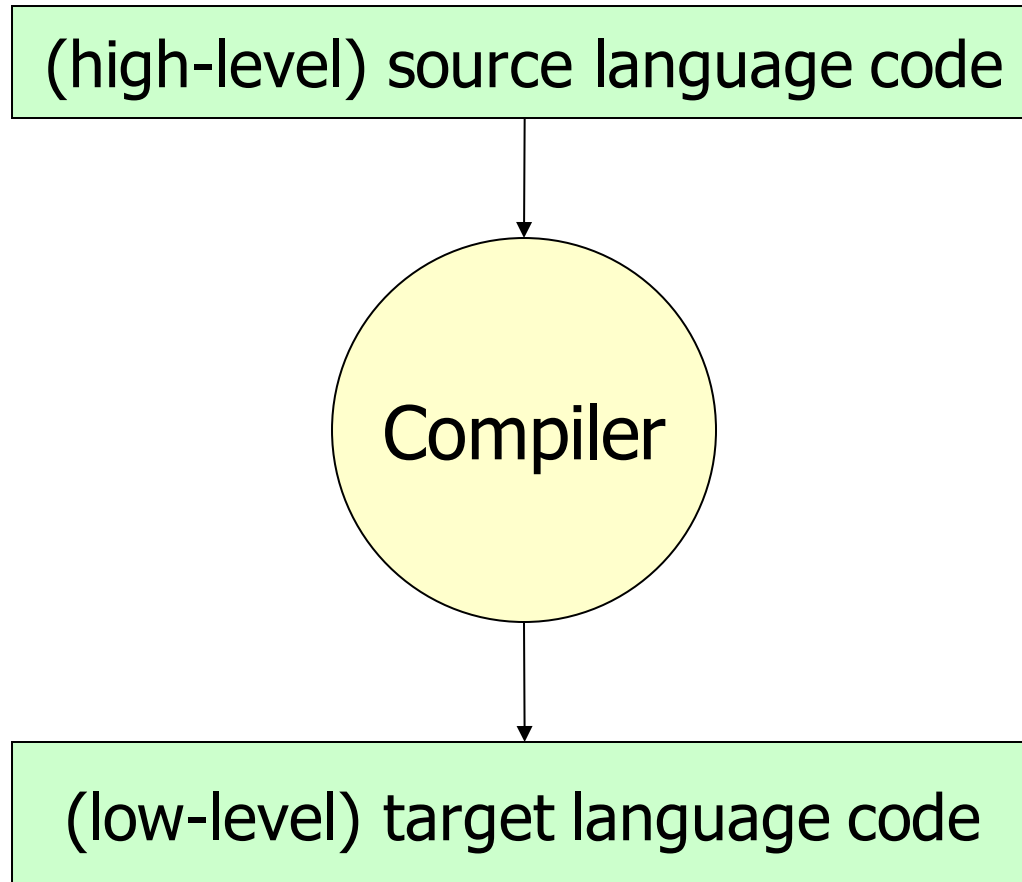


CS 301:
Compiler and Runtime System Design

What is a compiler?

What is a runtime system?

A compiler is a program translator.



A compiler is a program translator.

```
int expr(int n) {  
    int d = 4 * n * n * (n + 1) * (n + 1);  
    return d;  
}
```

```
lda $30,-32($30)          addq $3,1,$4  
stq $26,0($30)            mull $2,$4,$2  
stq $15,8($30)            ld1 $3,16($15)  
bis $30,$30,$15           addq $3,1,$4  
bis $16,$16,$1            mull $2,$4,$2  
stl $1,16($15)            stl $2,20($15)  
lds $f1,16($15)           ld1 $0,20($15)  
sts $f1,24($15)           br $31,$33  
ld1 $5,24($15)            $33:  
bis $5,$5,$2              bis $15,$15,$30  
s4addq $2,0,$3            ldq $26,0($30)  
ld1 $4,16($15)            ldq $15,8($30)  
mull $4,$3,$2             addq $30,32,$30  
ld1 $3,16($15)            ret $31,($26),1
```

A compiler is a program improver.

Unoptimized Code

```
    lda $30,-32($30)
    stq $26,0($30)
    stq $15,8($30)
    bis $30,$30,$15
    bis $16,$16,$1
    stl $1,16($15)
    lds $f1,16($15)
    sts $f1,24($15)
    ldl $5,24($15)
    bis $5,$5,$2
    s4addq $2,0,$3
    ldl $4,16($15)
    mull $4,$3,$2
    ldl $3,16($15)
    addq $3,1,$4
    mull $2,$4,$2
    ldl $3,16($15)
    addq $3,1,$4
    mull $2,$4,$2
    stl $2,20($15)
    ldl $0,20($15)
    br $31,$33
$33:
    bis $15,$15,$30
    ldq $26,0($30)
    ldq $15,8($30)
    addq $30,32,$30
    ret $31,($26),1
```

Optimized Code

```
s4addq $16,0,$0
mull $16,$0,$0
addq $16,1,$16
mull $0,$16,$0
mull $0,$16,$0
ret $31,($26),1
```

A compiler is a program checker.
A compiler is a programmer's assistant.

Static program analysis

Type checking

Scope checking

Control flow checking

Instrumentation for dynamic program analysis

A runtime system is
a compiler's co-conspirator.

(high-level) source language code



Compiler



(low-level) target language code

Runtime
System



A runtime system is

{ a program translator, a program improver, a program checker,
a programmer's assistant, a compiler's co-conspirator }

**Memory management +
Garbage Collection (GC)**

Dynamic program analysis

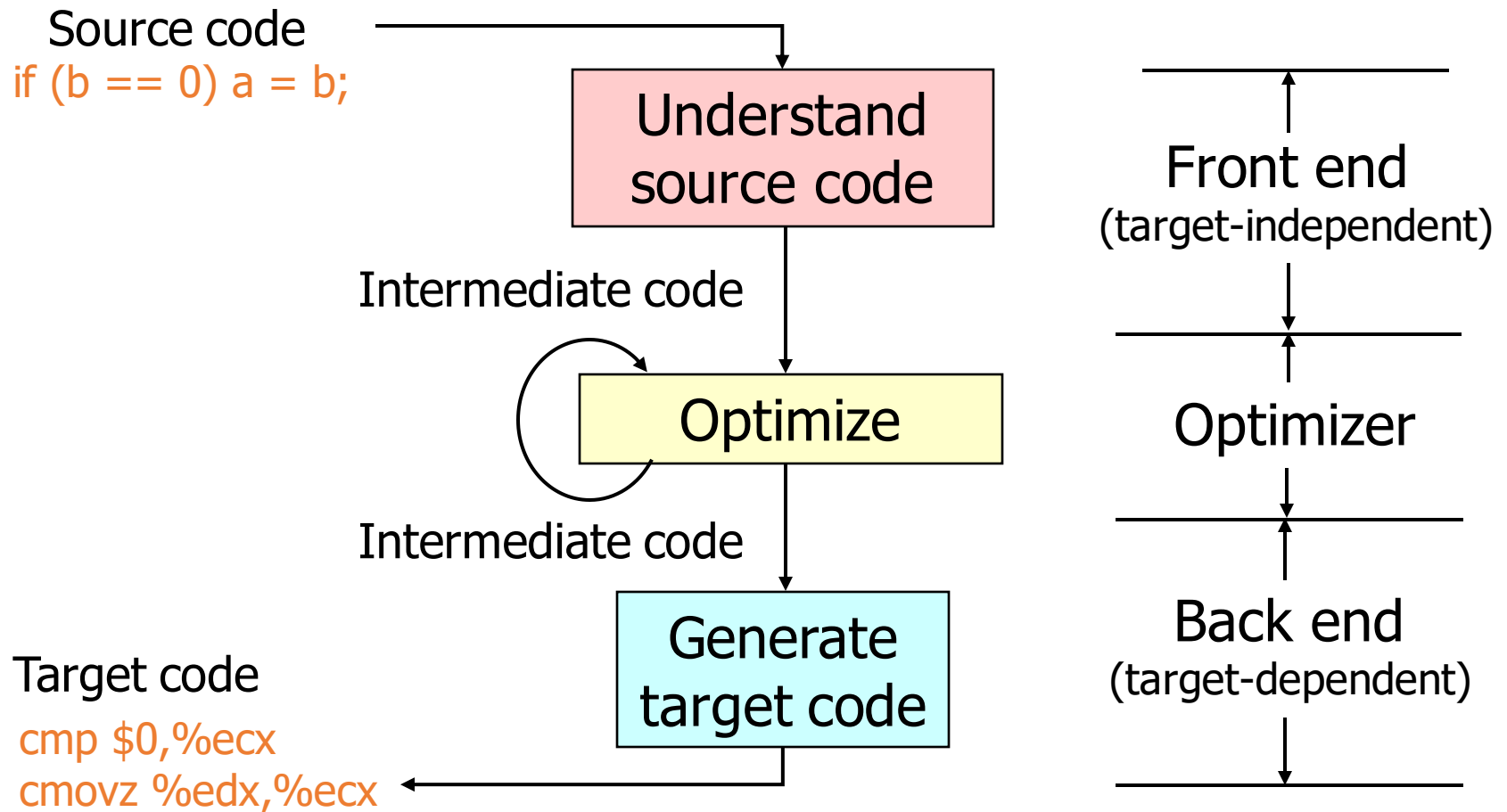
Just in Time compilation (JIT)

Feedback-directed optimization

Profiling + profile guided optimization (PGO)

Resource management

Simplified Compiler Structure



Compiler Front End

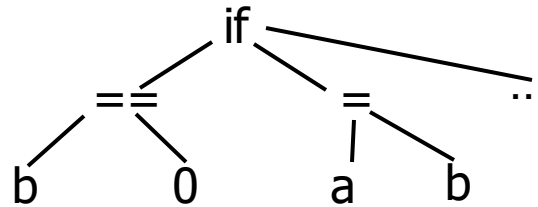
Source code
(character stream)

if (b == 0) a = b;

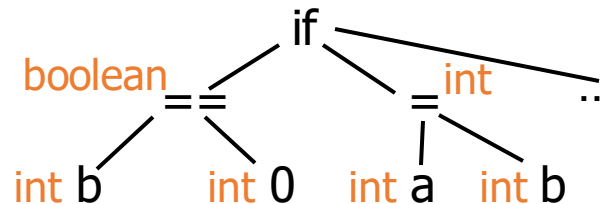
Token
stream

if	(b	==	0)	a	=	b	;
----	---	---	----	---	---	---	---	---	---

Abstract syntax
tree (AST)



Decorated
AST



Lexical Analysis

(Lexing)

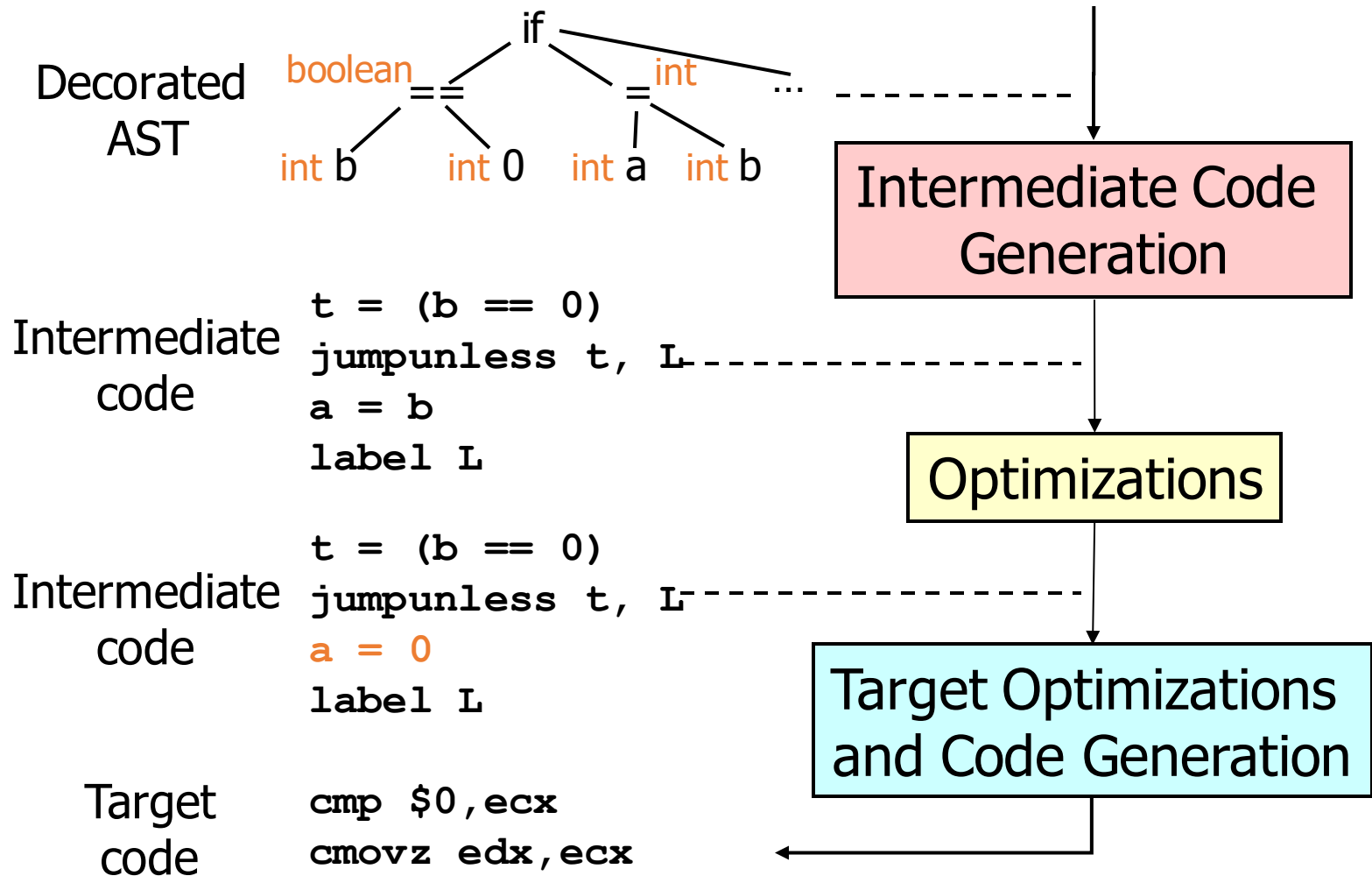
Syntax Analysis

(Parsing)

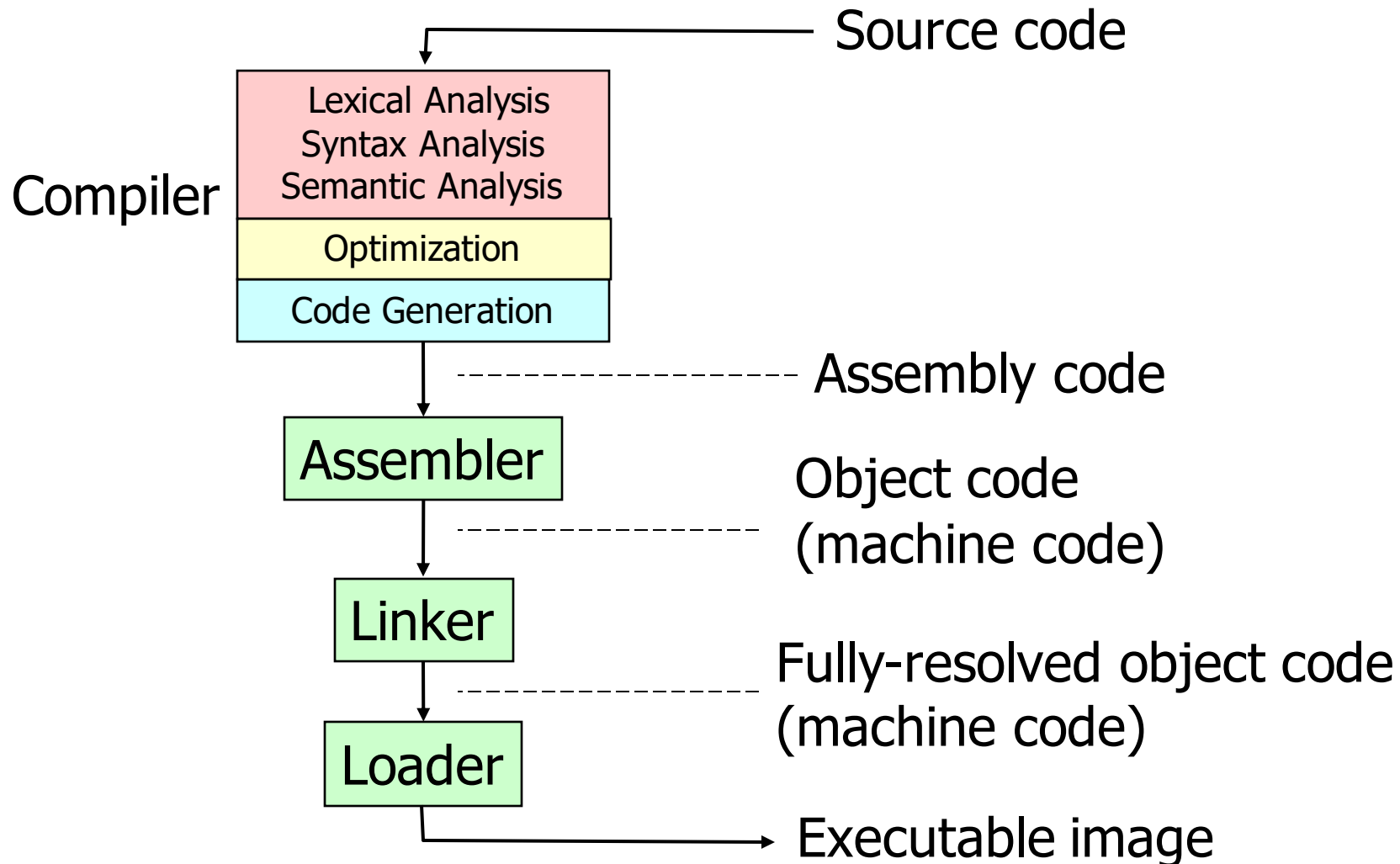
Semantic Analysis

(Name Resolution, Type Checking)

Compiler Middle / Back End



Zooming out: compiling to machine code...



Course Material

- Theoretical Foundations
 - Implementation
 - Synthesis
 - Programming Languages (251), Theory (235), Algorithms (231), Computer Systems (240), Software Construction
 - Semester project: full compiler
 - for small statically typed language
 - mostly from scratch
 - implemented in Scala
 - other tools: git, IntelliJ, gcc, gdb, ...
 - Tutorial
- Project Groups
 - groups of 3 (or 2)
 - you choose, I can help
 - work with new people on the Tiny compiler this week
 - Collaboration and Honor Code

Tutorial Meetings

I form trios with your schedule/people input, end of this week.

Weekly assignment for 1-hour small-group meeting with me.

Preparation: Read, work on all problems.

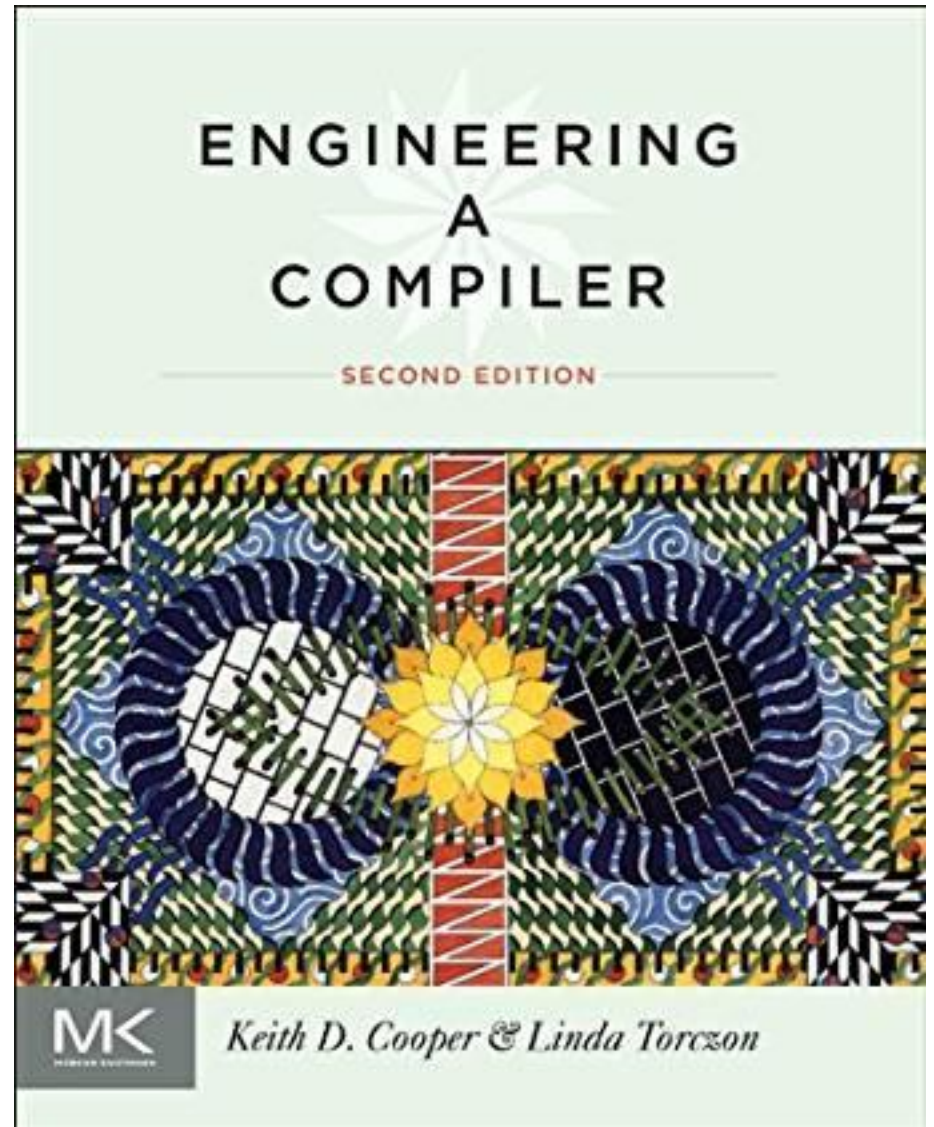
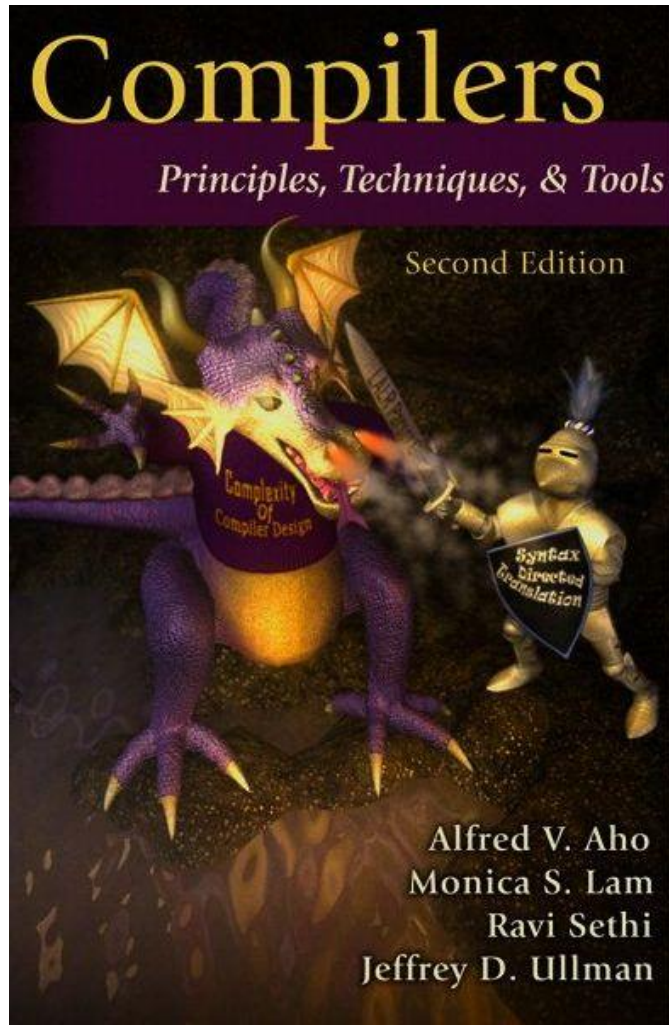
Participation: Discuss any of the problems in detail.

- work through solution on board
- extend problem in new directions
- explain where you got stuck and why

Reviews: Write clean solutions to a few questions I choose.

- in LaTeX
- due 24 hours after your meeting
- better prep before meeting => fewer written solutions after

Books



Weekly Schedule

Tutorial Day: Monday/Tuesday/Wednesday

- tutorial meetings
- next tutorial assignment posted
- Tuesday scheduled class time not required after today
 - May be used for tutorial
 - Otherwise: guaranteed project group availability, code review, etc.

Lab/Work Day: Friday scheduled class time (+ drop-in hours after)

- Introduce and start new project stages/checkpoints
- Project work time
- Code reviews
- Occasional mini-lectures framing upcoming tutorial material
- Project checkpoints / due dates

Other Days:

- Code reviews?
- Drop-in hours, appointments

Tiny compiler!

Source code
(character stream)

if (b == 0) a = b;

Lexical Analysis

Token
stream

if	(b	==	0)	a	=	b	;
----	---	---	----	---	---	---	---	---	---

Identifiers:

x **y11** **elsen** **_i00**

Integers:

2 **1000** **5L**

Floating point:

2.0 **.02** **1.** **1e5** **0.e-10**

Strings:

"x" **"She said, \"Hey!\""**

Comments:

/* don't change this */ **// or this**

Keywords:

if **else** **while** **break**

Symbols:

+ ***** **{** **}** **=** **<** **<<** **==** **[** **]** **>=**

Regular Expressions (*regex*, *RE*)

A **language** is a set of words: { SCI, KSC }, { a,b,c,d,... }

A **regular expression** describes a (regular) language

abab **a|b** **(a|b)*** **[1-9][0-9]*** **[a-z][a-z0-9]***

Regular expression primitives:

a	ordinary character stands for itself
ϵ	the empty string (<i>epsilon</i>)
R S	either R or S (<i>alternation</i>), where R,S are REs
RS	R followed by S (<i>concatenation</i>)
R*	R repeated 0 or more times (<i>Kleene star</i>)

$L(R)$ = the language defined by regular expression R

$L(\mathbf{a(SCI|KSC)}) = \{ aSCI, aKSC \}$

$L(\mathbf{[1-9][0-9]*}) = \{ 1,2,3,4,5,6,7,8,9,10,11,12,13,... \}$

Regular Expression Extensions

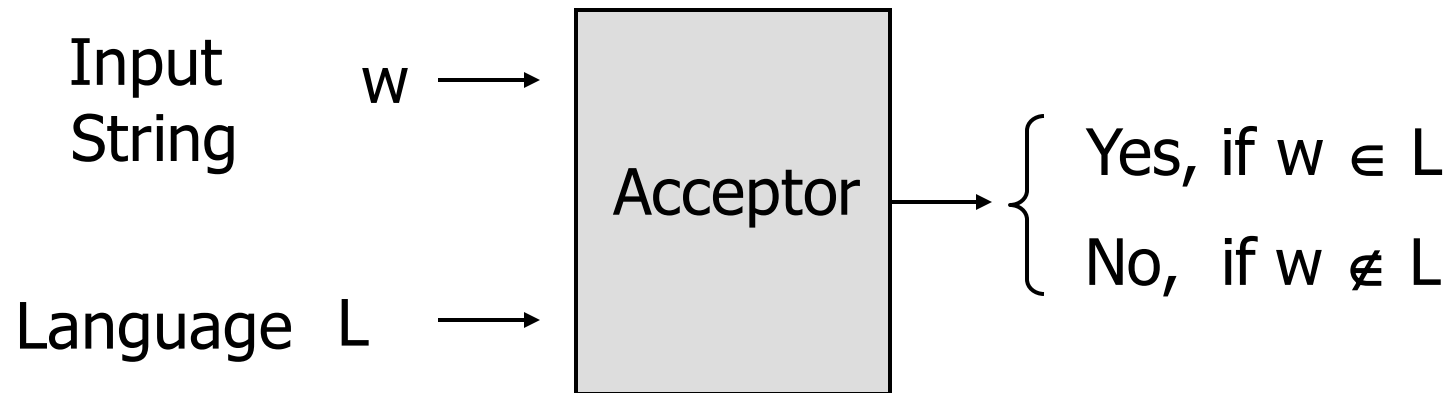
If R is a regular expressions, so are:

$R?$	$= \varepsilon \mid R$	<i>(zero or one R)</i>
R^+	$= RR^*$	<i>(one or more R's)</i>
(R)	$= R$	<i>(no effect: grouping)</i>
$[abc]$	$= a \mid b \mid c$	<i>(any of the listed characters)</i>
$[a-e]$	$= a \mid b \mid \dots \mid e$	<i>(character ranges)</i>
$[^ab]$	$= c \mid d \mid \dots$	<i>(any character except the listed characters)</i>
$.$	$= a \mid b \mid c \mid d \mid \dots$	<i>(any character except newline)</i>
$\backslash .$	$=$	<i>a literal . character</i>
$\backslash *$	$=$	<i>a literal * character</i>
$\backslash \backslash$	$=$	<i>a literal \ character</i>
\dots		

Acceptors:

(a.k.a. recognizers)

Abstract machines that determine if an input string belongs to a language, answering Yes/No.



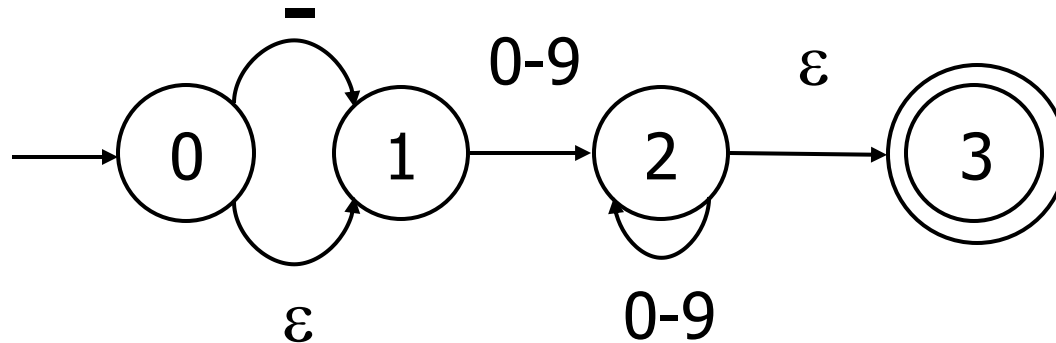
Finite Automata:

acceptors for languages described by regular expressions

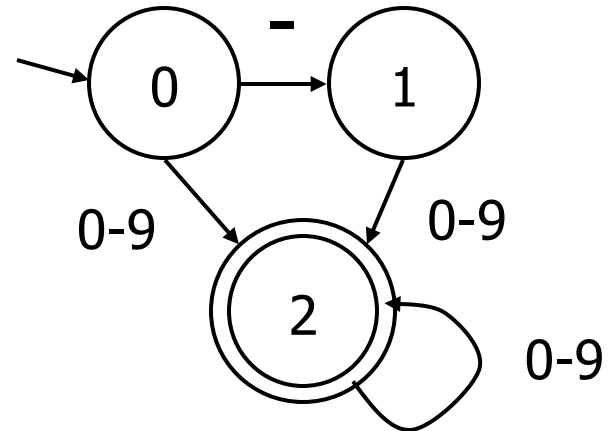
Finite Automata

Regular Expression: $(-|\epsilon)[0-9][0-9]^*$

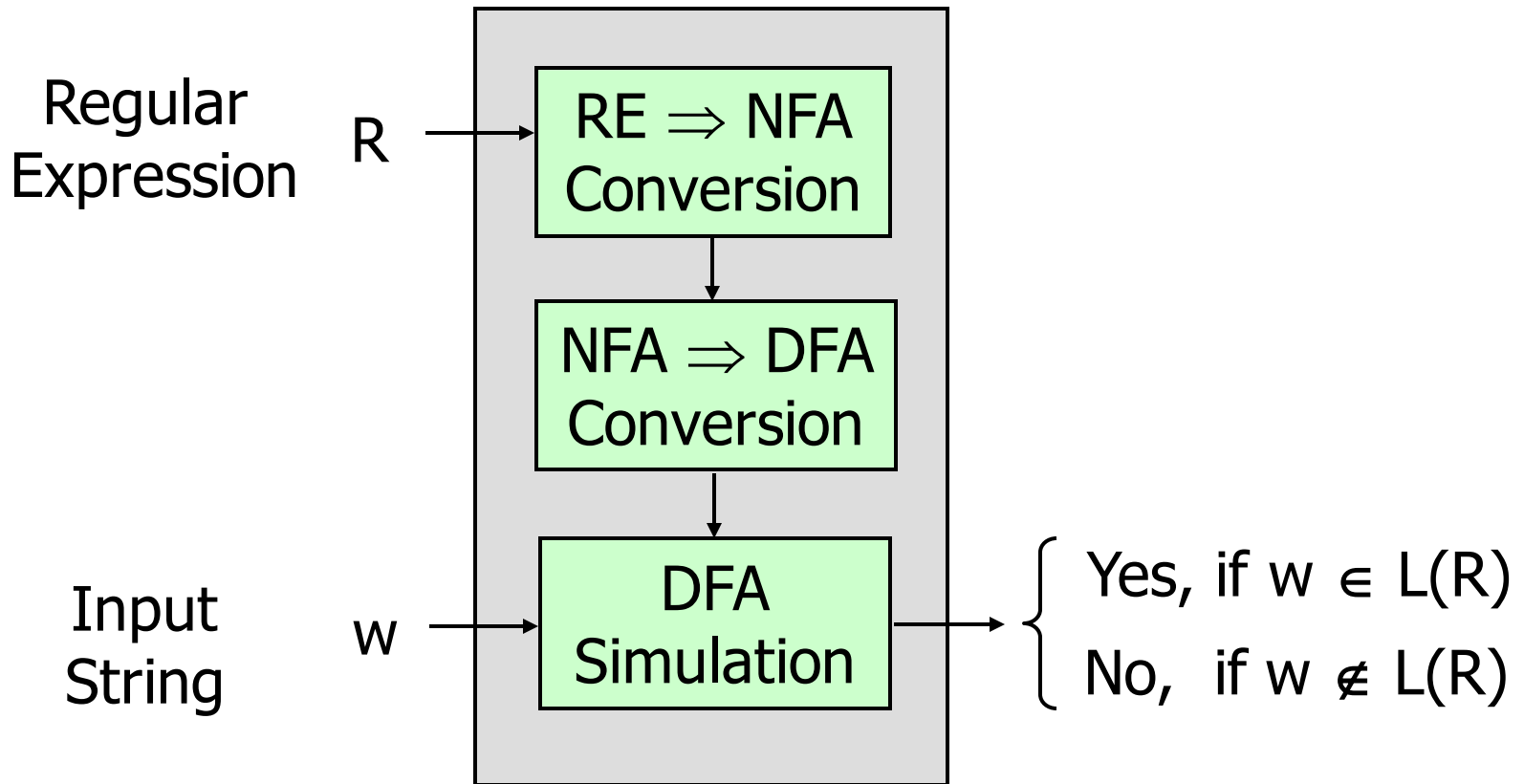
Nondeterministic Finite Automaton:



Deterministic Finite Automaton:



Building an acceptor for a regular expression:



Example Lexer Generator Specification

%%

```
digits      = 0|[1-9][0-9]*
letter      = [A-Za-z]
identifier  = {letter}({letter}|[0-9_])*
whitespace  = [\ \t\n\r]+
```

%%

```
{whitespace} { /* discard */ }
{digits}      { return new Token(INT, parseInt(yytext())); }
"if"          { return new Token(IF, yytext()); }
"while"       { return new Token(WHILE, yytext()); }
...
{identifier}  { return new Token(ID, yytext()); }
```