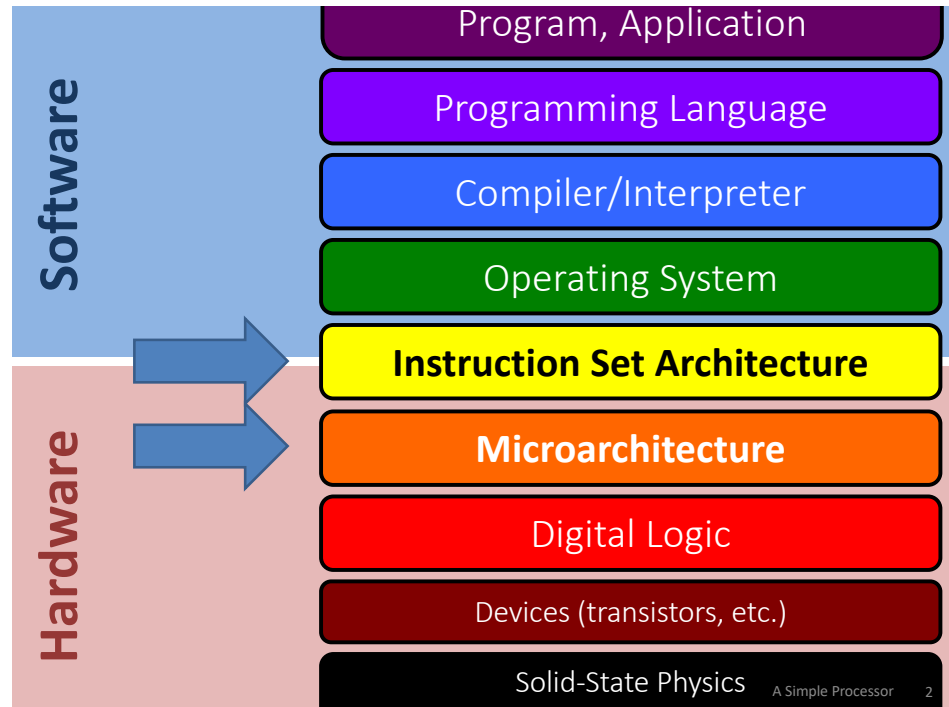


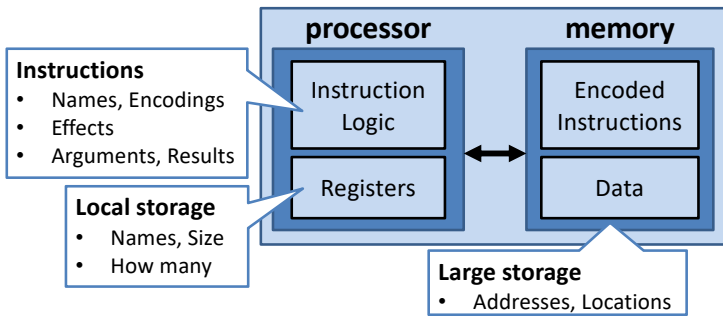


A Simple Processor

1. A simple Instruction Set Architecture
2. A simple microarchitecture (implementation):
Data Path and Control Logic



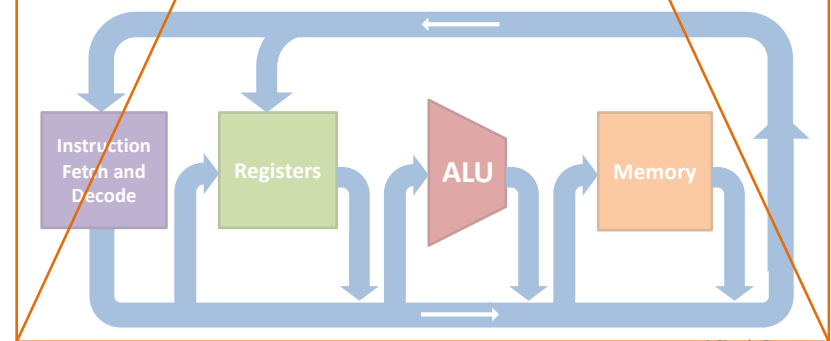
Instruction Set Architecture (HW/SW Interface)



Computer

Computer

Microarchitecture (Implementation of ISA)



HW ISA

An example made-up instruction set architecture

Word size = 16 bits

- Register size = 16 bits.
- ALU computes on 16-bit values.

Memory is byte-addressable, accesses full words (byte pairs).

16 registers: R0 - R15

- R0 always holds hardcoded 0
- R1 always holds hardcoded 1
- R2 – R15: general purpose

Instructions are 1 word in size.

Separate *instruction memory*.

Program Counter (PC) register

- holds address of next instruction to execute.

Address	Contents
0	First instruction, low-order byte
1	First instruction, high-order byte
2	Second instruction, low-order byte
...	...

R: Register File

Reg	Contents	Reg	Contents
R0	0x0000	R8	
R1	0x0001	R9	
R2		R10	
R3		R11	
R4		R12	
R5		R13	
R6		R14	
R7		R15	

M: Data Memory

Address	Contents
0x0 – 0x1	
0x2 – 0x3	
0x4 – 0x5	
0x6 – 0x7	
0x8 – 0x9	
0xA – 0xB	
0xC – 0xD	
...	

Program Counter

PC

Processor Loop

1. $ins \leftarrow IM[PC]$
2. $PC \leftarrow PC + 2$
3. Do ins

IM: Instruction Memory

Address	Contents
0x0 – 0x1	
0x2 – 0x3	
0x4 – 0x5	
0x6 – 0x7	
0x8 – 0x9	
...	

HW ISA Abstract Machine

HW ISA

Instructions

MSB 16-bit Encoding LSB

Assembly Syntax	Meaning	Opcode	Rs	Rt	Rd
ADD R_s, R_t, R_d	$R[d] \leftarrow R[s] + R[t]$	0010	s	t	d
SUB R_s, R_t, R_d	$R[d] \leftarrow R[s] - R[t]$	0011	s	t	d
AND R_s, R_t, R_d	$R[d] \leftarrow R[s] \& R[t]$	0100	s	t	d
OR R_s, R_t, R_d	$R[d] \leftarrow R[s] R[t]$	0101	s	t	d
LW $R_t, offset(R_s)$	$R[t] \leftarrow M[R[s] + offset]$	0000	s	t	$offset$
SW $R_t, offset(R_s)$	$M[R[s] + offset] \leftarrow R[t]$	0001	s	t	$offset$
BEQ $R_s, R_t, offset$	If $R[s] == R[t]$ then $PC \leftarrow PC + offset * 2$	0111	s	t	$offset$
JMP $offset$	$PC \leftarrow offset * 2$	1000	o	f	f

(R = register file, M = memory)

JMP offset is unsigned
All others are signed

R: Register File

Reg	Contents	Reg	Contents
R0	0x0000	R8	
R1	0x0001	R9	
R2		R10	
R3		R11	
R4		R12	
R5		R13	
R6		R14	
R7		R15	

M: Data Memory

Address	Contents
0x0 – 0x1	0x46 0x12
0x2 – 0x3	0xA9 0x56
0x4 – 0x5	
0x6 – 0x7	
0x8 – 0x9	
0xA – 0xB	
0xC – 0xD	
...	

Program Counter

PC

Processor Loop

1. $ins \leftarrow IM[PC]$
2. $PC \leftarrow PC + 2$
3. Do ins

IM: Instruction Memory

Address	Contents
0x0 – 0x1	LW R3, 0(R0)
0x2 – 0x3	LW R4, 2(R0)
0x4 – 0x5	AND R3, R4, R5
0x6 – 0x7	SW R5, 4(R0)
0x8 – 0x9	
0xA – 0xB	

HW ISA Abstract Machine

R: Register File

Reg	Contents	Reg	Contents
R0	0x0000	R8	
R1	0x0001	R9	2
R2		R10	3
R3		R11	
R4		R12	
R5		R13	
R6		R14	
R7		R15	

M: Data Memory

Address	Contents
0x0 – 0x1	
0x2 – 0x3	
0x4 – 0x5	
0x6 – 0x7	
0x8 – 0x9	
0xA – 0xB	
0xC – 0xD	
...	

Program Counter

PC

Processor Loop

1. $ins \leftarrow IM[PC]$
2. $PC \leftarrow PC + 2$
3. Do ins

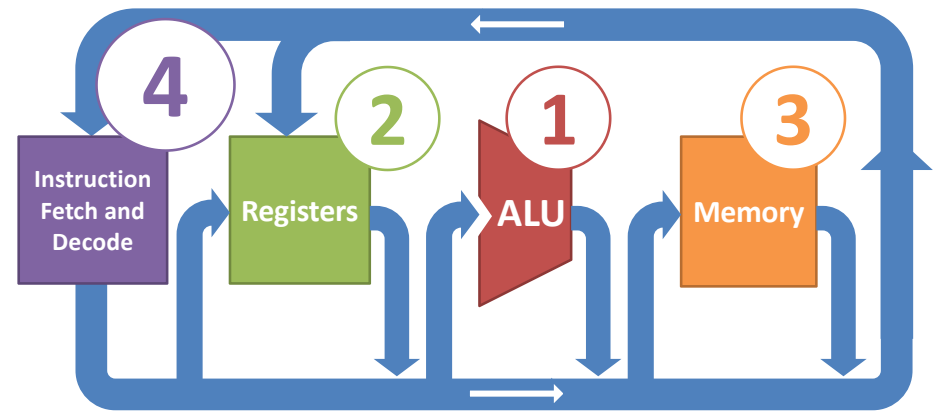
IM: Instruction Memory

Address	Contents
0x0 – 0x1	SUB R8, R8, R8
0x2 – 0x3	BEQ R9, R0, 3
0x4 – 0x5	ADD R10, R8, R8
0x6 – 0x7	SUB R9, R1, R9
0x8 – 0x9	JMP 1
0xA – 0xB	HALT

HW ISA Abstract Machine

A Simple Processor 9

HW ARCH microarchitecture



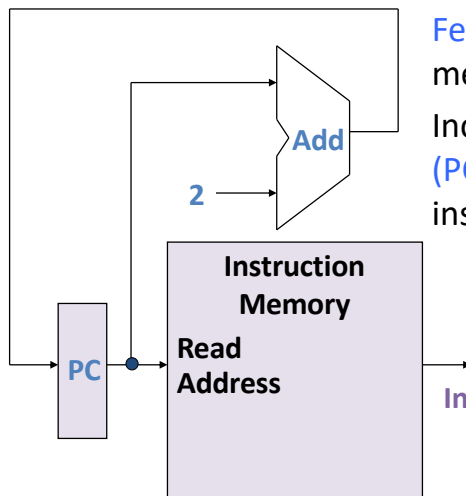
One possible hardware implementation of the HW ISA

A Simple Processor 11

Instruction Fetch

Processor Loop

1. $ins \leftarrow IM[PC]$
2. $PC \leftarrow PC + 2$
3. Do ins



Fetch instruction from memory.

Increment program counter (PC) to point to the next instruction.

A Simple Processor 12

Arithmetic Instructions

ADD R3, R6, R8

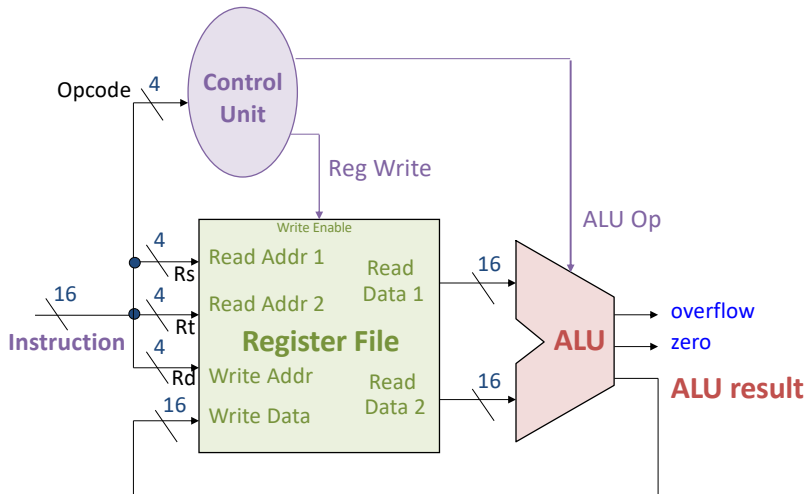
Opcode	Rs	Rt	Rd
0010	0011	0110	1000

16-bit Encoding

Instruction	Meaning	Opcode	Rs	Rt	Rd
ADD R_s, R_t, R_d	$R[d] \leftarrow R[s] + R[t]$	0010	0-15	0-15	0-15
SUB R_s, R_t, R_d	$R[d] \leftarrow R[s] - R[t]$	0011	0-15	0-15	0-15
AND R_s, R_t, R_d	$R[d] \leftarrow R[s] \& R[t]$	0100	0-15	0-15	0-15
OR R_s, R_t, R_d	$R_d \leftarrow R[s] R[t]$	0101	0-15	0-15	0-15
...					

A Simple Processor 13

Instruction Decode, Register Access, ALU



Memory Instructions

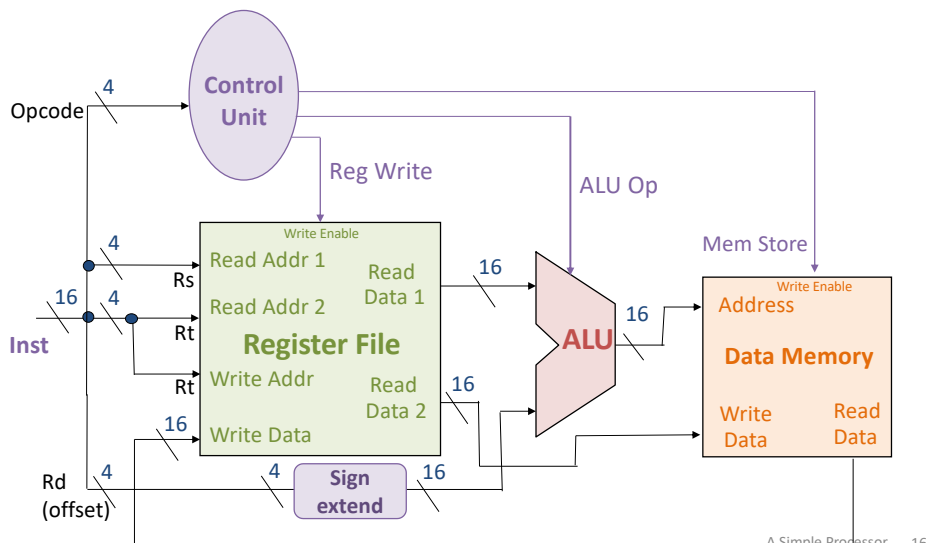
SW R6, -8(R3)

Opcode	Rs	Rt	Rd
0001	0011	0110	1000

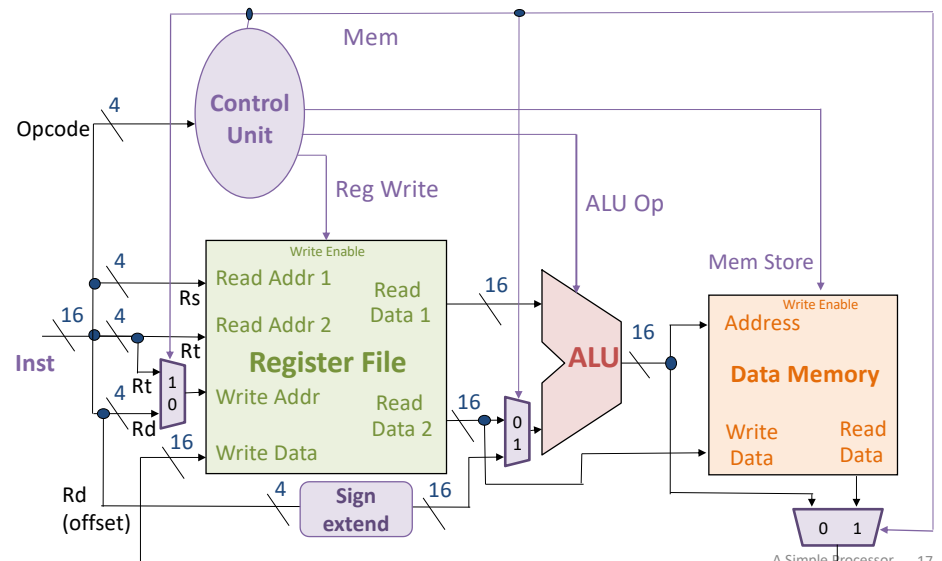
Instruction	Meaning	Op	Rs	Rt	Rd
LW Rt, offset(Rs)	$R[t] \leftarrow Mem[R[s] + offset]$	0000	0-15	0-15	offset
SW Rt, offset(Rs)	$Mem[R[s] + offset] \leftarrow R[t]$	0001	0-15	0-15	offset
...					

Memory access

How can we support arithmetic and memory instructions?
What's shared?



Choose with MUXs



Control-flow Instructions

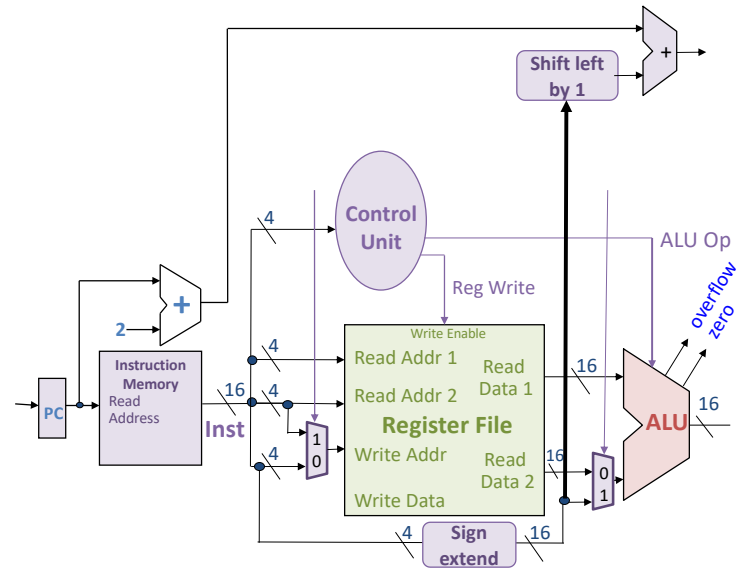
BEQ R1, R2, -2

Op	Rs	Rt	Rd
0111	0001	0010	1110

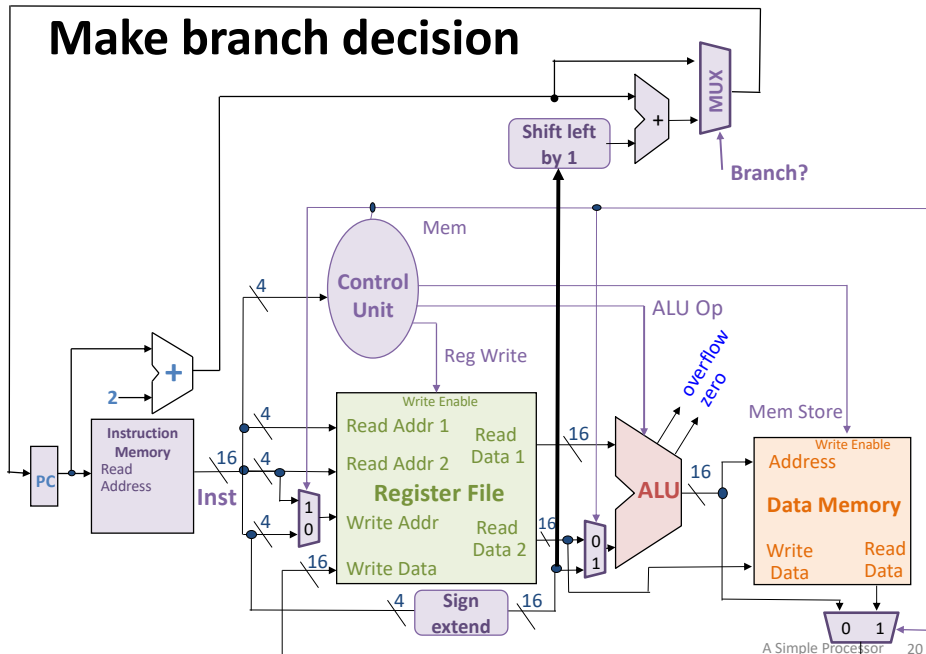
16-bit Encoding

Instruction	Meaning	Op	Rs	Rt	Rd
BEQ Rs, Rt, offset	If $R[s] == R[t]$ then $PC \leftarrow PC + offset * 2$	0111	0-15	0-15	offset
...					

Compute branch target



Make branch decision



HW ARCH: not the only implementation

Single-cycle architecture

- Simple, "easily" fits on a slide (and in your head).
- One instruction takes one clock cycle.
- Slowest instruction determines minimum clock cycle.
- Inefficient.

Could it be better?

- Performance, energy, debugging, security, reconfigurability, ...
- Pipelining
- OoO: out-of-order execution
- SIMD: single instruction multiple data
- Caching
- Microcode vs. direct hardware implementation
- ... enormous, interesting design space of **Computer Architecture**