

## TAC: Three-Address Code

---

This document summarizes a simple three-address code (TAC) targeted as an intermediate code representation in a compiler. You will likely wish to change or extend this instruction set to develop a TAC language for use as an intermediate format in your IC compiler. There are many potential design choices, and you should not treat this specification as final. See Dragon 6 (especially 6.2) or EC 5 for further discussions of three-address code and other intermediate representations.

TAC is a “flat” language without nested expressions. Every instruction references or *addresses* three or fewer distinct variables (*a*, *b*, *c*, *etc.*, more formally called addresses), constants (*301*, *false*, *"hello"* *etc.*), or *labels* (markers in the code: *L*, *etc.*).

## Instruction Forms

---

There are four basic types of instructions.

- **Arithmetic and Logic Instructions.**

Basic instruction forms include:

- unary operators  $a = OP\ b$ , where *OP* may be a unary operator: *-*, *!*
- binary operators  $a = b\ OP\ c$ , where *OP* can be
  - an arithmetic operator: *+*, *-*, */*, *\**
  - a logic operator: *&&*, *||*
  - a comparison operator: *==*, *!=*, *<*, *<=*, *>*, *>=*

- **Data Movement Instructions.**

Copy:  $a = b$   
 Load/store:  $a = *b$      $*a = b$   
 Array load/store:  $a = b[i]$      $a[i] = b$   
 Field load/store:  $a = b.f$      $a.f = b$

- **Branch Instructions.**

Label:                    *label L*  
 Unconditional jump:    *jump L*  
 Conditional jump:      *cjump a L*    (jump to *L* if *a* is true)

- **Function Call Instructions.**

Call with no result:    *call f(a<sub>1</sub>, ..., a<sub>n</sub>)*  
 Call with result:      *a = call f(a<sub>1</sub>, ..., a<sub>n</sub>)*

(Note: this TAC design abstract the representation of parameter passing, stack frames, *etc.* These details will emerge when doing machine code generation.)