



CS 240 Stage 2

Hardware-Software Interface

Memory addressing, C language, pointers

Assertions, debugging

Machine code, assembly language, program translation

Control flow

Procedures, stacks

Data layout, security, linking and loading



Programming with Memory

the memory model
pointers and arrays in C

Software

Program, Application

Programming Language

Compiler/Interpreter

Operating System

Instruction Set Architecture

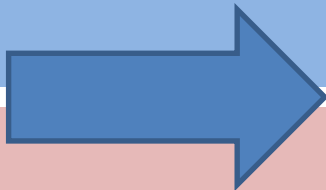
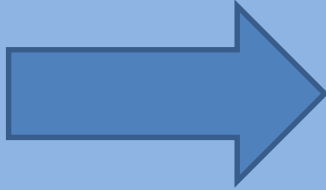
Microarchitecture

Digital Logic

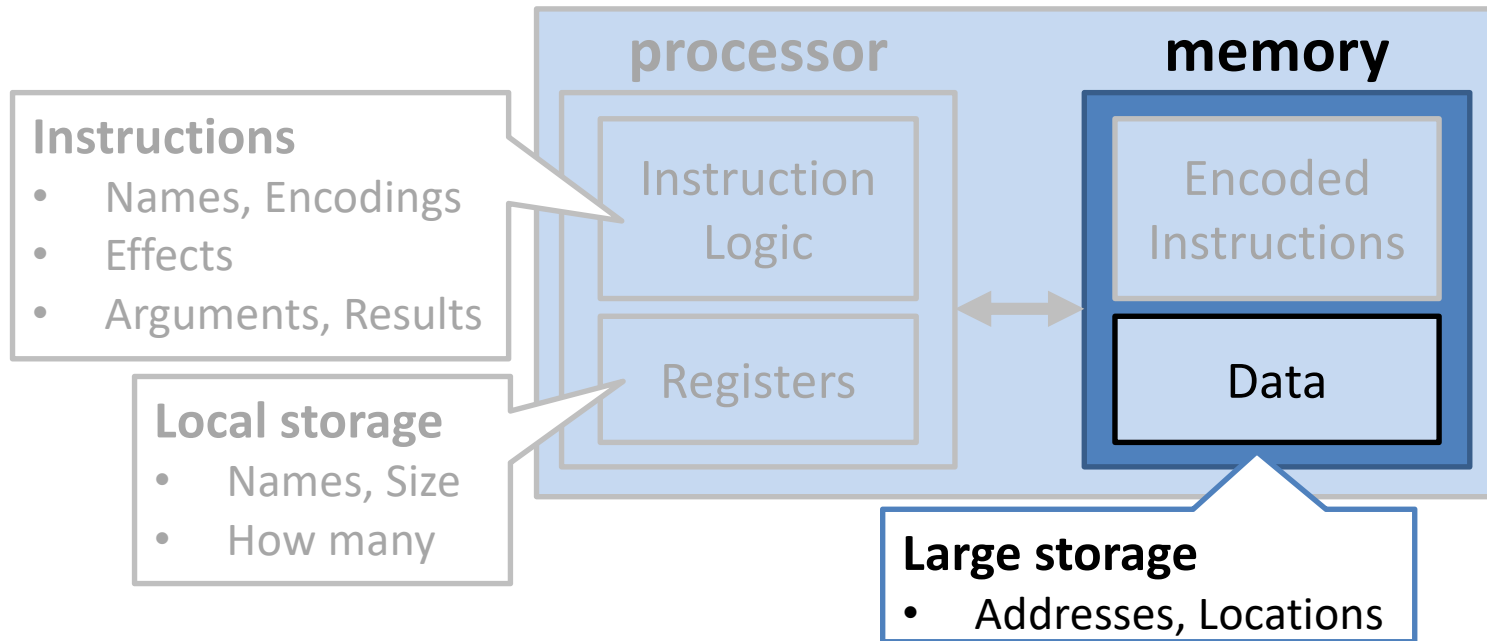
Devices (transistors, etc.)

Solid-State Physics

Hardware

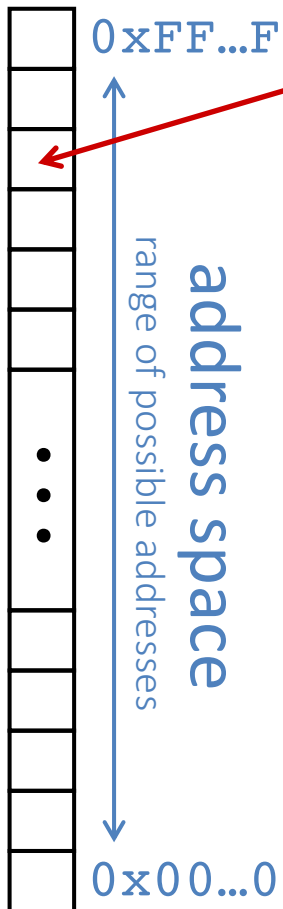


Instruction Set Architecture (HW/SW Interface)



Computer

Byte-addressable memory = mutable byte array



Location / cell = element

- Identified by unique numerical **address**
- Holds one byte

Address = index

- Unsigned number
- Represented by one word
- Computable and storable as a value

Operations:

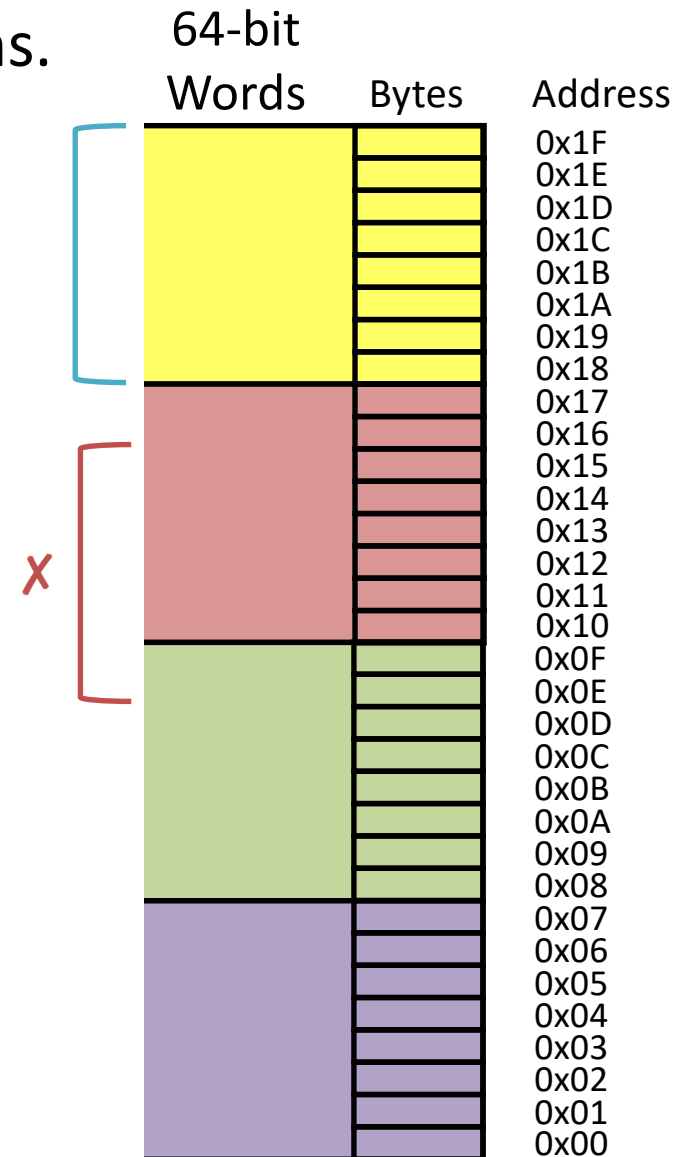
- **Load:** read contents at given address
- **Store:** write contents at given address

Multi-byte values in memory

Store across contiguous byte locations.

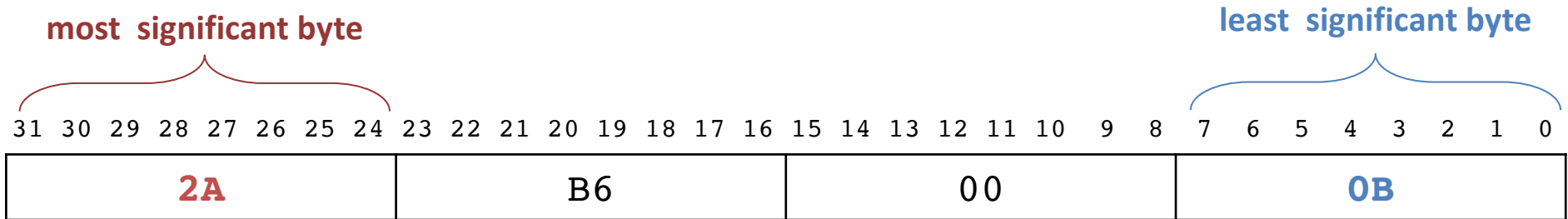
Alignment (Why?)

Bit order within byte always same.
Byte ordering within larger value?



Endianness

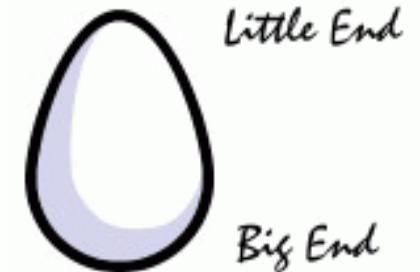
In what order are the individual bytes of a multi-byte value stored in memory?



Address	Contents
03	2A
02	B6
01	00
00	0B

Little Endian: least significant byte first

- low order byte at low address
- high order byte at high address
- used by x86, ...



Address	Contents
03	0B
02	00
01	B6
00	2A

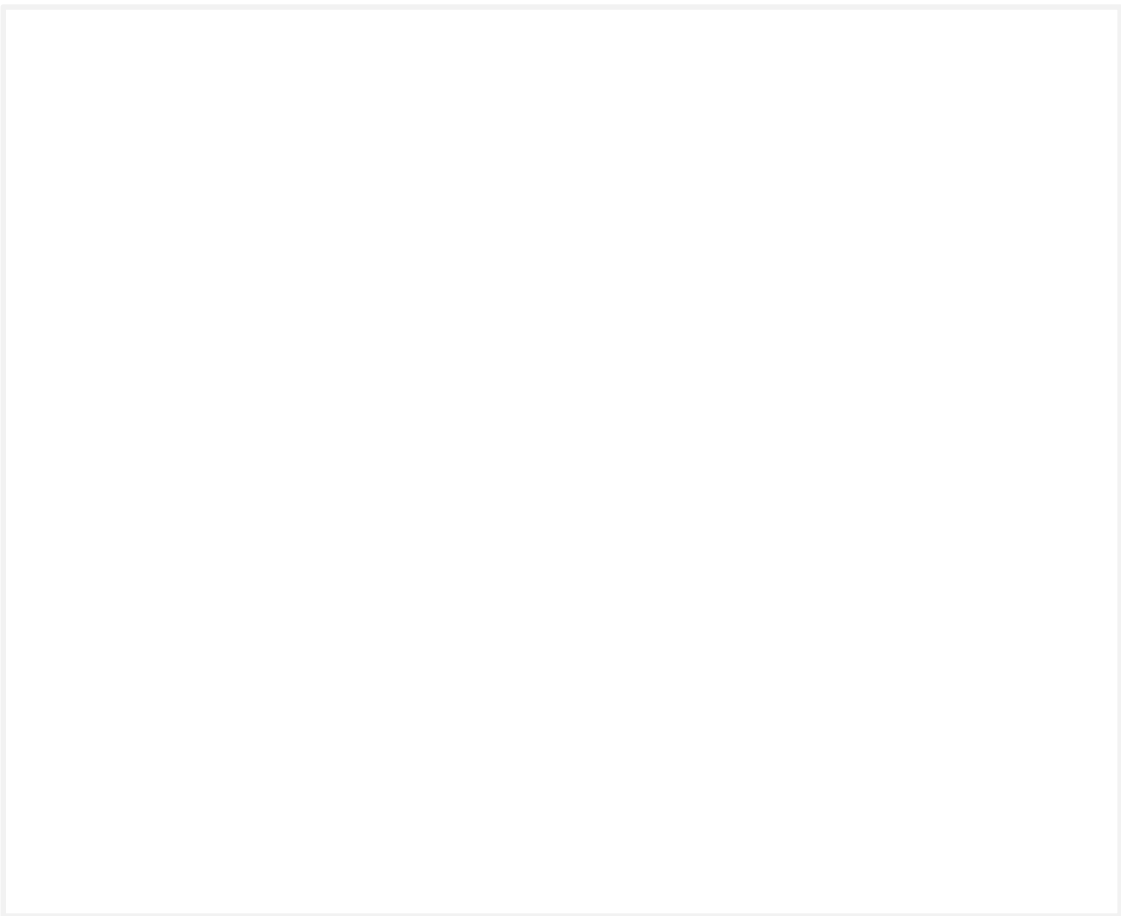
Big Endian: most significant byte first

- high order byte at low address
- low order byte at high address
- used by networks, SPARC, ...

Data, addresses, and pointers

address = index of a location in memory

pointer = a reference to a location in memory,
represented as an address stored as data



				0x24
00	00	00	F0	0x20
				0x1C
				0x18
				0x14
00	00	00	0C	0x10
				0x0C
00	00	00	20	0x08
				0x04
00	00	00	08	0x00

+3 +2 +1 +0

memory drawn as 32-bit values,
little endian order

C: Variables are locations

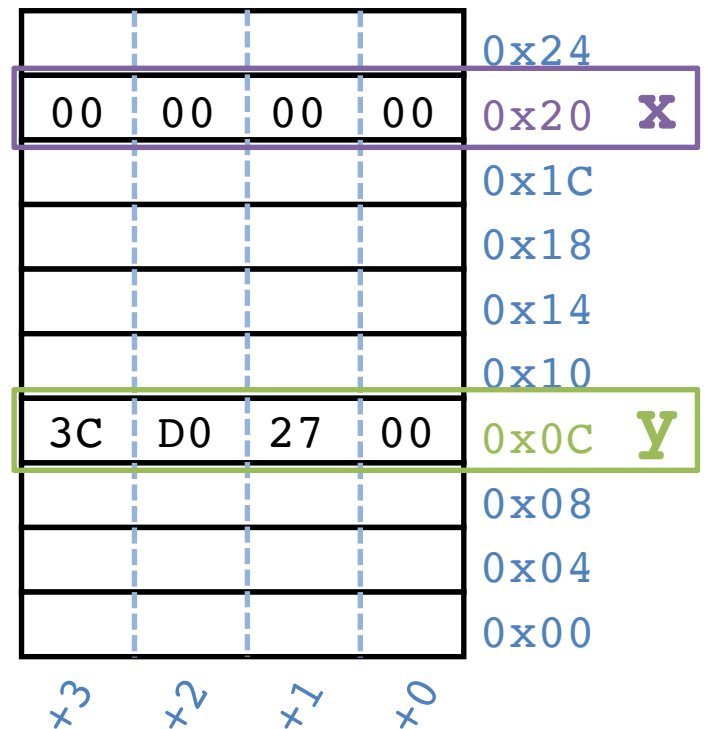
Compiler maps variable name → location.

Declarations do not initialize!

```
int x; // x @ 0x20
int y; // y @ 0x0C

x = 0; // store 0 @ 0x20

// store 0x3CD02700 @ 0x0C
y = 0x3CD02700;
```



C: Variables are locations

Compiler maps variable name → location.

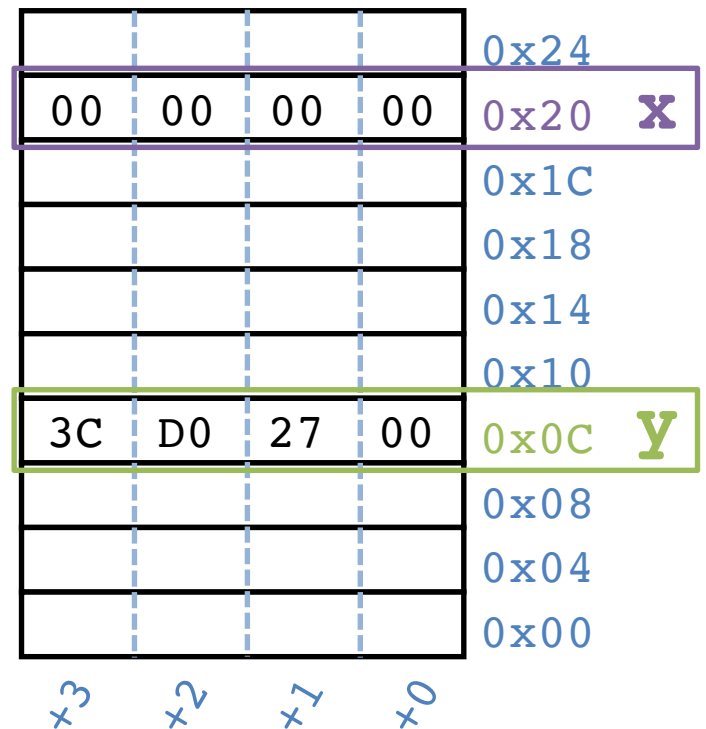
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int x; // x @ 0x20
int y; // y @ 0x0C

x = 0; // store 0 @ 0x20

// store 0x3CD02700 @ 0x0C
y = 0x3CD02700;

// 1. load the contents @ 0x0C
// 2. add 3
// 3. store sum @ 0x20
x = y + 3;
```



C: Variables are locations

Compiler maps variable name → location.

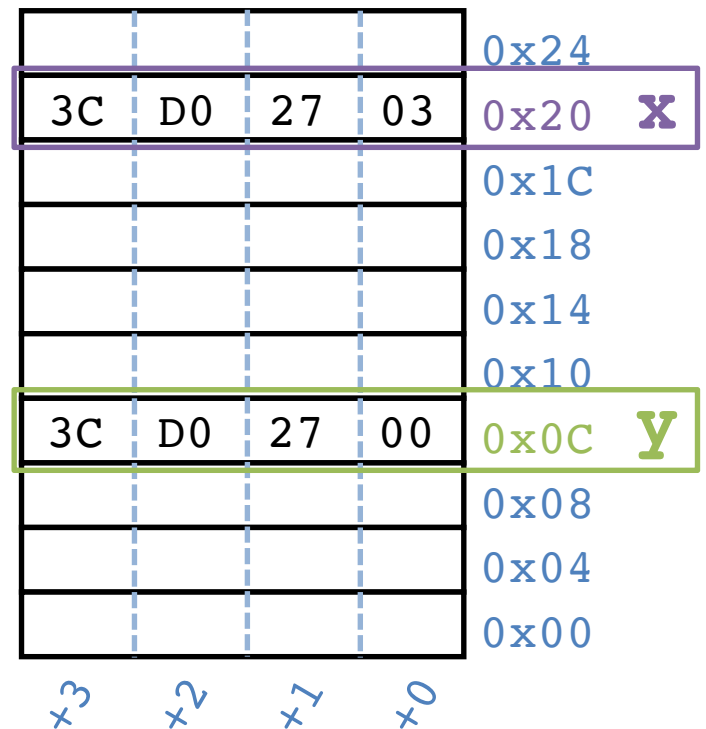
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// 3. store sum @ 0x20
x = y + 3;
```



C: Pointer operations and types

address = index of a location in memory

pointer = a reference to a location in memory,
represented as an address stored as data

Expressions using addresses and pointers:

&____ **address of** the memory location representing ____
a.k.a. "reference to ____"

*____ **contents at** the memory address given by ____
a.k.a. "dereference ____"

Pointer types:

____* address of a memory location holding a ____
a.k.a. "a reference to a ____"

C: Types determine sizes

Sizes of data types (in bytes)

Java Data Type	C Data Type	32-bit word	64-bit word
boolean	<i>bool</i>	1	1
byte	char	1	1
char		2	2
short	short int	2	2
int	int	4	4
float	float	4	4
	long int	4	8
double	double	8	8
long	long long	8	8
	long double	8	16
(reference)	(pointer) *	4	8

address size = word size

& = address of
*** = contents at

C: Pointer example

```
int* p;
```

Declare a variable, *p*

that will hold the address of a memory location holding an int

```
int x = 5;  
int y = 2;
```

Declare two variables, *x* and *y*, that hold ints, and store 5 and 2 in them, respectively.

```
p = &x;
```

Take the address of the memory location

representing *x*

... and store it in the memory location representing *p*.
Now, "*p* points to *x*."

Add 1 to the contents of memory at the address

```
y = 1 + *p;
```

given by the contents of the memory location representing *p*

... and store it in the memory location representing *y*.

C: Pointer type syntax

Spaces between base type, *, and variable name mostly do not matter.

The following are **equivalent**:

```
int* ptr;
```

I prefer this

I see: "The variable **ptr** holds an **address of an int** in memory."

```
int * ptr;
```

```
int *ptr;
```

more common C style

Looks like: "Dereferencing the variable **ptr** will yield an **int**."

Or "The **memory location** where the variable **ptr** points holds an **int**."

Caveat: do not declare multiple variables unless using the last form.

```
int* a, b; means int *a, b; means int* a; int b;
```


C: Arrays

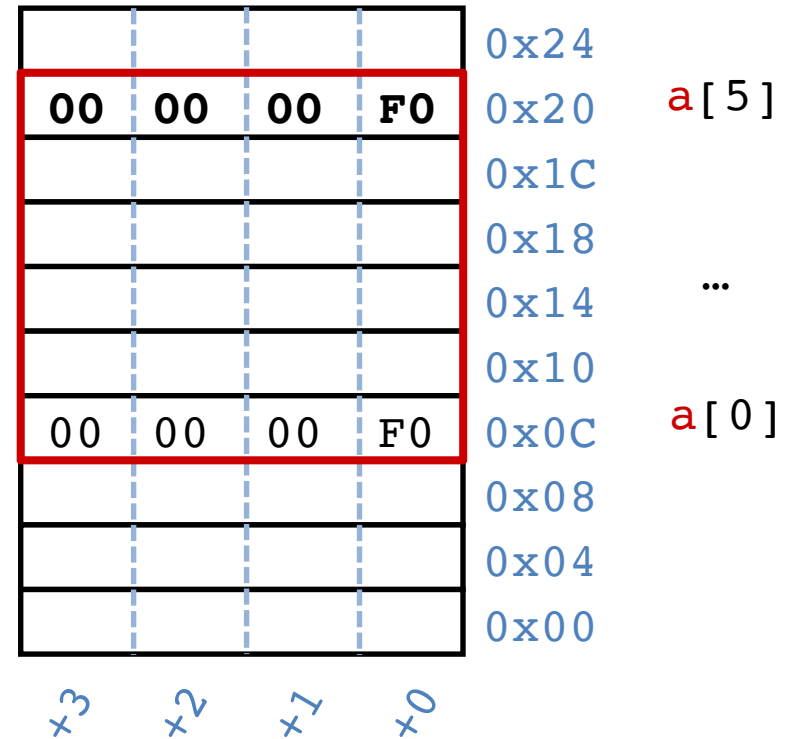
Declaration: `int a[6];`

Indexing: `a[0] = 0xf0;`
`a[5] = a[0];`

Arrays are adjacent memory locations storing the same type of data.

`a` is a name for the array's base address, can be used as an *immutable* pointer.

Address of `a[i]` is base address `a` plus `i` times element size in bytes.



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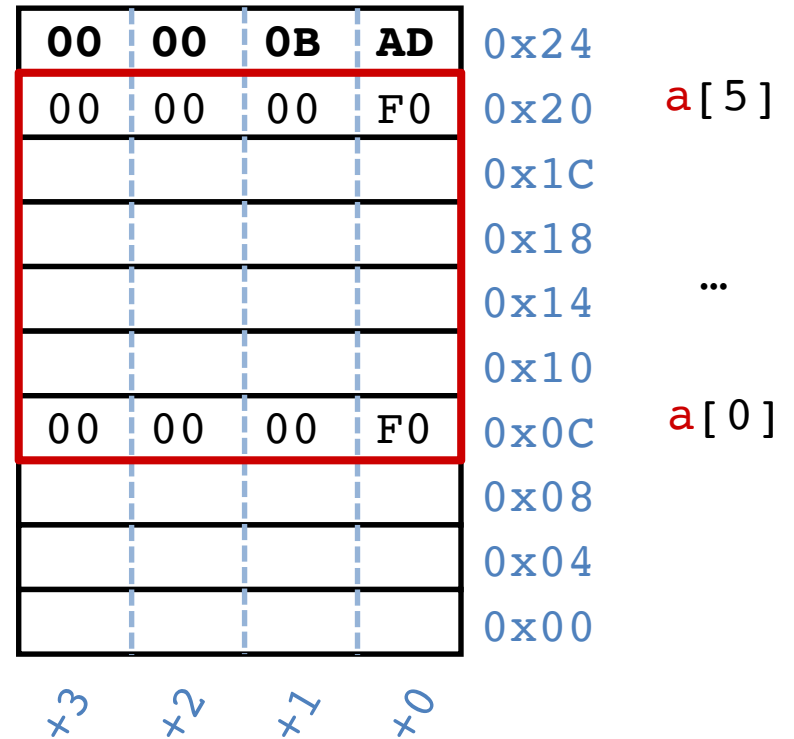
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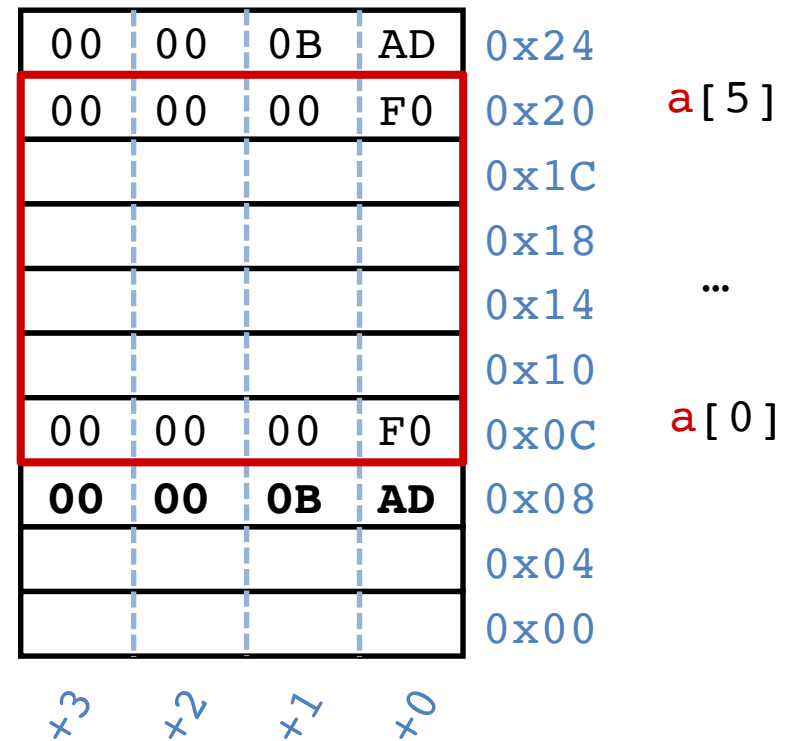
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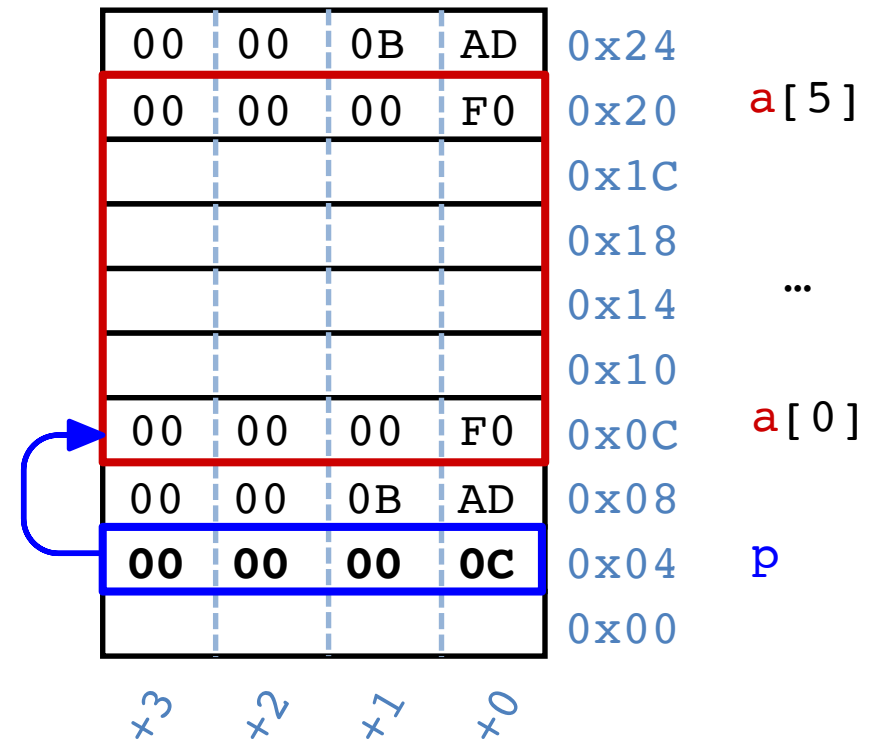
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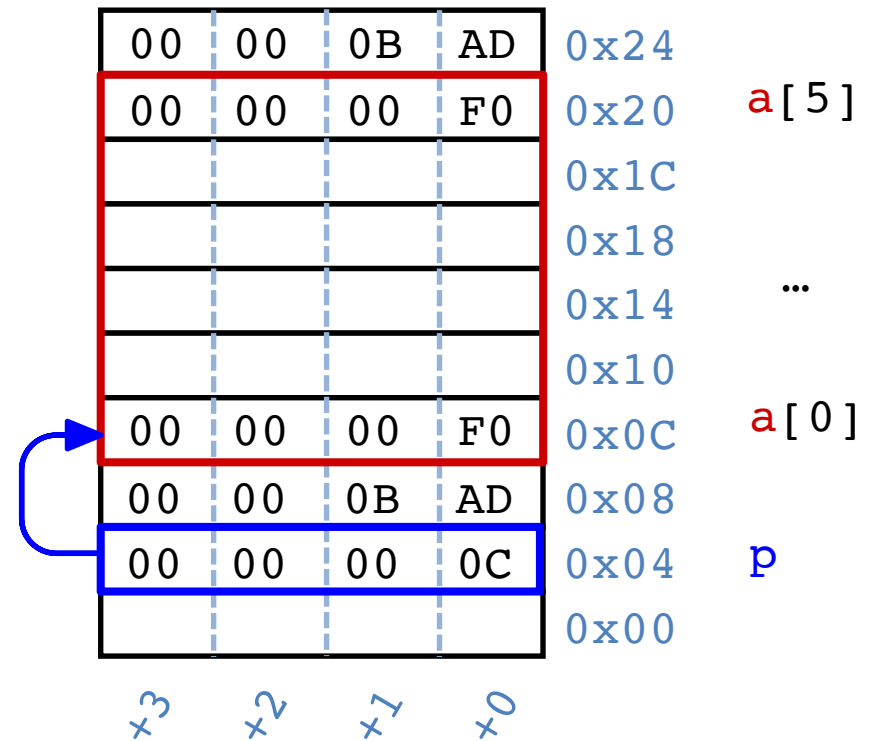
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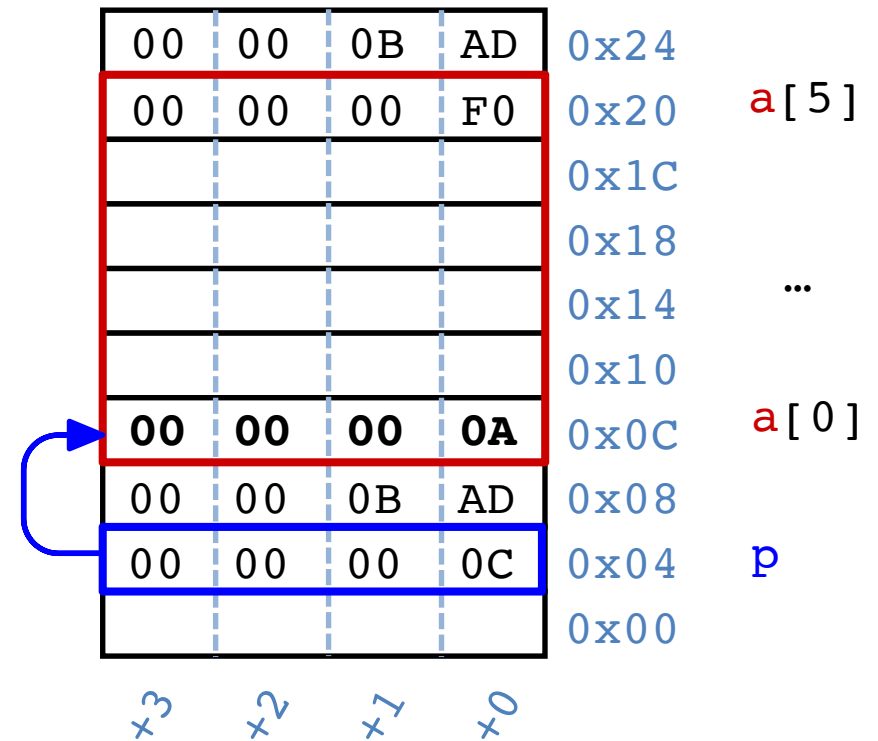
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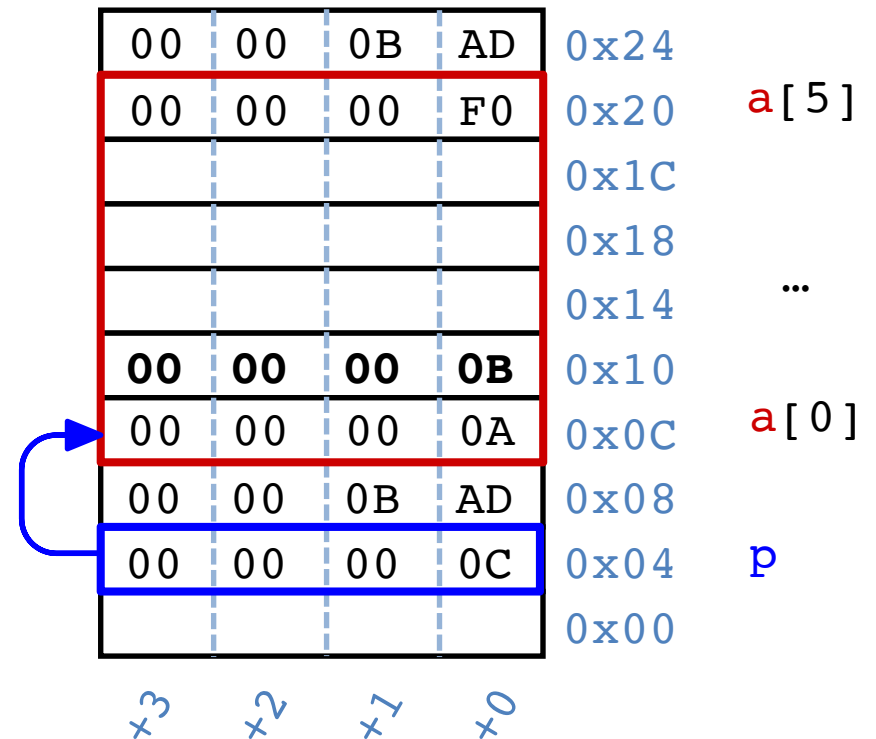
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Both are scaled by the size of the type.

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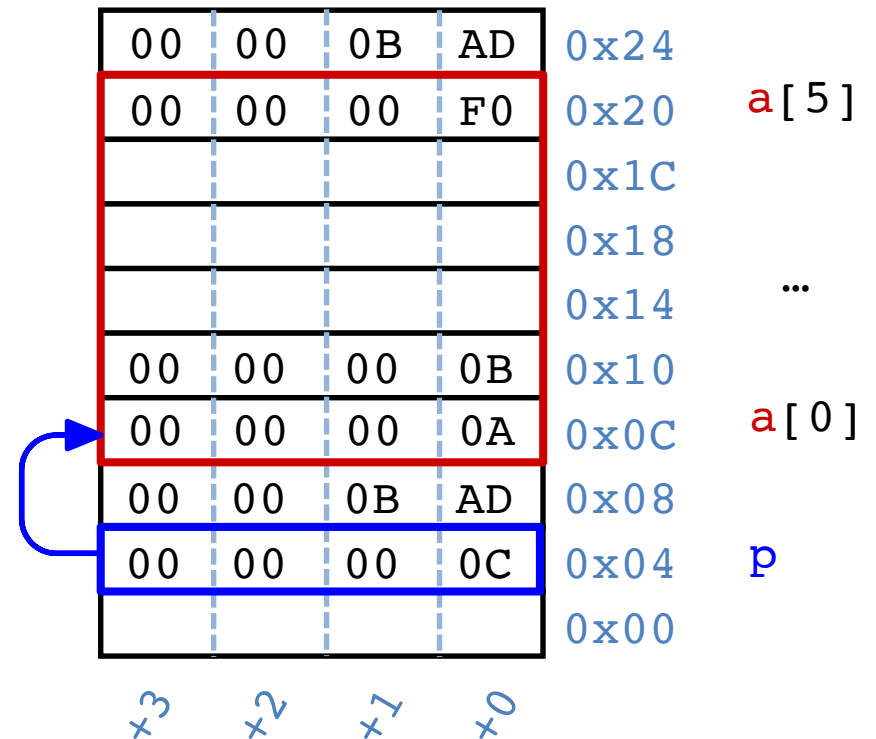
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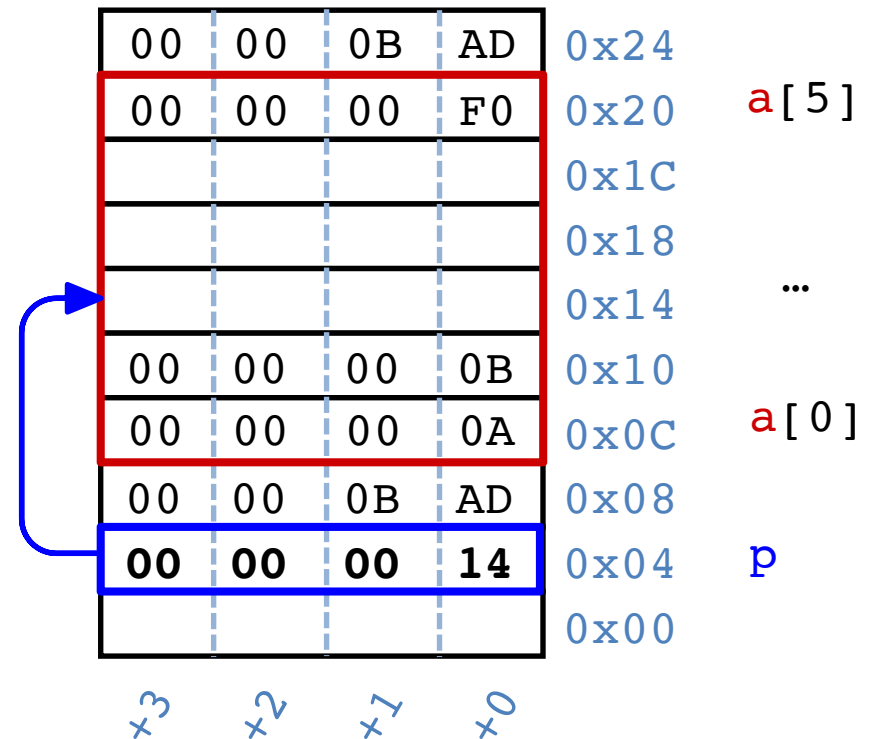
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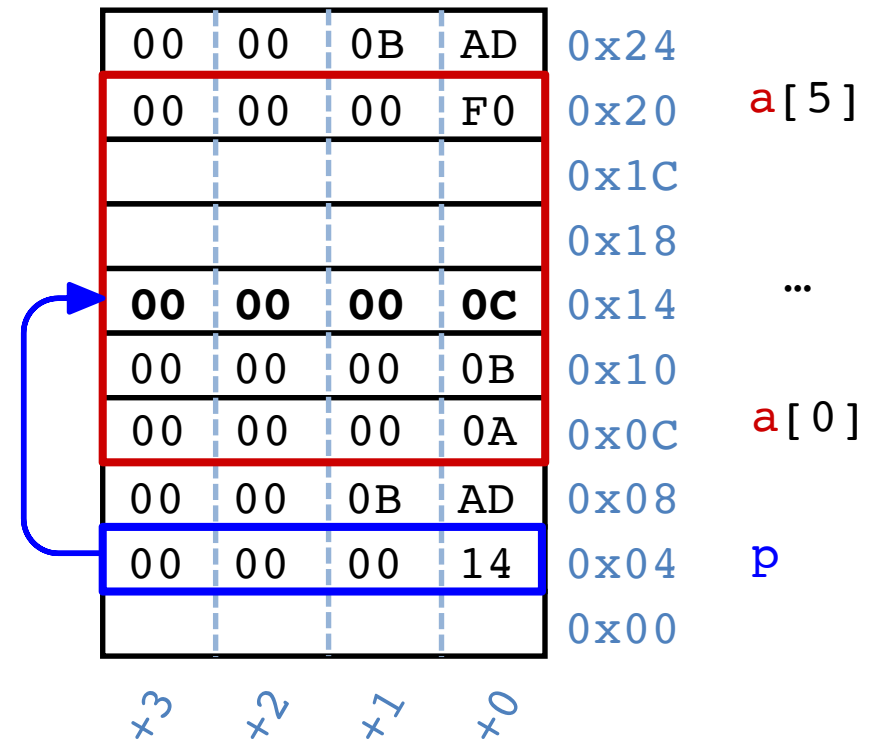
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`*p = a[1] + 1;`

C: Array allocation

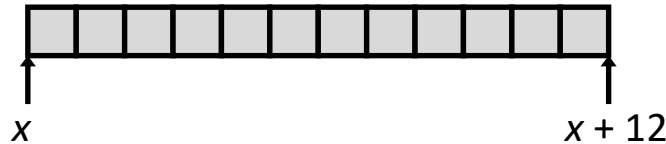
Basic Principle

T $A[N];$

Array of length N with elements of type T and name A

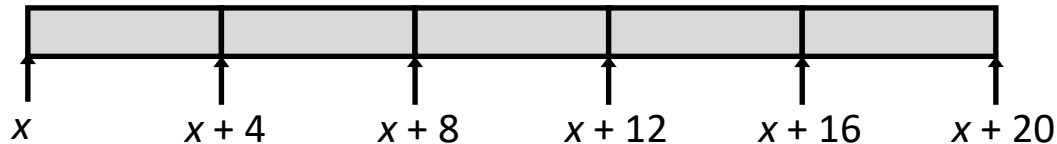
Contiguous block of $N * \text{sizeof}(T)$ bytes of memory

`char string[12];`

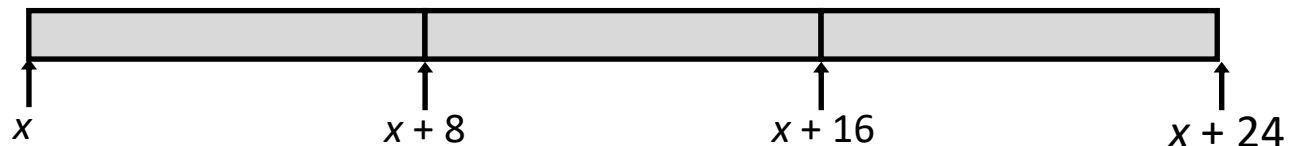


Use *sizeof* to determine proper size in C.

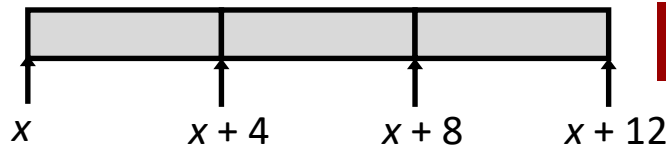
`int val[5];`



`double a[3];`



`char* p[3];`
(or `char *p[3];`)



IA32



x86-64

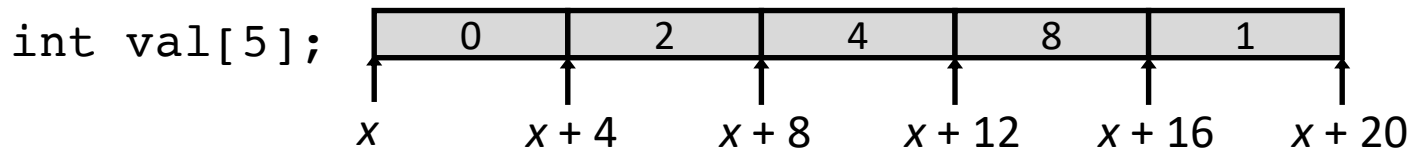
C: Array access

Basic Principle

$T \ A[N];$

Array of length N with elements of type T and name A

Identifier A has type T^*



Expression	Type	Value
<code>val[4]</code>	<code>int</code>	1
<code>val</code>	<code>int *</code>	
<code>val+1</code>	<code>int *</code>	
<code>&val[2]</code>	<code>int *</code>	
<code>val[5]</code>	<code>int</code>	
<code>*(val+1)</code>	<code>int</code>	
<code>val + i</code>	<code>int *</code>	

Representing strings

A C-style string is represented by an array of bytes (`char`).

- Elements are one-byte **ASCII codes** for each character.
- ASCII = American Standard Code for Information Interchange

32	space	48	0	64	@	80	P	96	`	112	p
33	!	49	1	65	A	81	Q	97	a	113	q
34	”	50	2	66	B	82	R	98	b	114	r
35	#	51	3	67	C	83	S	99	c	115	s
36	\$	52	4	68	D	84	T	100	d	116	t
37	%	53	5	69	E	85	U	101	e	117	u
38	&	54	6	70	F	86	V	102	f	118	v
39	'	55	7	71	G	87	W	103	g	119	w
40	(56	8	72	H	88	X	104	h	120	x
41)	57	9	73	I	89	Y	105	i	121	y
42	*	58	:	74	J	90	Z	106	j	122	z
43	+	59	;	75	K	91	[107	k	123	{
44	,	60	<	76	L	92	\	108	l	124	
45	-	61	=	77	M	93]	109	m	125	}
46	.	62	>	78	N	94	^	110	n	126	~
47	/	63	?	79	O	95	_	111	o	127	del

C: * and []



C programmers often use * where you might expect []:

e.g., char*:

- pointer to a char
- pointer to the first char in a string of unknown length

```
int strcmp(char* a, char* b);
```

C: 0 vs. '\0' vs. NULL

0

Name: zero
Type: int
Size: 4 bytes
Value: 0x00000000
Usage: The integer zero.

'\0'

Name: null character
Type: char
Size: 1 byte
Value: 0x00
Usage: Terminator for C strings.

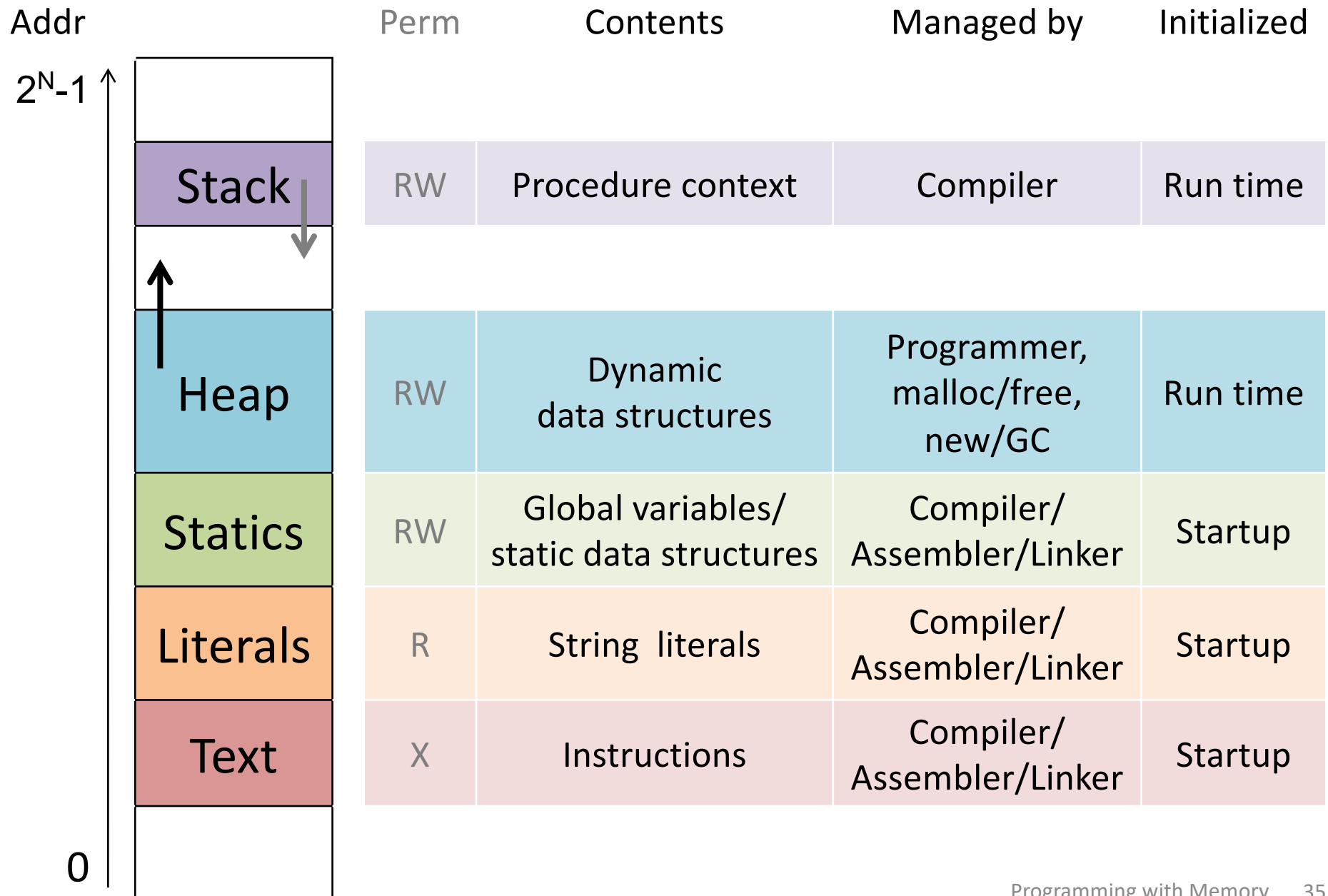
NULL

Name: null pointer / null reference / null address
Type: void*
Size: 1 word (= 8 bytes on a 64-bit architecture)
Value: 0x0000000000000000
Usage: The absence of a pointer where one is expected.
Address 0 is inaccessible, so *NULL is invalid; it crashes.

Is it important/necessary to encode the null character or the null pointer as 0x0?

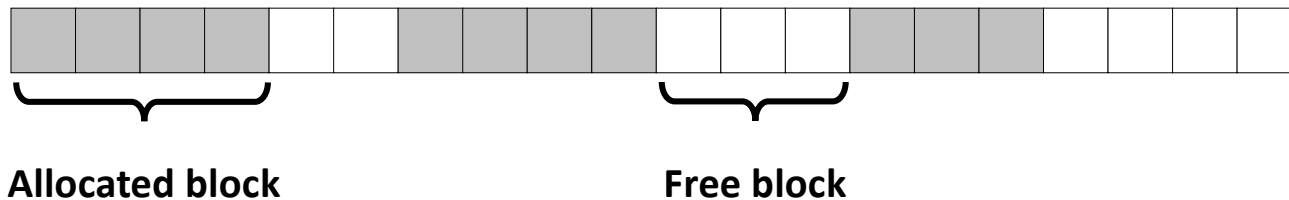
What happens if a programmer mixes up these "zeroey" values?

Memory address-space layout



C: Dynamic memory allocation in the heap

Heap:



Managed by memory allocator:

pointer to newly allocated block
of at least that size

number of contiguous bytes required

```
void* malloc(size_t size);
```

```
void free(void* ptr);
```

pointer to allocated block to free

C: standard memory allocator

```
#include <stdlib.h>    // include C standard library
```

```
void* malloc(size_t size)
```

Allocates a memory block of at least `size` bytes and returns its address.

If error (no space), returns `NULL`.

Rules:

Check for error result.

Cast result to relevant pointer type.

Use `sizeof(...)` to determine size.

```
void free(void* ptr)
```

Deallocates the block referenced by `ptr`, making its space available for new allocations.

`ptr` **must** be a `malloc` result that has not yet been freed.

Rules:

`ptr` must be a `malloc` result that has not yet been freed.

Do not use `*ptr` after freeing.

C: Dynamic array allocation

```
#define ZIP_LENGTH 5
int* zip = (int*)malloc(sizeof(int)*ZIP_LENGTH);
if (zip == NULL) { // if error occurred
    perror("malloc"); // print error message
    exit(0); // end the program
}
```

```
zip[0] = 0;
zip[1] = 2;
zip[2] = 4;
zip[3] = 8;
zip[4] = 1;
```

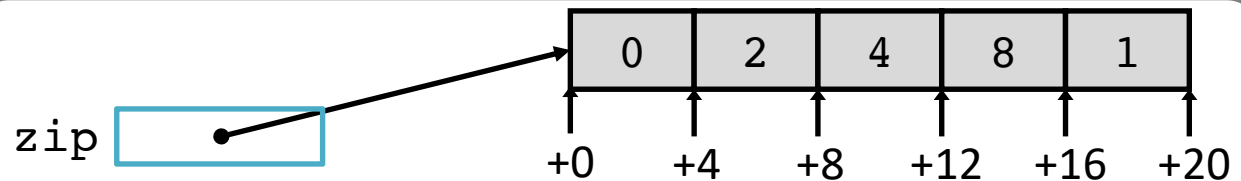
```
printf("zip is");
for (int i = 0; i < ZIP_LENGTH; i++) {
    printf(" %d", zip[i]);
}
```

```
printf("\n");
```

```
free(zip);
```

zip 0x7fedd2400dc0 0x7fff58bdd938

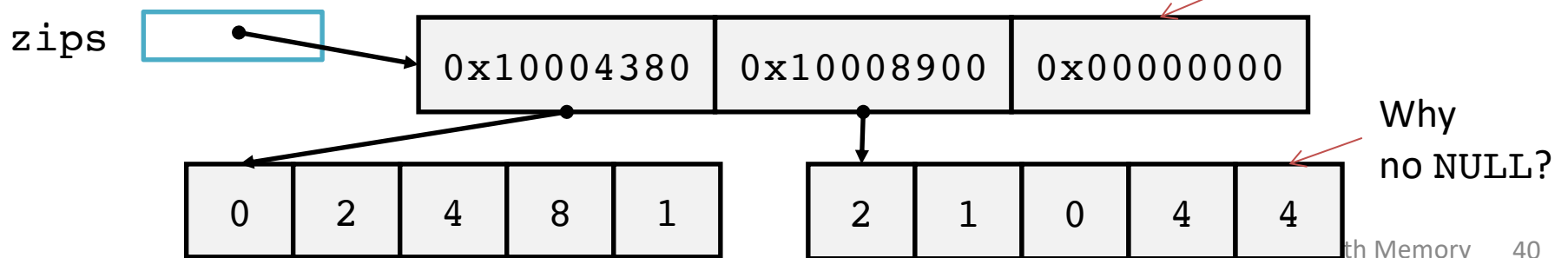
1	0x7fedd2400dd0
8	0x7fedd2400dcc
4	0x7fedd2400dc8
2	0x7fedd2400dc4
0	0x7fedd2400dc0



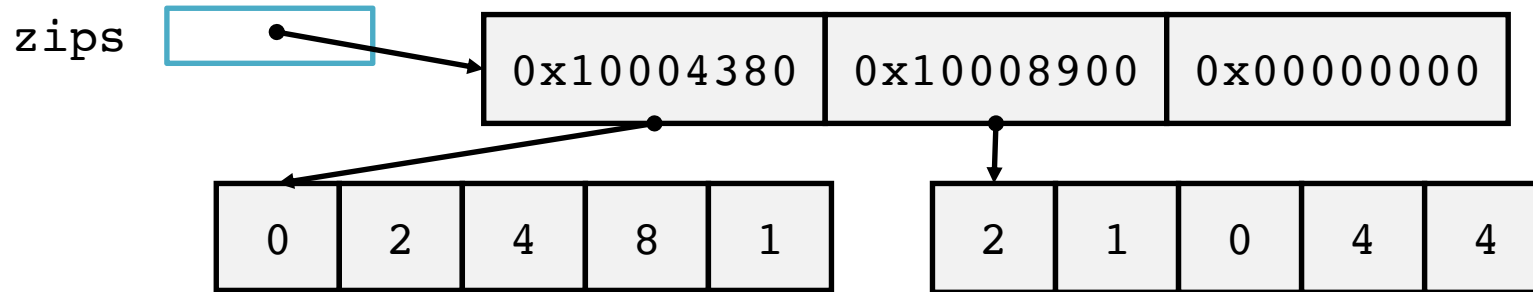
C: Array of pointers to arrays of ints

```
int** zips = (int**)malloc(sizeof(int*) * 3);  
  
zips[0] = (int*)malloc(sizeof(int)*5);  
int* zip0 = zips[0];  
zip0[0] = 0;  
zips[0][1] = 2;  
zips[0][2] = 4;  
zips[0][3] = 8;  
zips[0][4] = 1;  
  
zips[1] = (int*)malloc(sizeof(int)*5);  
zips[1][0] = 2;  
zips[1][1] = 1;  
zips[1][2] = 0;  
zips[1][3] = 4;  
zips[1][4] = 4;  
  
zips[2] = NULL;
```

Why terminate
with NULL?



Zip code



```
// return a count of all zips that end with digit endNum  
int zipCount(int* zips[], int endNum) {
```

```
}
```




<http://xkcd.com/138/>

C: scanf reads formatted input

```
int val;
```

Declared, but not initialized.
Holds anything.

```
...
```

```
scanf("%d", &val);
```

Read one `int`
in decimal₁₀ format
from input.

Store it in memory
at this address.

Store in memory at the address
given by the **address of `val`**:
store input @ `0x7F...F38`.

`val`

CA	FE	12	34	

`0x7FFFFFFFFFFFFFFF3C`

`0x7FFFFFFFFFFFFFFF38`

`0x7FFFFFFFFFFFFFFF34`

C: Classic bug using scanf



```
int val;  
...  
scanf("%d", val);
```

Declared, but not initialized.
Holds anything.

Read one int
in decimal₁₀ format
from input.

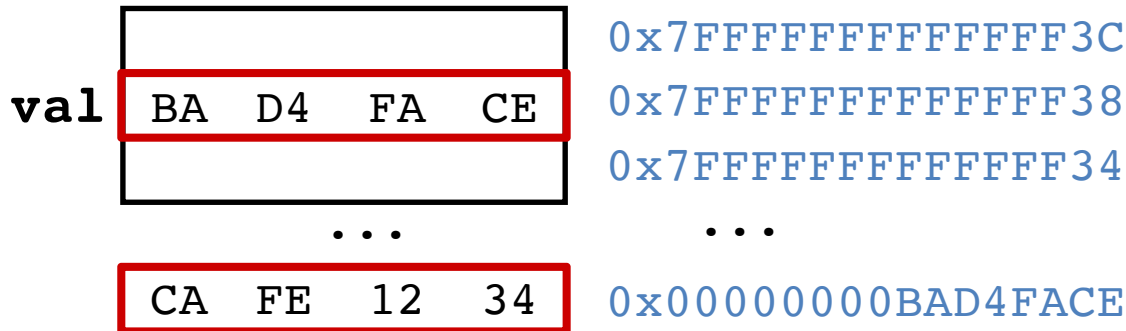
Store it in memory
at this address.

Store in memory at the address
given by the **contents of val**
(implicitly cast as a pointer):
store input @ 0xBAD4FACE.

Best case: 🤦 crash immediately
with segmentation fault/bus error.

Bad case: 😡 silently corrupt data
stored @ 0xBAD4FACE,
fail to store input in val,
and keep going.

Worst case: 💻🔥🚀 program does literally anything.



C: Memory error messages

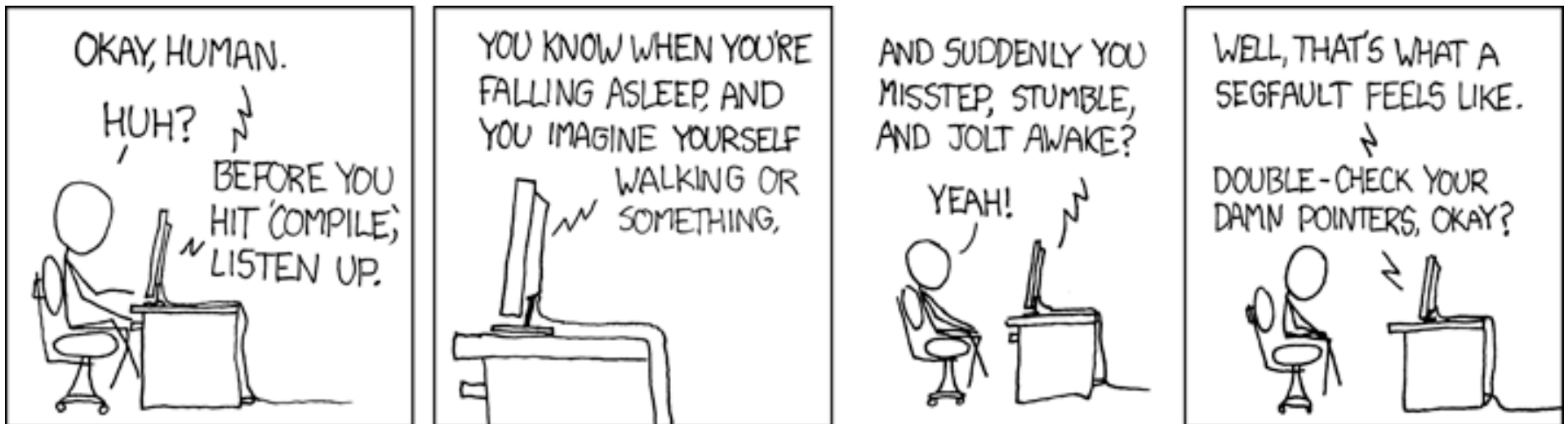
11: **segmentation fault** ("segfault", SIGSEGV)

accessing address outside legal area of memory

10: **bus error** (SIGBUS)

accessing misaligned or other problematic address

More to come on debugging!



<http://xkcd.com/371/>

C: Why?

Why learn C?

- Think like actual computer (abstraction close to machine level) without dealing with machine code.
- Understand just how much Your Favorite Language provides.
- Understand just how much Your Favorite Language might cost.
- Classic.
- Still (more) widely used (than it should be).
- Pitfalls still fuel devastating reliability and security failures today.

Why not use C?

- Probably not the right language for your next personal project.
- It "gets out of the programmer's way" even when the programmer is unwittingly running toward a cliff.
- Many advances in programming language design since then have produced languages that fix C's problems while keeping strengths.