



Dynamic Dispatch

semantic essence
of "object-oriented" programming languages
(OOP)

How are names resolved?

Key piece of semantics in any language.

- ML, Racket:
 - Just one kind of *variables*.
 - Lexical scope – unambiguous binding
 - Record *field names* are not variables: no "lookup"
- Java,:
 - Local variables: lexical scope (more limited)
 - Instance variables, methods
 - Look up in terms of special `self` / `this` "variable"
 - it's more complicated...

Method lookup in OO languages

Two key questions for Java:

– General case:

What `m` is run by `____.m()` ?

– Specific case:

What `m` is run by `this.m()` ?

dynamic dispatch

```
class Point {
    double x, y;
    Point(double x, double y) {
        this.x = x; this.y = y;
    }
    double getX() { return this.x; }
    double getY() { return y; }
    double distFromOrigin() {
        return Math.sqrt(this.getX() * this.getX()
            + getY() * getY());
    }
}
```

implicit this.

```
class PolarPoint extends Point { // poor design, useful example
    double r, theta;
    PolarPoint(double r, double theta) {
        super(0.0, 0.0);    this.r = r;    this.theta = theta;
    }
    double getX() { return this.r * Math.cos(this.theta); }
    double getY() { return r * Math.sin(theta); }
}
```

```
Point p = ...;           // ???
p.getX();                 // ???
p.distFromOrigin();     // ???
```

overriding

Method lookup: example

```
Point p = ...;           // ???  
p.getX();                // ???  
p.distFromOrigin();     // ???
```

Key questions:

- Which **distToOrigin** is called?
- Which **getX**, **getY** methods does it call?

Dynamic dispatch

(a.k.a. late binding, virtual methods)

The unique OO semantics feature.

Method call: `e.m()`

Evaluation rule:

1. Under the current environment, evaluate `e` to value `v`.
2. Let `C` refer to the class of the receiver object `v`.
3. Until class `C` contains a method definition `m() { body }`
let `C` refer to the superclass of the current `C` and repeat step 3.
4. Under the environment of class `C`, extended with the binding `this ↦ v`, evaluate the `body` found in step 3.

Note: `this` refers to *current receiver object*, not containing class.

- `this.m()` uses dynamic dispatch just like other calls.
- NOT lexical scope, not dynamic scope

Dynamic Dispatch is not ...

$$\begin{aligned} & \text{obj0.m}(\text{obj1}, \dots, \text{objn}) \\ & \neq \\ & \text{m}(\text{obj0}, \text{obj1}, \dots, \text{objn}) \end{aligned}$$

Is **this** just an implicit parameter that captures a first argument written in a different spot?

NO!

"What **m** means" is determined by run-time class of **obj0**!

Must inspect **obj0** before starting to execute **m**.

this is different than any other parameters.

Key artifacts of dynamic dispatch

- Why **overriding** works...
`distFromOrigin` in `PolarPointA`
- Subclass's definition of **m** "shadows" superclass's definition of **m** when dispatching on object of subclass (or descendant) in all contexts, **even if dispatching from method in superclass.**
- More complicated than the rules for closures
 - Must treat `this` specially
 - May seem simpler only if you learned it first
 - Complicated != inferior or superior

Closed vs. open

ML: closures are, well, *closed*.

```
fun even x = if x=0 then true  else odd  (x-1)
and odd  x = if x=0 then false else even (x-1)
```

May shadow `even`, but calls to `odd` are unaffected.

```
(* does not change odd: too bad, would help *)
fun even x = (x mod 2) = 0
```

```
(* does not change odd: good, would break *)
fun even x = false
```

Closed vs. open

Most OOP languages: subclasses can change the behavior of superclass methods they do not override.

```
class A {
  boolean even(int x) {
    if (x == 0) return true;
    else return odd(x-1);
  }
  boolean odd(int x) {
    if (x == 0) return false;
    else return even(x-1);
  }
}
class B extends A { // improves odd in B objects
  boolean even(int x) { return x % 2 == 0; }
}
class C extends A { // breaks odd in C objects
  boolean even(int x) { return false; }
}
```

OOP trade-off: implicit extensibility

Any method that calls overridable methods may have its behavior changed by a subclass *even if it is not overridden*.

- On purpose, by mistake?
- Behavior depends on calls to overridable methods
- **Harder** to reason about “the code you're looking at.”
 - Sources of unknown behavior are pervasive:
all overridable methods transitively called by this method.
 - Avoid by disallowing overriding: “private” or “final”
- **Easier** for subclasses to extend existing behavior without copying code.
 - Assuming superclass method is not modified later

FP trade-off: explicit extensibility

A function that calls other functions may have its behavior affected *only where it calls functions passed as arguments*.

- *Easier* to reason about “the code you're looking at.”
 - Sources of unknown behavior are explicit: calls to argument functions.
- *Harder* for other code to extend existing behavior without copying code.
 - Only by functions as arguments.

Aside: *overloading* is static.

overloading:

> 1 methods in class have same name

overriding:

if and only if same number/types of arguments

Pick the "best" overloaded method using the *static* types of the arguments

- Complicated rules for “best”
- Some confusion when expecting wrong *over*-thing

static dispatch

(a.k.a early binding, non-virtual methods)

Requires
static types...

- Lookup method based on static type of receiver.
- Calls to `e.m2()` where `e` has declared class `C`
 - *(the lexically enclosing class is this's "declared class")*
 - **always resolve** to "closest" method `m2` defined in `C` or `C`'s ancestor classes
 - completely ignores run-time class of object result of `e`
- ... similar to lexical scope for method lookup with inheritance.
- Same method call **always** resolves to same method definition.
- Determined statically by type system *before* running program.
- used for `super` in Java, non-virtual methods in C++

static dispatch

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    }
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    double getY() { return y; }
    double distFromOrigin() {
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overriding