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# Programming Languages

**CS 251**  
*Fall 2021*

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Recap

# Language components

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- ◆ **Values**: expressions that cannot be reduced any further
- ◆ **Expressions**: bits of the language
- ◆ **Declarations**: bind variables to values

# Our big step semantics for Racket

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- ◆ **Values**: expressions that cannot be reduced any further

Value rule:  $v \downarrow v$

- ◆ **Expressions**: bits of the language

Addition rule:

$$\frac{\begin{array}{l} e1 \downarrow v1 \\ e2 \downarrow v2 \end{array}}{(+ e1 e2) \downarrow v}$$

where  $v1$  and  $v2$  are numbers and  $v$  is the sum of  $v1$  and  $v2$

# Big step semantics: practice

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Let's write the big step semantics for **and**

**And:**

????

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(and e1 e2) ↓ v

# Big step semantics: practice

Let's write the big step semantics for **and**

**And:**

$$\frac{\begin{array}{l} e1 \downarrow v1 \\ e2 \downarrow v2 \end{array}}{\text{(and } e1 \ e2) \downarrow v}$$

If  $v1$  and  $v2$  are Booleans, either  $v = v1 = v2 = \#t$  or  $v$  is  $\#f$ . Otherwise, a dynamic error is produced.

# Big step semantics: practice

Let's write the big step semantics for **and**:

**And-True:**

$$e1 \downarrow \#t$$
$$e2 \downarrow \#t$$

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$$(\text{and } e1 \ e2) \downarrow \#t$$

Otherwise, a dynamic error is produced.

**And-False:**

$$e1 \downarrow v1$$
$$e2 \downarrow v2$$

---

$$(\text{and } e1 \ e2) \downarrow \#f$$

If  $v1$  or  $v2$  is  $\#f$  and both are Booleans,  $(\text{and } e1 \ e2)$  evaluates to  $\#f$ . Otherwise, a dynamic error is produced.

# First class functions

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In Racket, functions are values. This is because Racket has **first class functions**: functions have all the rights and privileges of other values.

## Function Bill of Rights:

*We the Racketeers hereby declare that functions:*

- ◆ Do not need to be named (lambdas)
- ◆ Can be returned by functions
- ◆ Can be arguments to functions



# Functions as values

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Do we need to do anything special for functions?

No! Like other values, functions can be evaluated any further... *until they are applied.*

Syntax: (**lambda** ( $id_1, \dots, id_n$ )  $e$ )



**parameters**

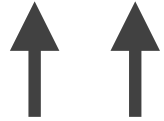
**function body**

Semantics:  $v \downarrow v$

# Function application

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Syntax:  $(e_1 e_2)$



**function**   **argument**

Semantics: ????

What happens when a function is **applied**?

# Function application

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When a function is applied to a value, that value gets **bound** to the function's **parameter** inside the **scope** of the function.

```
(define (id e) e)  
(id 5)
```

# Function application

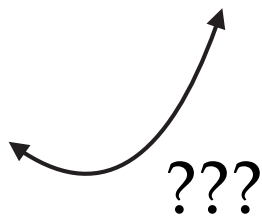
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When a function is applied to a value, that value gets **bound** to the function's **parameter** inside the **scope** of the function.

**But what does that mean?**

`(define (id e) e)`

`(id 5)`



# Variables and binding

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What are variables?

What is variable binding?

# Variables and binding

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What are variables?

- ◆ Variables **store values**

# Variable binding

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What is variable binding?

- ◆ Variable binding links the value to all occurrences of the variable within the **binder's scope**

# Variable binding

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What is variable binding?

- ◆ Variable binding links the value to all occurrences of the variable within the **binder's scope**

Ok, but what does it mean to “link” the value and variable?



# Variable binding

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What is variable binding?

- ◆ Variable binding links the value to all occurrences of the variable within the **binder's scope**

Ok, but what does it mean to “link” the value and variable?

- ◆ There are different ways of thinking about this! For now, we will use:

**The Substitution Model of Variable Binding:** when a value  $v$  is bound to an expression  $e$ , substitute the value  $v$  for every occurrence of  $e$  in the scope of the binder.

# Binding as substitution

**The Substitution Model of Variable Binding:** when a value  $v$  is bound to an expression  $e$ , substitute the value  $v$  for every occurrence of  $e$  in the scope of the binder.

$$(\text{id } 5) = ((\text{lambda } (e) e) 5) = 5$$

# Scope

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What is the scope of a variable binding?  
That depends on the binding construct!

# Discovering variable scope

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What is the scope of a variable binding?  
That depends on the binding construct!

We can figure out the scope of a variable through experimentation.

Starting premise:

- ◆ Referencing an **unbound** variable will throw an error
- ◆ Referencing a **bound variable** will return the value it is bound to

# Discovering Scope: **let**

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**Goal:** figure out the binding scope of **let**

# Tracking our discoveries

**Goal:** figure out the binding scope of **let**

We should keep track of what we are learning about Racket's behavior!

Context	Variable value
before <b>let</b>	
within <b>let</b>	
after <b>let</b>	

# Tracking our discoveries

**Goal:** figure out the binding scope of **let**

We should keep track of what we are learning about Racket's behavior!

Context	Variable value
before <b>let</b>	undefined
within <b>let</b>	
after <b>let</b>	

# Tracking our discoveries

**Goal:** figure out the binding scope of **let**

We should keep track of what we are learning about Racket's behavior!

Context	Variable value
before <b>let</b>	undefined
within <b>let</b>	5
after <b>let</b>	



# Discovering Scope: **let**

**Goal:** figure out the binding scope of **let**

**Hypothesis 1:** **let** binds all occurrences of the variable in its body.

Context	Variable value
before <b>let</b>	undefined
within <b>let</b>	5
after <b>let</b>	undefined

# Discovering Scope: **let**

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**Goal:** figure out the binding scope of **let**

**Hypothesis 1:** **let** binds all occurrences of the variable in its body.

What about nested **let**?

# Discovering Scope: **let**

**Goal:** figure out the binding scope of **let**

What about nested **let**?

Context	Variable value
before first <b>let</b>	
within first <b>let</b>	
within second <b>let</b>	
after both <b>lets</b>	

# Discovering Scope: **let**

**Goal:** figure out the binding scope of **let**

What about nested **let**?

Context	Variable value
before first <b>let</b>	undefined
within first <b>let</b>	
within second <b>let</b>	
after both <b>lets</b>	

# Discovering Scope: **let**

**Goal:** figure out the binding scope of **let**

What about nested **let**?

Context	Variable value
before first <b>let</b>	undefined
within first <b>let</b>	"Outer"
within second <b>let</b>	
after both <b>lets</b>	

# Discovering Scope: **let**

**Goal:** figure out the binding scope of **let**

What about nested **let**?

Context	Variable value
before first <b>let</b>	undefined
within first <b>let</b>	"Outer"
within second <b>let</b>	"Inner"
after both <b>lets</b>	

# Discovering Scope: **let**

**Goal:** figure out the binding scope of **let**

What about nested **let**?

Context	Variable value
before first <b>let</b>	undefined
within first <b>let</b>	"Outer"
within second <b>let</b>	"Inner"
after both <b>lets</b>	undefined

# Variable shadowing

We've run into a case of **variable shadowing**: there is another binder for **x** nested within the first binder.

If we want to be able to refer to the first value, we need to use two different variable names.

```
(let (x 7)
  (let (x 2)
    (+ 0 x)
  )
  x
)
```



# Discovering Scope: **let**

**Goal:** figure out the binding scope of **let**

**Hypothesis 2:** **let** binds all **unbound** occurrences of the variable in its body.

Context	Variable value
before first <b>let</b>	undefined
within first <b>let</b>	"Outer"
within second <b>let</b>	"Inner"
after both <b>lets</b>	undefined

# Binding as substitution

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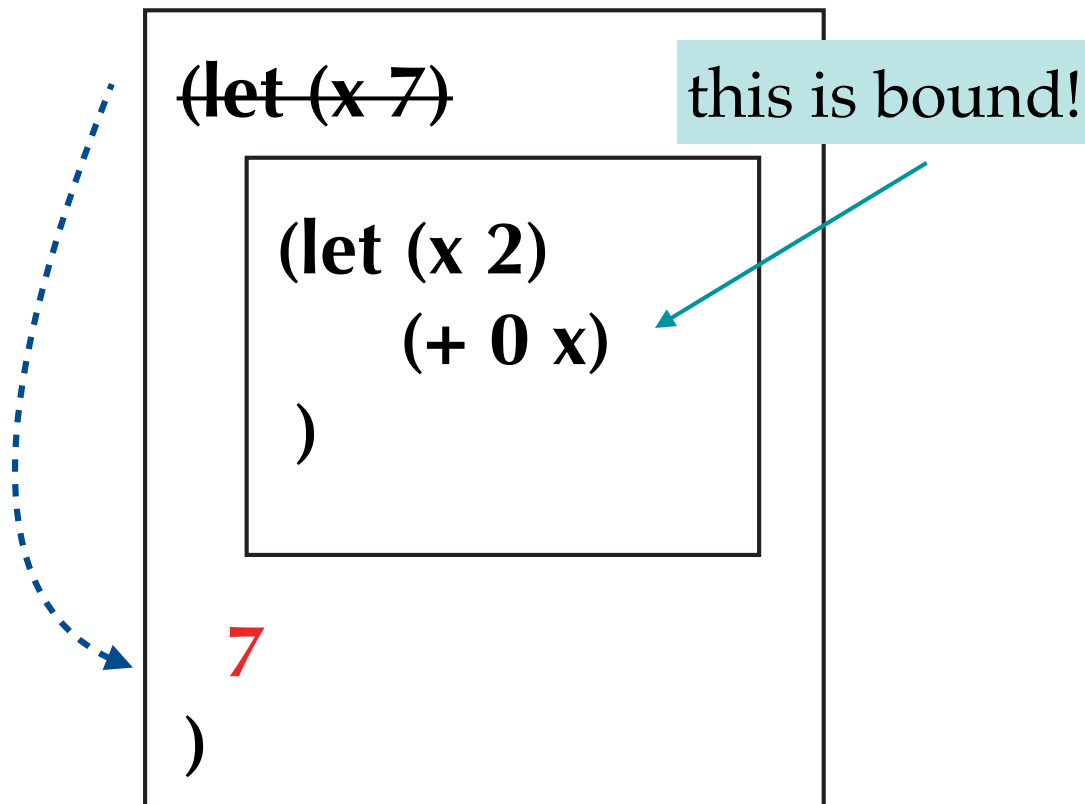
## **The Substitution Model of Variable Binding:**

When a value  $v$  is bound to an expression  $e$ , substitute the value  $v$  for every **unbound** occurrence of  $e$  in the scope of the binder.

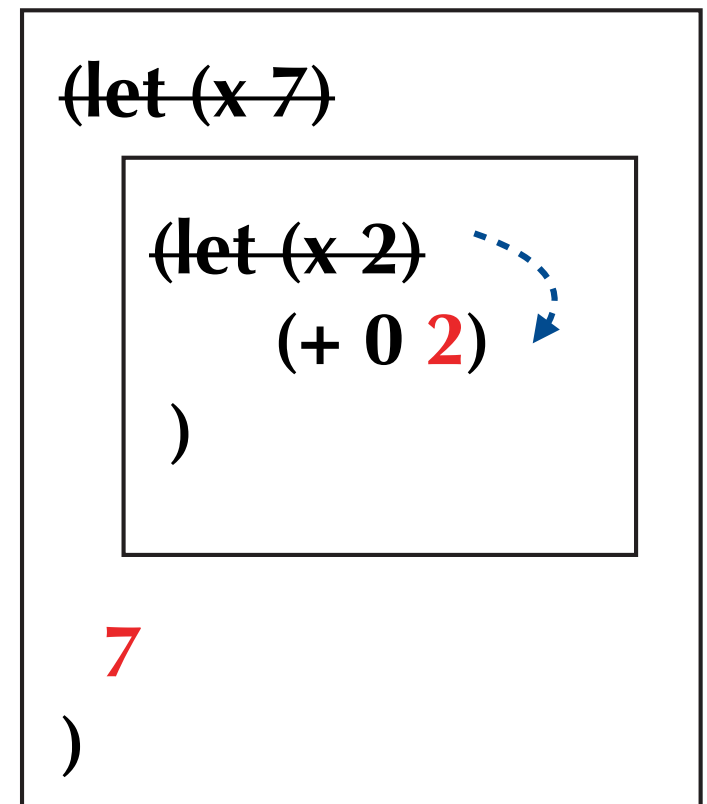
# Understanding variable shadowing

Variable shadowing may look like it mutates the variable. But doesn't. If we apply the substitution model of variable binding, we see:

## Step 1



## Step 2



# Variable binding and scope

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What is variable binding?

- ◆ Variable binding links the value to all **unbound** occurrences of the variable within the **binder's scope**

What is the scope of a variable binding?

- ◆ That depends on the binding construct!

The binding scope of a **let** is its body.

# Practice:

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Use the same technique to figure out the binding scope of **function application**.

# Discovering Scope: Application

**Goal:** figure out the binding scope of **function application**

Context	Variable value
before function application	undefined
within function body	"Outer"
outside of function call	undefined

# Variable binding and scope

---

What is variable binding?

- ◆ Variable binding links the value to all **unbound** occurrences of the variable within the **binder's scope**

What is the scope of a variable binding?

- ◆ That depends on the binding construct!

The binding scope of a **let** is its body.

The binding scope of a **function** is its body.

What is the binding scope of **define**?

# Practice:

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Use the same technique to figure out the binding scope of **define**.



# Discovering Scope: Application

**Goal:** figure out the binding scope of **define**

Context	Variable value
before define	
within function body	
within function body when name is shadowed	
after define	

# Discovering Scope: Application

**Goal:** figure out the binding scope of **define**

Context	Variable value
before define	undefined
within function body	
within function body when name is shadowed	
after define	

# Discovering Scope: Application

**Goal:** figure out the binding scope of **define**

Context	Variable value
before define	undefined
within function body	"donut"
within function body when name is shadowed	
after define	

# Discovering Scope: Application

**Goal:** figure out the binding scope of **define**

Context	Variable value
before define	undefined
within function body	"donut"
within function body when name is shadowed	"mocha"
after define	

# Discovering Scope: Application

**Goal:** figure out the binding scope of **define**

Context	Variable value
before define	undefined
within function body	"donut"
within function body when name is shadowed	"mocha"
after define	"donut"

# Define

What is the binding scope of **define**?

Similar to **let**, but there's no body. (**define** e x) scopes over all subsequent unbound occurrences of e within the current scope ( the scope it is called in).

```
(define x 7)
```

x : 7

```
(let ()
```

```
  (define x 2)
```

```
  (+ 0 x)
```

```
)
```

x : 2

x : 7

# Define

---

What is the binding scope of **define**?

Similar to **let**, but there's no body. (**define** e x) scopes over all subsequent unbound occurrences of e within the current scope.

**define** can be used within functions and local binding constructs like **let**. But this is considered poor style.

# What have we discovered?

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- ◆ Checked our big step semantics for functions
- ◆ Tried to think about the semantics of function application
- ◆ Explored **scope** and **variable binding**
- ◆ Learned the **substitution model** of variable binding

**Next class:** big step semantics for function application and variable binding constructs!