# Laboratory 10 <br> Data Structures Representation Computer Science 240 

## One-dimensional arrays

Different languages use different implementations at the machine level to represent data structures.

In Java, arrays are actually implemented as arrays of addresses (pointers) to the elements, which are stored elsewhere in memory (not necessarily in contiguous locations).

In $C$, the elements of the array are stored in a contiguous block, starting at the base address of the array.

In the C model,
address of element in array $=$ base address + element size * index
If the size of the element is limited to 1 , 2 , or 4 bytes, what is another more efficient way to accomplish the multiplication?

In $C$, to define some arrays of 8 elements of different sizes:

```
int elements[] = {0x1, 0x3, 0x5, 0x7, 0x9, 0x11, 0x13, 0x15};
short welements[] = {0x23, 0x25, 0x27, 0x29, 0x31, 0x33, 0x35, 0x37}
byte belements[] = {0x20, 0x30, 0x40, 0x50, 0x60, 0x70, 0x80, 0x90}
```

The equivalent in X 86 is:
.data
elements: .long $0 \times 1,0 \times 3,0 \times 5,0 \times 7,0 \times 9,0 \times 11,0 \times 13,0 \times 15$
welements: .word $0 \times 23,0 \times 25,0 \times 27,0 \times 29,0 \times 31,0 \times 33,0 \times 35,0 \times 37$
belements: .byte $0 \times 20,0 \times 30,0 \times 40,0 \times 50,0 \times 60,0 \times 70,0 \times 80,0 \times 90$
Either would be displayed using gdb as:

| $0 \times 8049714<$ elements $>:$ | $0 \times 00000001$ | $0 \times 00000003$ | $0 \times 00000005$ | $0 \times 00000007$ |  |
| :--- | :--- | :--- | :--- | :--- | :--- |
| $0 \times 8049724<$ elements $+16>:$ | $0 \times 00000009$ | $0 \times 00000011$ | $0 \times 00000013$ | $0 \times 00000015$ |  |
| $0 \times 8049734$ | <welements $>:$ | $0 \times 00250023$ | $0 \times 00290027$ | $0 \times 00330031$ | $0 \times 00370035$ |

## Two-dimensional arrays

In C, when nested array of arrays are used, each row is stored contiguously in memory (row-major format), and the address of an element can be calculated by the following formula (size of row is the number of columns in a row):
address of element[row][col] =
base address of array +
(row * size of row * size of element) + (col * size of element)
-or-
base address of array + (row*size of row + col)*size of element

In $C$, to define a $4 x 4$ array of integers:
int twodarr[4][4] $=\{\{0 x 1,0 \times 2,0 x 3,0 x 4\}$,
$\{0 x 4,0 x 6,0 x 7,0 x 8\}$,
$\{0 \times 9,0 \times 10,0 \times 11,0 \times 12\}$,
$\{0 \times 13,0 \times 14,0 \times 15,0 \times 16\}\} ;$
The equivalent in X 86 is:

|  | .data |
| :--- | :--- |
| twodarr: $\quad$ | $. l o n g ~ 0 \times 1,0 \times 2,0 \times 3,0 \times 4$ |
|  | $. l o n g ~ 0 \times 5,0 \times 6,0 \times 7,0 \times 8$ |
|  | $. l o n g ~ 0 \times 9,0 \times 10,0 \times 11,0 \times 12$ |
|  | $. l o n g ~ 0 \times 13,0 \times 14,0 \times 15,0 \times 16$ |

Either would be displayed using gdb as:

| 0x80497a0 | <twodarr >: | 0x00000001 | 0x00000002 | 0x00000003 | 0x00000004 |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0x80497b0 | <twodarr+16>: | 0x00000004 | 0x00000006 | 0x00000007 | 0x00000008 |
| 0x80497c0 | <twodarr+32>: | 0x00000009 | 0x00000010 | 0x00000011 | 0x00000012 |
| 0x80497d0 | <twodarr+48>: | 0x00000013 | 0x00000014 | 0x00000015 | 0x00000016 |

