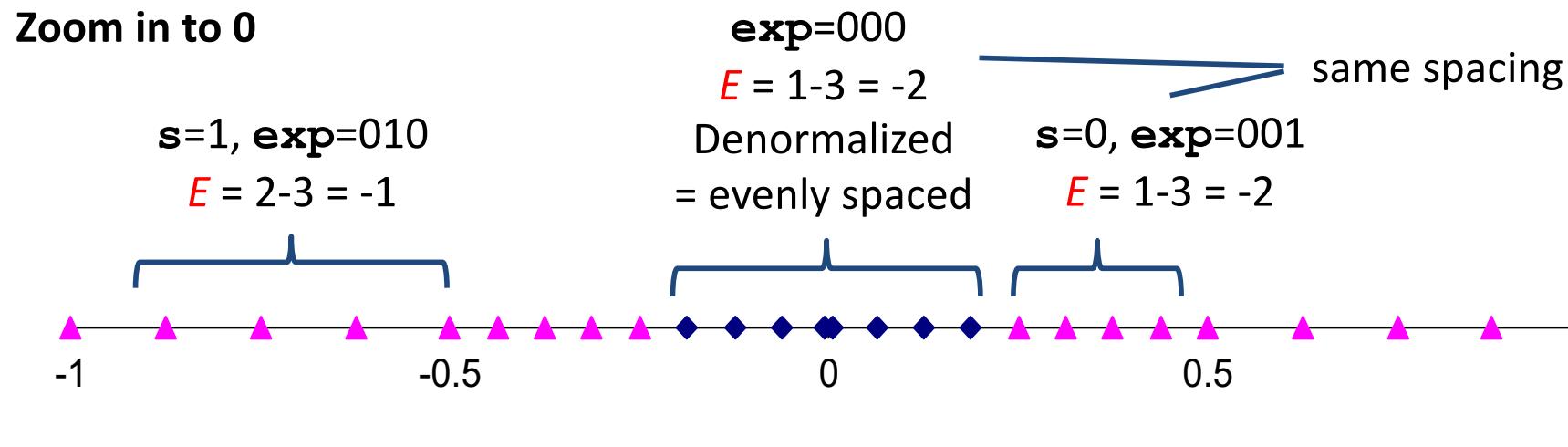
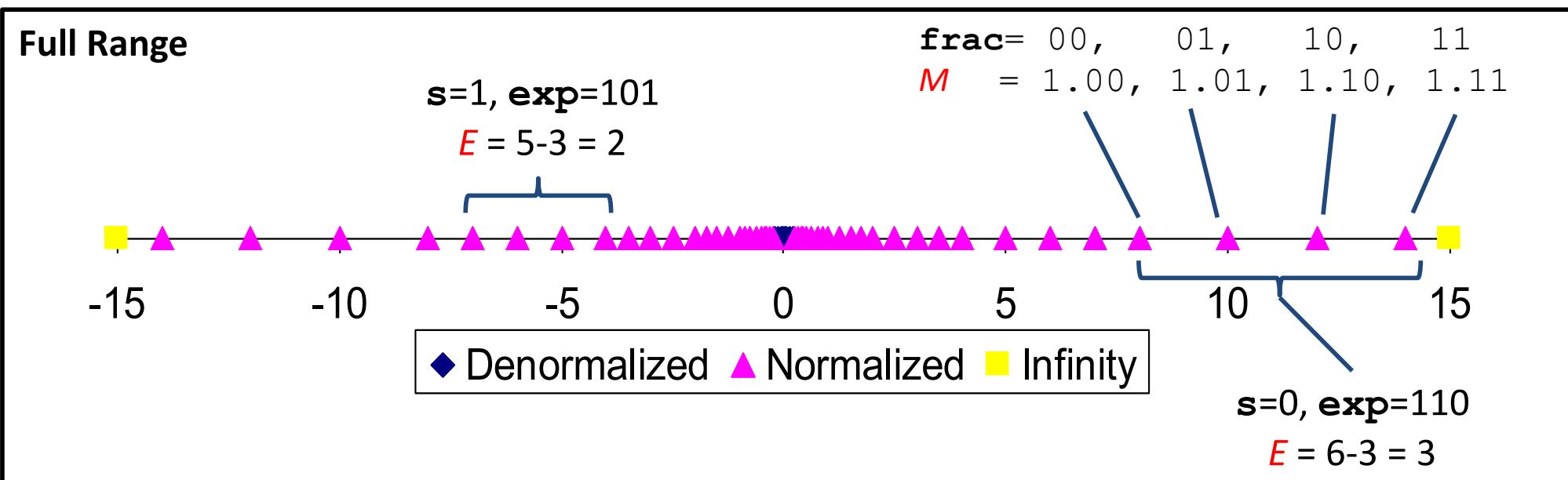
Value distribution example

6-bit IEEE-like format



$$\text{Bias} = 2^{3-1} - 1 = 3$$

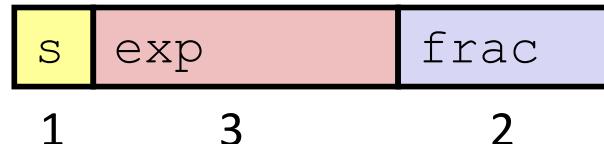
1 3 2



Try to represent 3.14, 6-bit example

6-bit IEEE-like format

$$\text{Bias} = 2^{3-1} - 1 = 3$$



Value: 3.14 ;

$$3.14 = 11.0010\ 0011\ 1101\ 0111\ 0000\ 1010\ 000\dots$$

$$= 1.1001\ 0001\ 1110\ 1011\ 1000\ 0101\ 0000\dots_2 \times 2^1 \quad (\text{normalized form})$$

Significand:

$$\textcolor{red}{M} = 1.10010001111010111011100001010000\dots_2$$

$$\text{frac} = \underline{10}_2$$

Exponent:

$$\textcolor{red}{E} = 1 \quad \text{Bias} = 3 \quad \text{exp} = 4 = 100_2$$

Result:

$$0\ 100\ 10 = 1.10_2 \times 2^1 = 3 \quad \text{next highest?}$$

Floating Point Arithmetic*

$$V = (-1)^S * M * 2^E$$



```
double x = . . . , y = . . . ;
```

```
double z = x + y;
```



1. Compute exact result.
2. Fix/Round, roughly:

Adjust M to fit in [1.0, 2.0)...

If $M \geq 2.0$: shift M right, increment E

If $M < 1.0$: shift M left by k , decrement E by k

Overflow to infinity if E is too wide for **exp**

Round* M if too wide for **frac**.

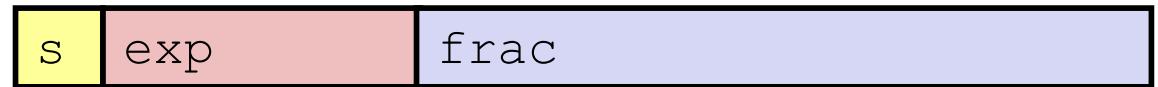
Underflow if nearest representable value is 0.

...

*complicated...

Lessons for programmers

$$v = (-1)^S \cdot M \cdot 2^E$$



`float` \neq real number \neq `double`

Rounding breaks associativity and other properties.

```
double a = ..., b = ...;
```

✗ ...

```
if (a == b) ...
```

```
if (abs(a - b) < epsilon) ...
```