Metaprogramming

These slides borrow heavily from Ben Wood's Fall '15 slides.

CS251 Programming Languages
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How to implement a programming language

Interpretation

An **interpreter** written in the **implementation language** reads a program written in the **source language** and **evaluates** it.

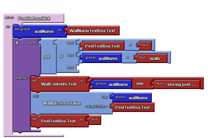
Translation (a.k.a. compilation)

An **translator** (a.k.a. **compiler)** written in the **implementation language** reads a program written in the **source language** and **translates** it to an equivalent program in the **target language**.

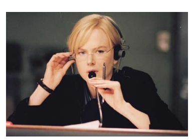
But now we need implementations of:

- **implementation language**
- **target language**

Metaprogramming: Interpretation



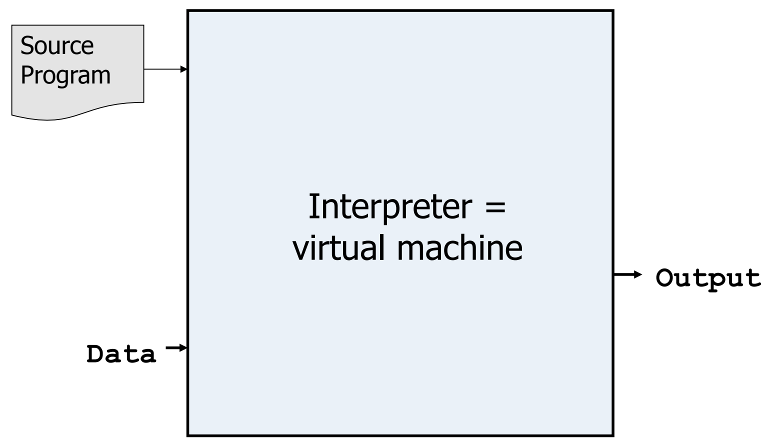
Program in language L



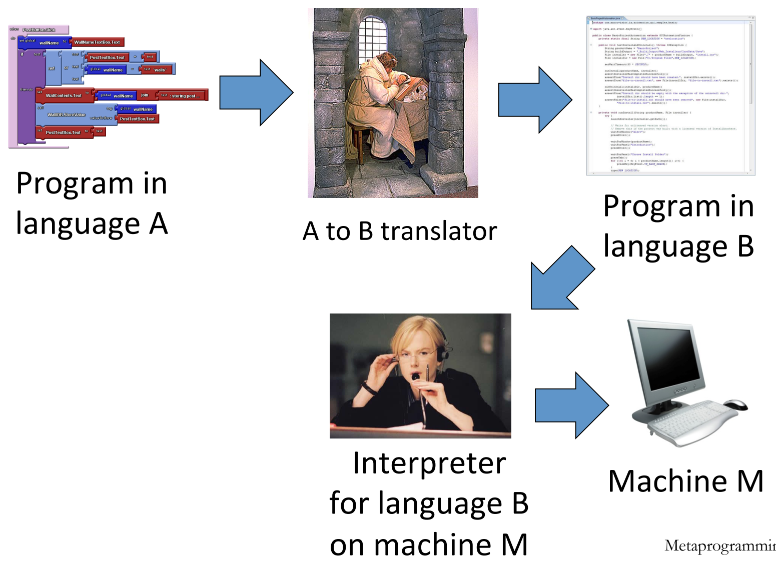
Interpreter for language L on machine M

Machine M

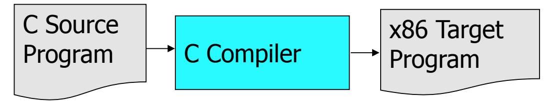
Interpreters



Metaprogramming: Translation

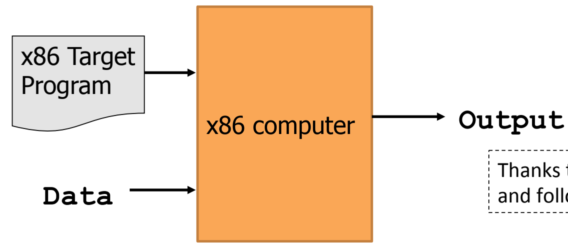


Compiler




```
if (x == 0) {
  x = x + 1;
}
...
```

```
cmp (1000), $0
bne L
add (1000), $1
L:
...
```



Thanks to **Ben Wood** for these and following pictures

Java Compiler


```
if (x == 0) {
  x = x + 1;
}
...
```

```
load 0
ifne L
load 0
inc
store 0
L:
...
```

`(compare compiled C to compiled Java)`

Interpreters vs Compilers

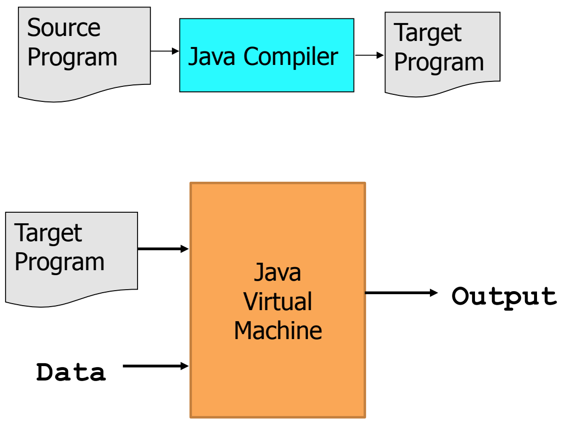
Interpreters

- No work ahead of time
- Incremental
- maybe inefficient

Compilers

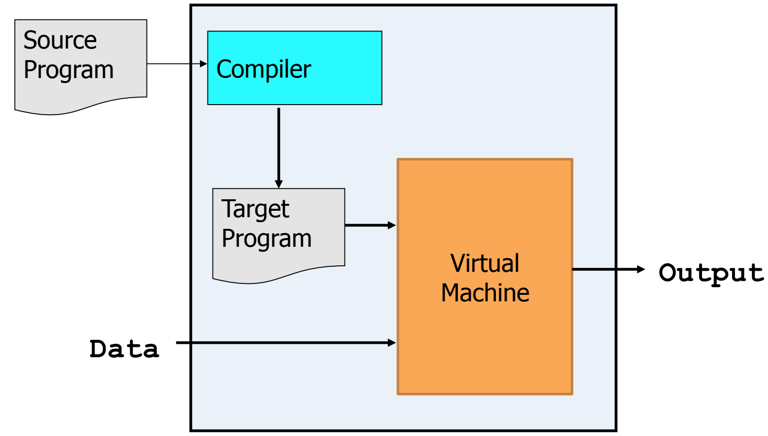
- All work ahead of time
- See whole program (or more of program)
- Time and resources for analysis and optimization

Compilers... whose output is interpreted

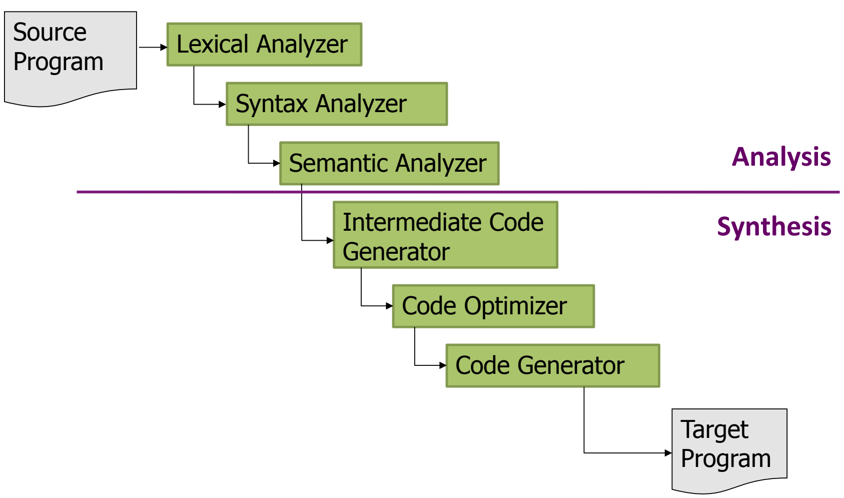


Doesn't this look familiar?

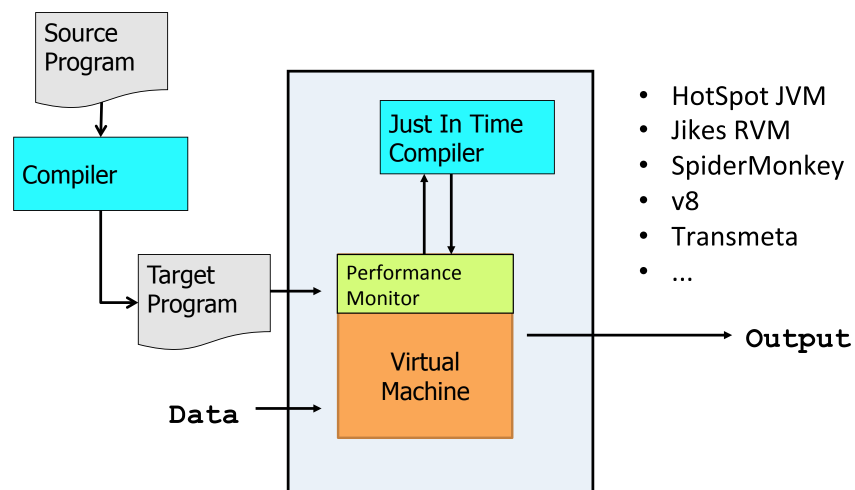
Interpreters... that use compilers.



Typical Compiler

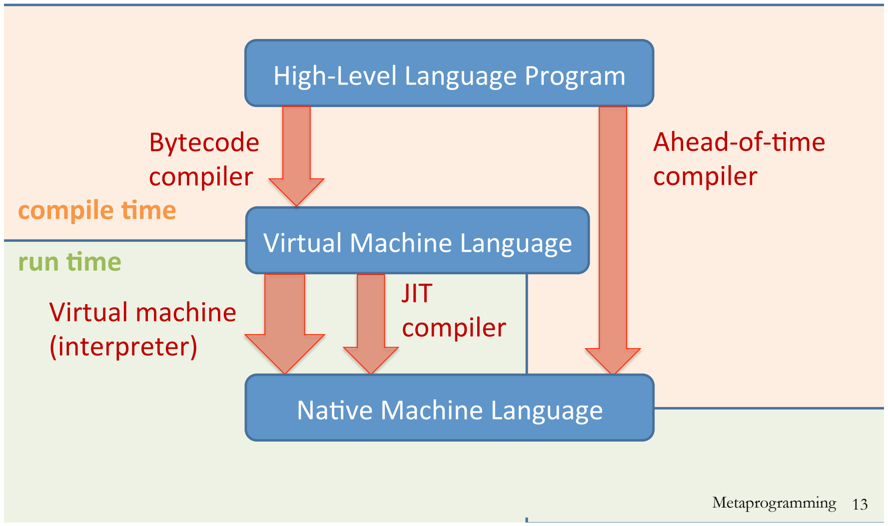


JIT Compilers and Optimization



- HotSpot JVM
- Jikes RVM
- SpiderMonkey
- v8
- Transmeta
- ...

Virtual Machine Model



How to implement a programming language

Can describe by deriving a "proof" of the implementation using these inference rules:

Interpreter Rule

P-in-L program   L interpreter machine

P machine

Translator Rule

P-in-S program   S-to-T translator machine

P-in-T program

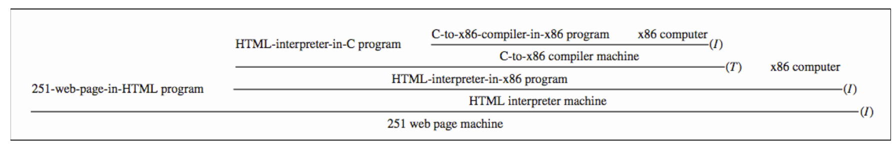
Implementation Derivation Example

Prove how to implement a "251 web page machine" using:

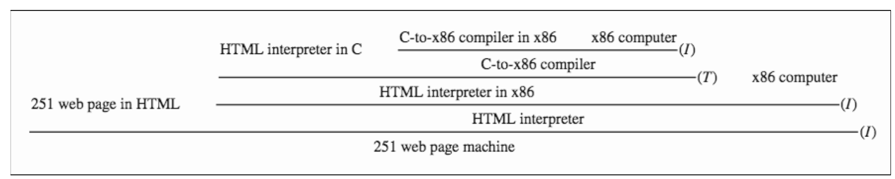
- 251-web-page-in-HTML program (a web page written in HTML)
- HTML-interpreter-in-C program (a web browser written in C)
- C-to-x86-translator-in-x86 program (a C compiler written in x86)
- x86 interpreter machine (an x86 computer)

No peaking ahead!

Implementation Derivation Example Solution



We can omit some occurrences of "program" and "machine":



Implementation Derivation Are Trees

And so we can represent them as nested structures, like nested bulleted lists:

- ❑ 251-web-page-in-HTML program
    - o HTML-interpreter-in-C program
        - • C-to-x86 compiler-in-x86 program
        - • X86 computer
    - o C-to-x86 compiler machine (I)
  - ✧ HTML-interpreter-in-x86 program (T)
  - ✧ x86 computer
- ❑ HTML interpreter machine (I)

251 web page machine (I)

Version that shows conclusions below bullets. More similar to derivations with horizontal lines, but harder to create,and read

251 web page machine (I)

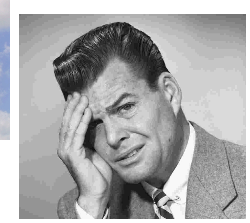
- ❑ 251-web-page-in-HTML program
- ❑ HTML interpreter machine (I)
  - ✧ HTML-interpreter-in-x86 program (T)
    - o HTML-interpreter-in-C program
    - o C-to-x86 compiler machine (I)
        - • C-to-x86 compiler-in-x86 program
        - • X86 computer
  - ✧ x86 computer

Preferred “top-down” version that shows conclusions above bullets.

Metaprogramming: Bootstrapping Puzzles

How can we write Scheme interpreter in Scheme?

How can we write a Java-to-x86 compiler in Java?



Metacircularity and Bootstrapping

Many examples:

- Lisp in Lisp / Scheme in Scheme/Racket in Racket
- Python in Python: PyPy
- Java in Java: Jikes RVM, Maxine VM
- ...
- C-to-x86 compiler in C
- `eval` construct in languages like Lisp, JavaScript

How can this be possible?

Key insights to bootstrapping:

- The first implementation of a language **cannot** be in itself, but must be in some other language.
- Once you have one implementation of a language, you can implement it in itself.

Metacircularity Example 1: Problem

Suppose you are given:

- Scheme-interpreter-in-Python program
- Python machine
- Scheme-interpreter-in-Scheme program

How do you create a Scheme interpreter machine using the Scheme-interpreter-in-Scheme program?

Metacircularity Example 1: Solution

Suppose you are given:

- Scheme-interpreter-in-Python program
- Python machine
- Scheme-interpreter-in-Scheme program

How do you create a Scheme interpreter machine using the Scheme-interpreter-in-Scheme program?

Scheme interpreter machine #2 (I)
- ❑ Scheme-interpreter-in-Scheme program
- ❑ Scheme-interpreter machine #1 (I)
  - ✧ Scheme-interpreter-in-Python program
  - ✧ Python machine

But why create Scheme interpreter machine #2 when you already have Scheme-interpreter machine #1?

Metacircularity Example 1: More Realistic

Suppose you are given:

- Scheme-subset-interpreter-in-Python program (implements only core Scheme features; no desugaring or other frills)
- Python machine
- Full-Scheme-interpreter-in-Scheme program

How do you create a Full-Scheme interpreter machine using the Full-Scheme-interpreter-in-Scheme program?

Full-Scheme interpreter machine (I)
- ❑ Scheme-interpreter-in-Scheme program
- ❑ Scheme-subset interpreter machine #1 (I)
  - ✧ Scheme-subset-interpreter-in-Python program
  - ✧ Python machine

Metacircularity Example 2: Problem

Suppose you are given:

- C-to-x86-translator-in-x86 program (a C compiler written in x86)
- x86 interpreter machine (an x86 computer)
- C-to-x86-translator-in-C-subset program

How do you compile the C-to-x86-translator-in-C ?

Metacircularity Example 2: Solution

Suppose you are given:

- C-to-x86-translator-in-x86 program (a C compiler written in x86)
- x86 interpreter machine (an x86 computer)
- C-to-x86-translator-in-C program

How do you compile the C-to-x86-translator-in-C ?

C-to-x86-translator machine #2 (I)
- ❑ C-to-x86-translator-in-x86 program #2 (T)
  - ✧ C-to-x86-translator-in-C
  - ✧ C-to-x86-translator machine #1 (I)
    - ○ C-to-x86-translator-in-x86 program #1
    - ○ x86 computer
- ❑ x86 computer

But why create C-to-x86-translator-in-x86 program #2 (T) when you already have C-to-x86-translator-in-x86 program #1?

Metacircularity Example 2: More Realistic

Suppose you are given:

- C-subset-to-x86-translator-in-x86 program
  (a compiler for a subset of C written in x86)
- x86 interpreter machine (an x86 computer)
- Full-C-to-x86-translator-in-C-subset program
  (a compiler for the full C language written in a subset of C)

How do you create a Full-C-to-x86-translator machine ?

Full-C-to-x86-translator machine (I)
- ❑ Full-C-to-x86-translator-in-x86 program (T)
  - ✧ Full-C-to-x86-translator-in-C-subset
  - ✧ C-subset-to-x86-translator machine (I)
    - ○ C-subset-to-x86-translator-in-x86 program
    - ○ x86 computer
- ❑ x86 computer

A long line of C compilers

C-version_n-to-target_n-translator machine (I)
- ❑ C-version_n-to-target_n-translator program in target_n-1 (T)
  - ✧ C-version_n-to-target_n-translator program in C-version_n-1
  - ✧ C-version_n-1-to-target_n-1 translator machine (I)
    - ○ C-version_n-1-to-target_n_1-translator program in target_n-2 (T)
      - ⋮
      - ➢ C-version_2-to-target_2-translator-program in target_1 (T)
        - ▪ C-version_2-to-target_2-translator program in C-version_1
        - ▪ C-version_1-to-target_1 translator machine (I)
          - • C-version_1-to-target_1-translator program in assembly_0
          - • assembly_0 computer
      - ➢ target_1 computer
      - ⋮
    - ○ target_n-2 computer
- ❑ target_n-1 computer

- ○ The versions of C and target languages can change at each stage.
- ○ Trojan horses from earlier source files can remain in translator machines even if they're not in later source file! See Ken Thompson's *Reflection on Trusting Trust*

More Metaprogramming in SML

- We've already seen PostFix and Intex SML
- A sequences of expression languages implemented in SML that look closer and closer to Racket:
  - Bindex: add naming
  - Valex: add more value types, dynamic type checking, desugaring
  - HOFL: first class function values, closure diagrams

Remember: language != implementation

- Easy to confuse "the way this language is usually implemented" or "the implementation I use" with "the language itself."
- Java and Racket can be compiled to x86
- C can be interpreted in Racket
- x86 can be compiled to JavaScript
- Can we compile C/C++ to Javascript?
  http://kripken.github.io/emscripten-site/